

# Announcements

- Program #2
  - Due next Thursday (3/3/16)

# Writers Have Priority

## reader

```
repeat
  P(z);
  P(rsem);
  P(x);
  readcount++;
  if (readcount == 1) then
    P(wsem);

  V(x);
  V(rsem);
V(z);
readunit;
P(x);
  readcount- -;
  if readcount == 0 then
    V (wsem)

V(x)
forever
```

## writer

```
repeat
  P(y);
  writecount++;
  if writecount == 1 then
    P(rsem);

  V(y);
  P(wsem);
writeunit
  V(wsem);
  P(y);
  writecount--;
  if (writecount == 0) then
    V(rsem);

  V(y);
forever;
```

# Notes on readers/writers with writers getting priority

Semaphores  $x, y, z, wsem, rsem$  are initialized to 1

```
P(z);  
  P(rsem);  
  P(x);  
    readcount++;  
    if (readcount==1) then  
      P(wsem);  
  V(x);  
  V(rsem);  
V(z);
```



readers queue up on semaphore  $z$ ; this way only a single reader queues on  $rsem$ . When a writer signals  $rsem$ , only a single reader is allowed through

# Sample Synchronization Problem

- Class Exercise:
  - **CMSC 412 Midterm #1 (Spring 1998) Q#3**
- Went over master solution
- Variables:
  - Semaphore mutex = 1
  - Semaphore writer = 0
  - Semaphore reader = 0
  - int nReader = 0
  - int nWriter = 0
  - int wReader = 0
  - int wWriter = 0

# Sample Synchronization Problem

- **Class Exercise:**
  - **CMSC 412 Midterm #1 (Spring 1998) Q#3**
- Solve a variation of the readers-writers problem, in which multiple writers can write at the same time. Specifically, there are readers and writers. Up to 5 reads at the same time are allowed, but only one write at the same time are allowed. A read and a write at the same time is not allowed. Provide a solution using semaphores with the following properties:
  - no busy waiting.
  - starvation-free (i.e. a continuous stream of readers does not starve writers, and vice versa) is desirable but not compulsory (but you will lose some points).
  - you cannot use process ids and you cannot have a separate semaphore for every process.