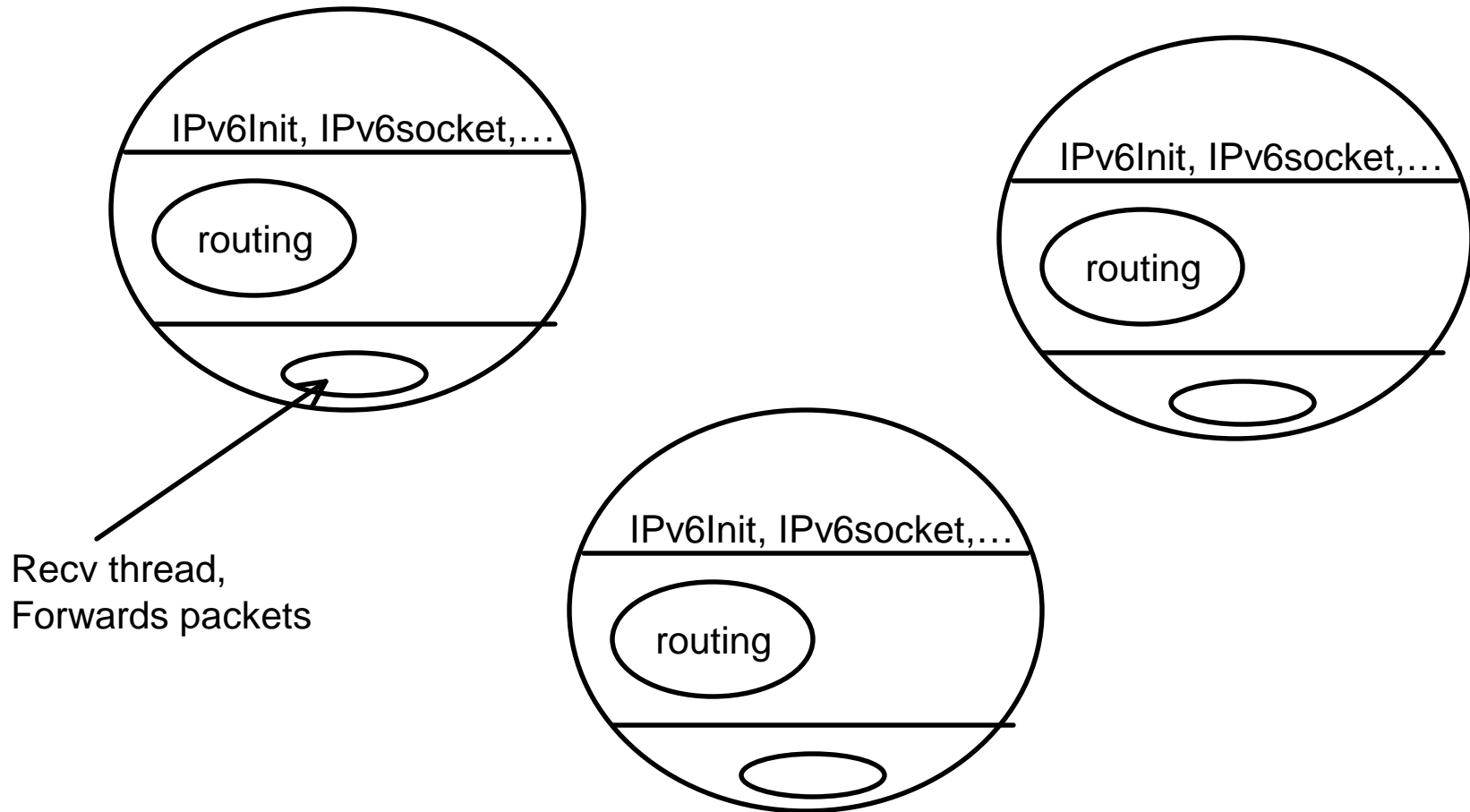


Announcements

- Reading
 - Chapter 4 (4.1-4.2)
- Project #4 is on the web
 - Note policy about project #3 missing components
- Homework #1
 - Due 11/6/01
 - Chapter 6: 4, 12, 24, 37
- Midterm #2
 - 11/8/01 in class

