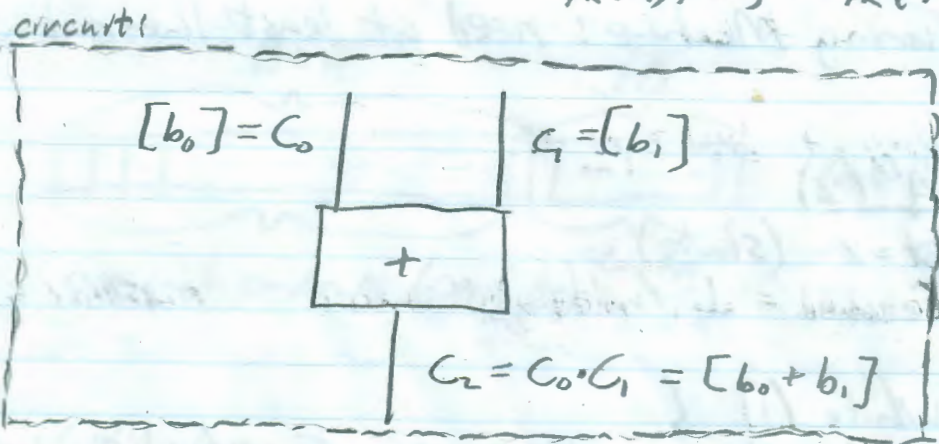


09/30/2013

Last week review:

$$\begin{array}{c} \underline{C} \\ (pk, sk) \leftarrow \text{Gen}(1^n) \end{array} \quad \begin{array}{c} \underline{S} \\ \xrightarrow{pk} \\ \text{Enc}_{pk}(x_1), \dots, \text{Enc}_{pk}(x_\ell) \end{array}$$

$\text{Enc}_{pk}(x_1), \dots, \text{Enc}_{pk}(x_\ell)$



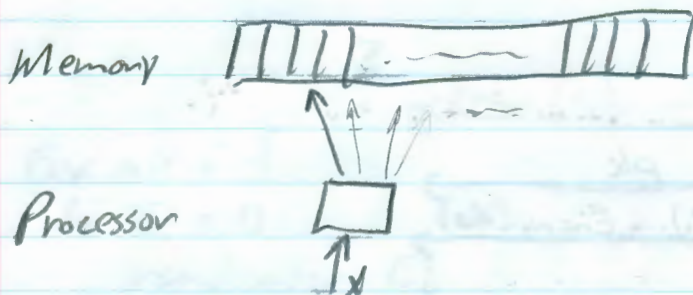
Fully homomorphic encryption (FHE):

$$\begin{array}{c} \underline{C} \\ \xrightarrow{\text{local hom. eval of circuit}} \\ \text{Enc}_{pk}(\text{out}_1), \dots, \text{Enc}_{pk}(\text{out}_\ell) \end{array}$$

+ clients can off load computations to server

- hom. enc is expensive

Secure computation in the RAM model:



Turing Machine: need at least linear time for computation

$A^M(x)$

$st = x$ (start)

Getched = \perp (nothing)

While (1) {

$(st, addr, val) \leftarrow \text{NextInst}(st, \text{Getched})$

if $st == \text{halt}$

output val

else {

$\text{Getched} = M(addr)$

if $(val \neq \text{None})$

$M(addr)$

}

}

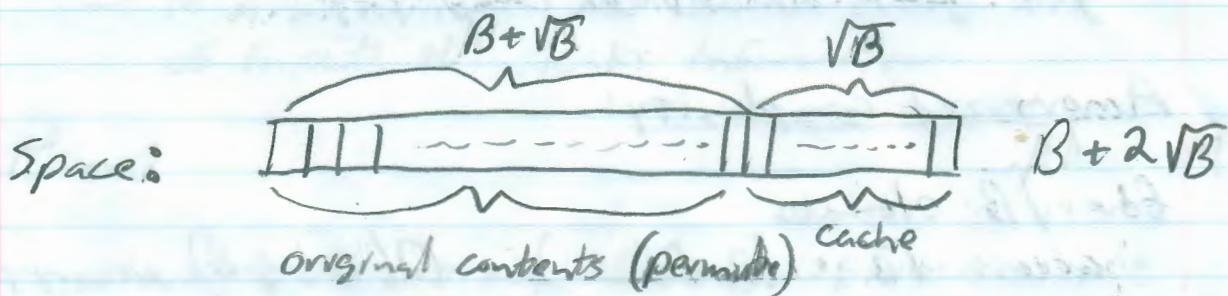
Oblivious RAM (ORAM)

Trivial ORAM:

- to read address v , do linear scan over M (and encrypt contents)

Non-trivial ORAM scheme:

- (Goldreich - Ostrovsky '94)



Initialization:

- Store in CPU key k for pseudorandom permutation over $\{1, \dots, B + \sqrt{B}\}$

- In virtual memory, store $\{(i, M[i])\}_{i=1}^B$
in position $F_k(i)$

- In virtual memory, store $(i, \text{null})_{i=B+1}^{B+\sqrt{B}}$

To read address i do:

- 1) Perform linear scan of cache for element of the form (i, v)
 \rightarrow if found, set $val = v$
- 2) If not found, probe position $F_k(i)$ to get value
If found, probe position $F_k(B + \text{ctr})$, $\text{ctr}++$
- 3) Linear scan of cache to write the (new) value at $M[i]$

\sqrt{B} times
then
refresh

Refresh (every \sqrt{B} steps)

- choose fresh key k'
- (obviously) re-shuffle elements according to k'
- empty cache

$O(B \log B)$ sorting +
 $O(n)$ additional operations

- get good amortized complexity

Amortized complexity:

for \sqrt{B} elements:

access $\sqrt{B} \cdot (2\sqrt{B} \pm 1) + O(B \log B)$ memory problems

$$\frac{O(B \log B)}{\sqrt{B}} = \boxed{O(\sqrt{B} \log B)}$$

cache B^ϵ :

$$\frac{B^\epsilon \cdot O(B^\epsilon)}{B^\epsilon} + O(B \log B) = \boxed{O(B^{1-\epsilon})}$$