

Lecture 25

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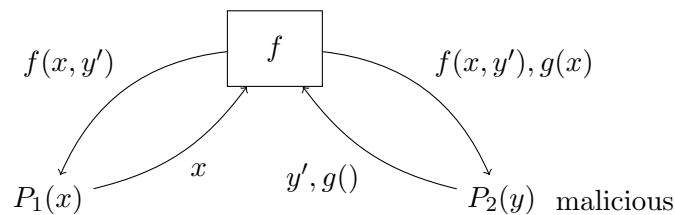
1 Summary

In this lecture we introduced a much more efficient protocol for malicious security given a weaker notation of security. In particular, we talk about efficient GC under 1-bit leakage[1].

In general, we define security by comparing real world to ideal world. When we say a weaker security, we can do the followings:

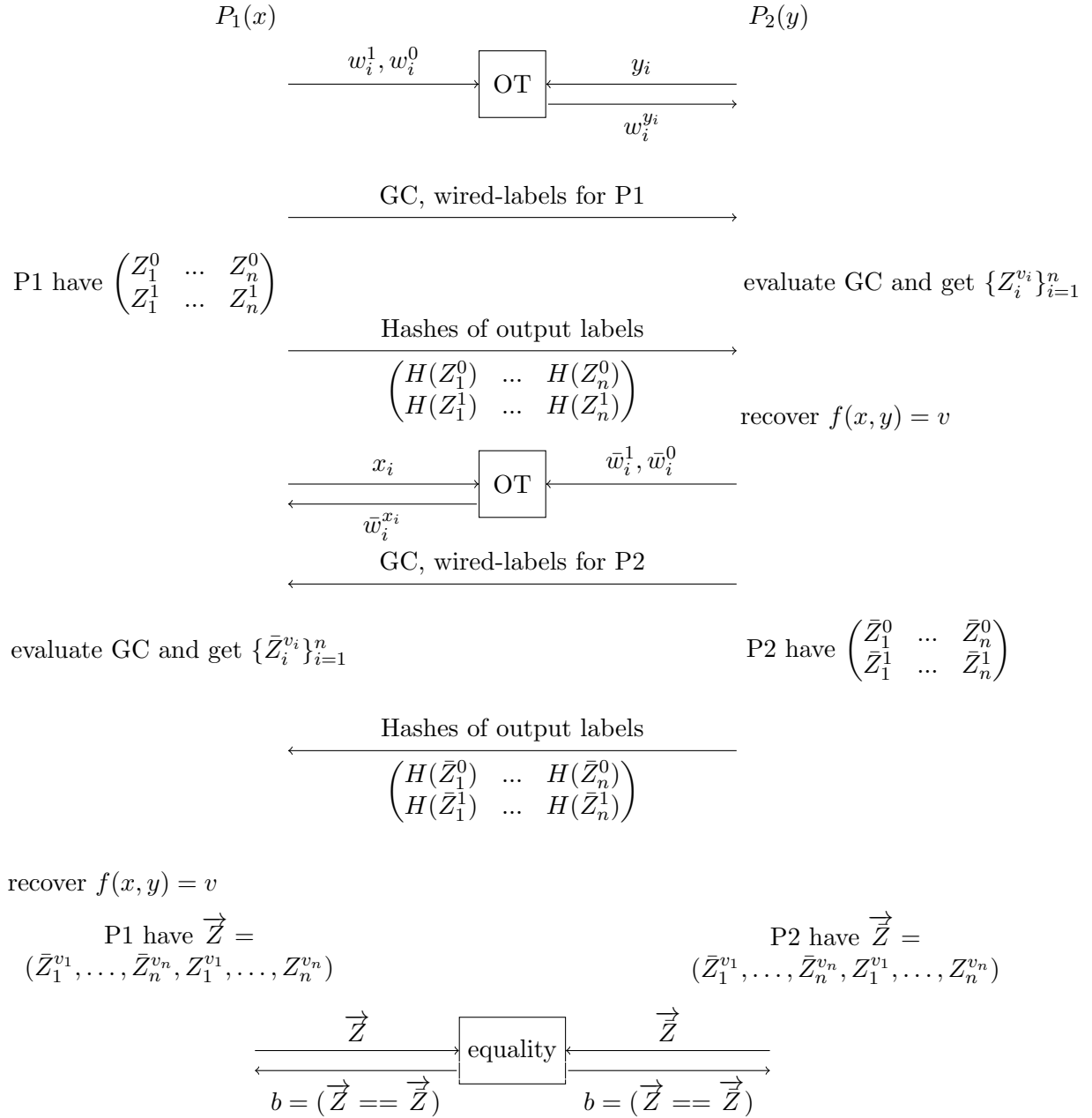
- Weaken the notation of "comparison"
- Weaken the ideal world
 - 1-bit leakage
 - covert security

weaker model of security The malicious party can send a function $g()$ and get $g(x)$ when receiving result from ideal functionality.



2 The protocol for 1-bit-leakage

The protocol is as follows, where Z_i^b is the label for i -th output wire when the value is b .



3 Proof using simulation

WLOG, we assume that P_2 is corrupted and we have simulator S :

- 1) Extract P_2 's input to OT in first phase \Rightarrow this defines input y
- 2) send y to ideal functionality and get back v

- 3) use semi-honest simulation to generate all input-wire labels, Garbled Circuits, to give to P2. We also output $\{Z_i^{v_i}\}_{i=1}^n$. We choose uniformly random for $\{Z_i^{\bar{v}_i}\}_{i=1}^n$ and send the matrix of hashes.
- 4) Extract input wired-labels for P2's circuit from OT; receive GC, input wired-labels and matrix.
Extract P2's input to equality test.
- 5) Define the following $g()$ on input x :
 - use the bits of x to select $\{\bar{w}_i^{x_i}\}$
 - run GC evaluation as P1 would to get v'
 - Define vector \vec{Z} that P1 would use in equality test
 - return 1 iff $\vec{Z} == \vec{Z}'$
- 6) receive $g(x)$ and give it to P2.
- 7) if $g(x) == 0$ or P2 abort, send "abort" to ideal functionality otherwise send "continue".

References

- [1] Huang, Yan, Jonathan Katz, and David Evans. "Quid-pro-quo-tocols: Strengthening semi-honest protocols with dual execution." In *Security and Privacy (SP)*, 2012 IEEE Symposium on, pp. 272-284. IEEE, 2012.