Project #4 Notes



Transmission: Satellites

- **Different Orbits Possible**
 - orbit affects many communication properties
- **Geosynchronous**
 - always over the same spot on the earth
 - 36,000 Km orbit is required
 - only 180 slots possible
 - uses one uplink and one down link frequency
 - large round-trip latencies
- **LEO (Low Earth Orbit)**
 - each satellite keeps moving into and out of range
 - solution: use a large number of satellites
 - sort of like cells, but the cells are the ones moving
 - lower round-trip latency

Sending More Than one Signal At Once

- Called multiplexing
 - original goal of Bell was to MUX multiple telegraph signals
- Time Division Multiplexing
 - everyone gets whole bandwidth
 - but only when its their turn

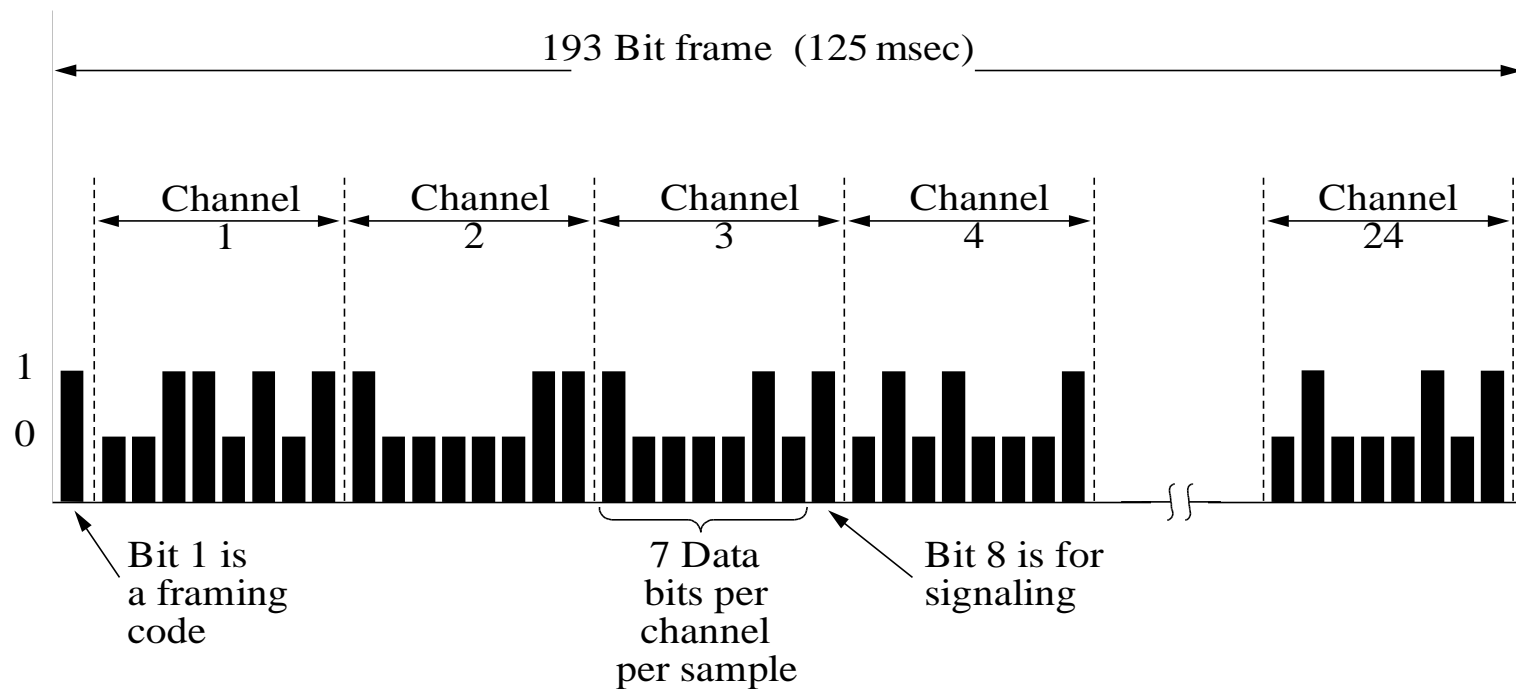


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Frequency Division Multiplexing

