

# Subexponential Time Algorithms via Concentration Bounds

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# Outline

- 1 Althofer's Lemma
- 2 Extension to Non-linear Mappings
- 3 Market Equilibrium

# Motivation

- Computation of equilibria in 2-player games is a hard problem.
- However, if strategies are **sparse** and **uniform** then the search space is sub-exponential.
- Althofer Lemma guarantees existence of such a grid in strategy space of players.
- We will give a novel application of this sparsification technique in the second half.

# Sparse Solutions to Linear Systems: Althofer's Lemma

- Let  $p = (p_1, \dots, p_n)$  be any probability vector, i.e.,  $\sum_{i=1}^n p_i = 1, p_i \geq 0$ .
- Let  $A$  be any  $m \times n$  matrix with all entries between 0 and 1. Let  $A_i$  be the  $i^{\text{th}}$  row of  $A$ .
- Then the linear transform  $Ap$  can be approximated by  $Aq$ , where probability vector  $q$  is:

approximate representation of  $p$

$$|A_i \cdot (p - q)| \leq \epsilon, i \in \{1, \dots, m\}.$$

sparse at most  $k = \frac{\log 2m}{2\epsilon^2}$  entries are non-zero.

uniform all entries are of form  $q_i = \frac{k_i}{k}$ ,  $k_i$  is integral

# Proof Outline

- Construct a sparse representation of  $p$  by sampling.
- Use Hoeffding bound and union bound to prove that with non-zero probability, the constructed representation is an approximate representation of  $p$ .
- This proved existence of a sparse representation.

# Proof Outline : Construction of Sparse Representation

- Let  $p = (p_1, \dots, p_n)$  be a given probability vector.
- Sample  $k$  times with replacement from set  $\{1, \dots, n\}$  with probability of sampling  $i$  given by  $p_i$ .
- Let  $S$  be the multiset representing the result of the above sampling.
- Let  $n_i$  be the number of times  $i$  occurs in  $S$ .
- Then  $q_i = \frac{n_i}{k}$ .

# Proof Outline : Properties of Sparse Representation

Probability vector  $q = (q_1, \dots, q_n)$ :

is sparse at most  $k$  non-zero entries.

is uniform all entries are multiples of  $\frac{1}{k}$ .

preserves dot products  $E(a \cdot q) = a \cdot p$ .

# Proof Outline

- To prove existence of sparse  $q$ , we need to choose appropriate  $k$ .
- Dot products are preserved in expectation, therefore, the linear transform  $Ap$  is preserved in expectation.
- We use Hoeffding bound and union bound to choose appropriate value of  $k$ .