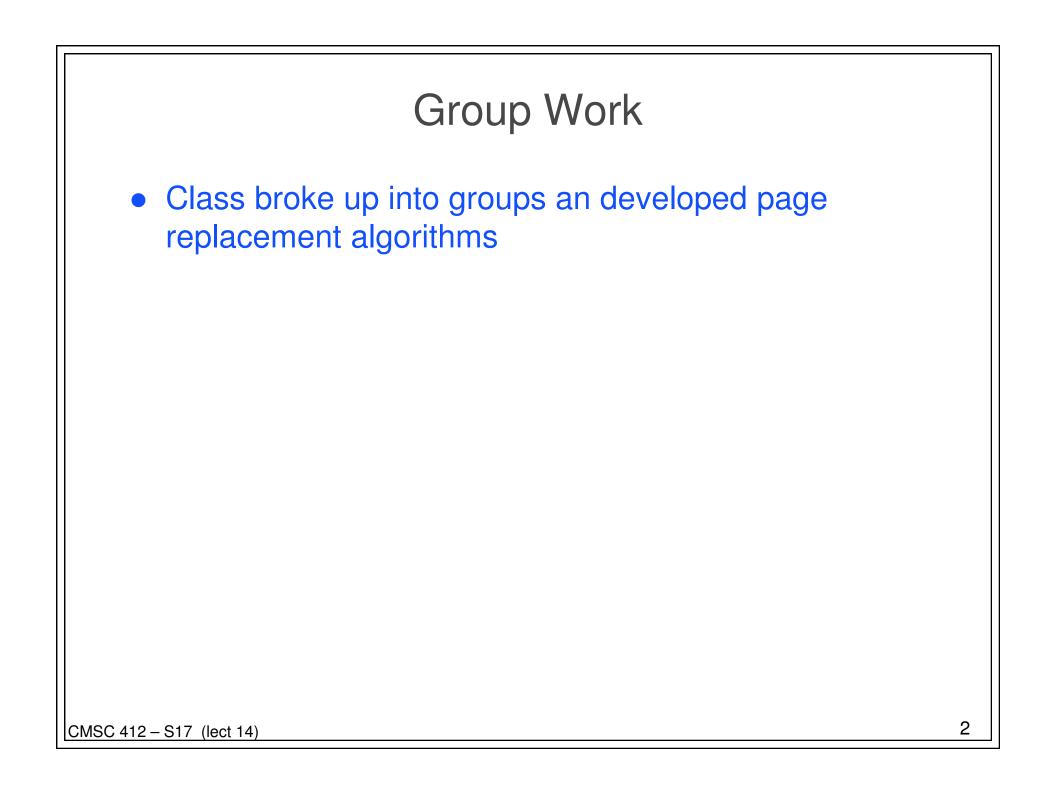
Announcements

- Project #4 teams emailed out
- Git groups setup
- Reminder about re-grade deadline
 - Thursday March 30th (11:00 AM)
- Reading Chapter 10 (in 8th Ed)



What happens when we fault and there are no more physical pages?

- Need to remove a page from main memory
 - if it is "dirty" we must store it to disk first.
 - dirty pages have been modified since they were last stored on disk.
- How to we pick a page?
 - Need to choose an appropriate algorithm
 - should it be global?
 - should it be local (one owned by the faulting process)

Page Replacement Algorithms

• FIFO

- Replace the page that was brought in longest ago
- However
 - old pages may be great pages (frequently used)
 - number of page faults may increase when one increases number of page frames (discouraging!)
 - called belady's anomaly
 - 1,2,3,4,1,2,5,1,2,3,4,5 (consider 3 vs. 4 frames)

• Optimal

- Replace the page that will be used furthest in the future
- Good algorithm(!) but requires knowledge of the future
- With good compiler assistance, knowledge of the future is sometimes possible

Page Replacement Algorithms

• LRU

- Replace the page that was actually used longest ago
- Implementation of LRU can be a bit expensive
 - e.g. maintain a stack of nodes representing pages and put page on top of stack when the page is accessed
 - maintain a time stamp associated with each page
- Approximate LRU algorithms
 - maintain reference bit(s) which are set whenever a page is used
 - at the end of a given time period, reference bits are cleared