- Frequency Division

- everyone gets to talk at once
- but only in their own frequency

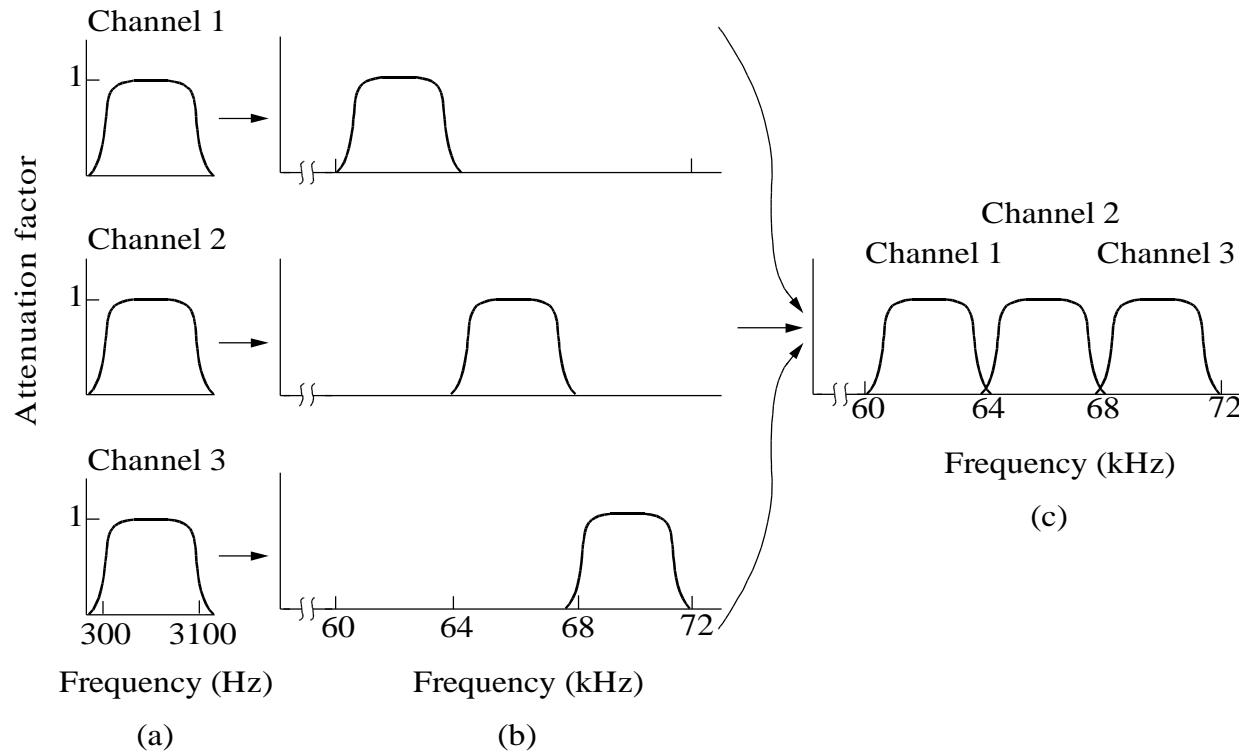


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ATM Switching

- Requirements

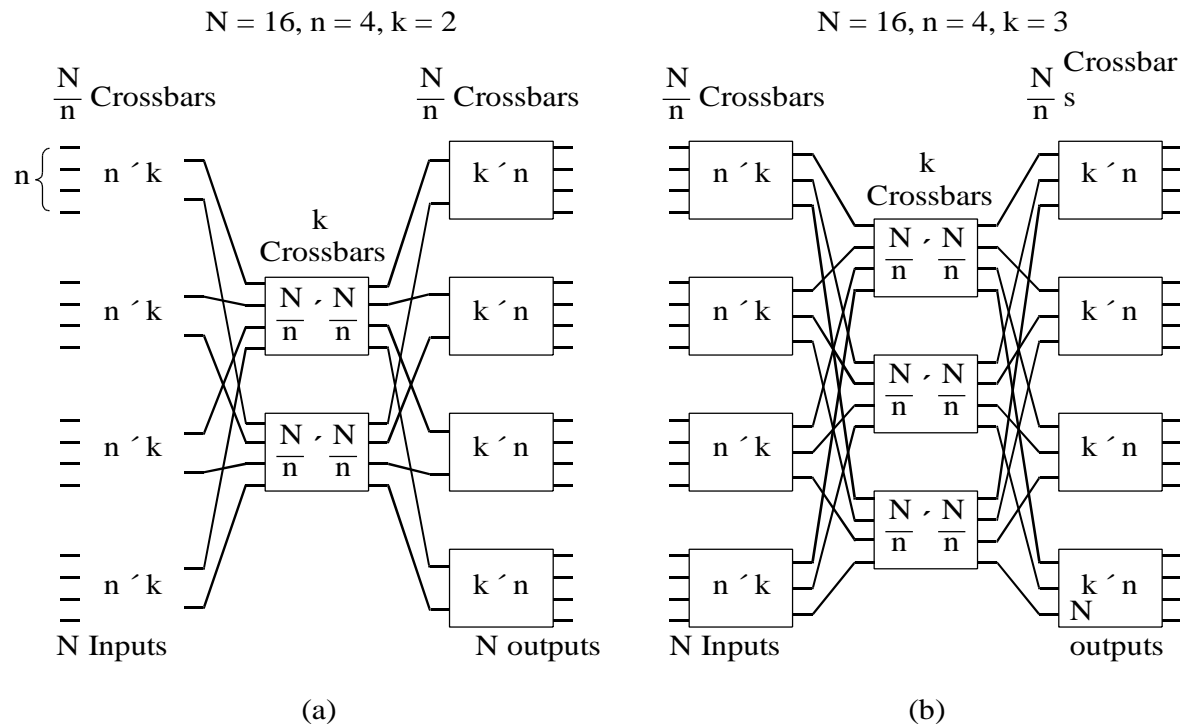
- be able to switch 360,000 cells/sec per input link
- switch cells with as low a discard rate as possible
- never reorder the cells on a virtual circuit

- Issues

- multiple cells destined for the same output at once
 - need to buffer one of them
 - must ensure fairness is maintained
- head-of-line blocking
 - possible that a blocked output is holding up cells that could be delivered

Switching Fabric (space division)

- Cross bars are great, but require $O(n^2)$ wires
- Can use a collection of smaller cross bar switches
 - penalty: a request to connect may **block**



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Batcher-banyan Switching

- Banyan
 - can do a “good” or “poor” job of switching due to collisions
 - if the inputs are sorted, we get performance
- Batcher
 - sorts traffic base on full address of destination
 - compares two colliding packets and uses final destination to select output port
 - requires $O(n \log^2 n)$ nodes (2x2 switching elements)

Title: (Adobe Illustrator (R) Version 5.0 Level 2 Emulation)
Creator: Adobe Illustrator(TM) 5.0
CreationDate: (04/10/93) ()

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Medium Access Layer

- **Broadcast Networks**
 - share a common resource for communication
 - bus, wire, air, etc.
 - need to coordination access to this resource
- **Limits of Static Channel Allocation**
 - suitable for constant rate traffic of similar speeds
 - however, bursty traffic results in poor channel utilization
 - consider one queue vs. separate queues for each person
 - n queues with bursty arrival have mean delay n times
1 queue
- **Dynamic Allocation**
 - only use channel when have something to send
 - need to control access to the channel

Shared Channel Model

- **Station model**
 - N independent stations
 - each wants to send λ frames per second
 - a station may not send another frame until the first is sent
- **Single Channel Assumption**
 - all stations communicate over a single shared channel
- **Collisions: two stations attempt to send at once**
 - neither transmission succeeds
- **Time**
 - continuous time: frame transmissions can start anytime
 - discrete time: clock ensures all sends initiate at the start of a slot
- **Carrier Sense**
 - stations can tell if channel is in use before sending
 - stations must wait to know if channel was in use

Aloha

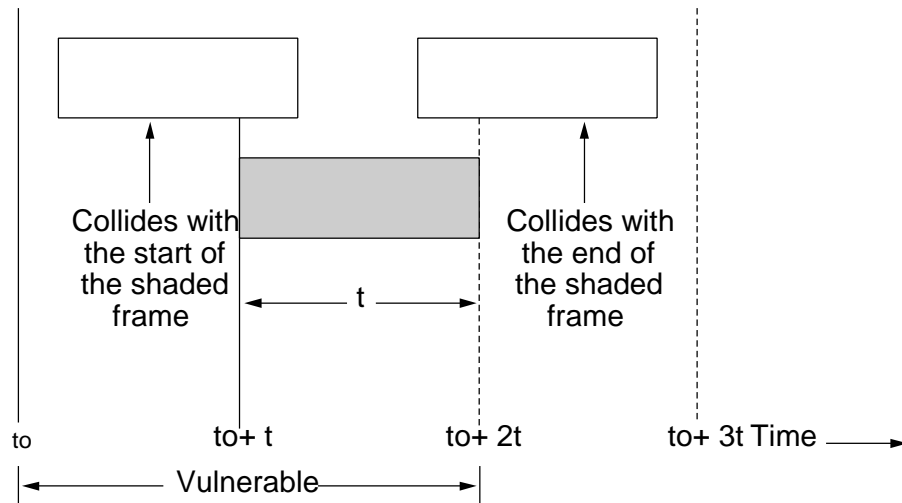
- Stations

- ground based radio stations on islands

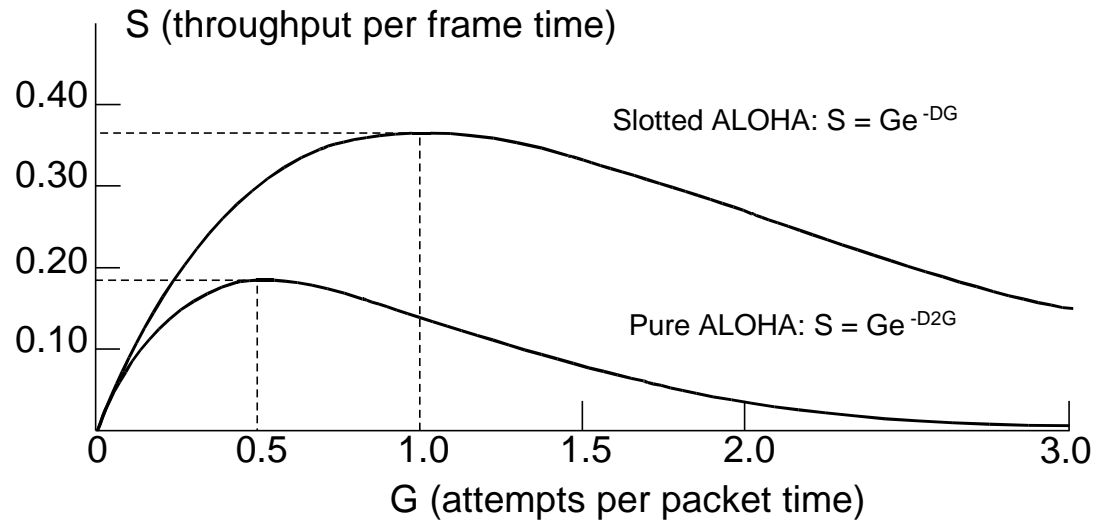
- Pure Aloha

- If send data at will, collisions will happen
- on collision, wait a random amount of time & try again
- use standard, fixed size packets
- what is channel efficiency (assuming Poisson distribution)?
 - S **new** frames per frame time
 - assume G total frames trying to be sent per frame time
 - P_0 = probability of exactly one sender sending
 - $S = G P_0$
 - probability of k frames generated during a frame time
 - $\Pr[k] = G^k e^{-G} / k!$
 - $P_0 = e^{-G} e^{-G} = e^{-2G}$, so $S = G e^{-2G}$

Performance of Aloha



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Aloha (cont.)

- Slotted Aloha

- Use a central clock
- Each station only sends at the start of frame
- Reduces collision window by 1/2
 - $S = G e^{-G}$