

09/25/2013

GMW vs. Yao

- preprocessing ?

* GMW:

preprocess all OTs

Online phase:

- does not involve any crypto

- has low communication

drawback:

- round complexity linear in circuit depth

overall computation:

- 1 OT per gate

- 4-in-1 H-evaluations per gate

(using load-balancing)

* Yao:

preprocessing

- preprocess OTs

- (preprocess garbled-circuit generation)

online phase:

- garbled-circuit evaluation

- GC generation

⇒ high communication

round-complexity constant

overall computation:

- cost of generating / evaluating garbled gates

- 3-in-1 AES operations per gate

Semi-honest MPC implementations:

- Fairplay (2004)
 - Yao
- Tasty (2010)
 - mixed-mode protocols
 - Yao or homomorphic encryption
- Huang et al. (2011)
 - pipelining
 - significantly more efficient / scalable
 - prog. framework
 - Yao
- Choi et al.
 - GMW
- Schneider-Zohner (2013)
 - GMW vs. Yao

TASTY

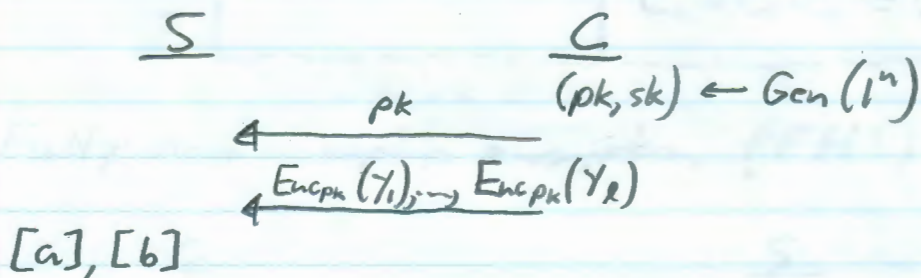
hom. enc:

given encryption of a, b ,
can generate encryption of $a+b$

$$[a], [b] \Rightarrow [a+b]$$

$$[a], r \Rightarrow [r \cdot a]$$

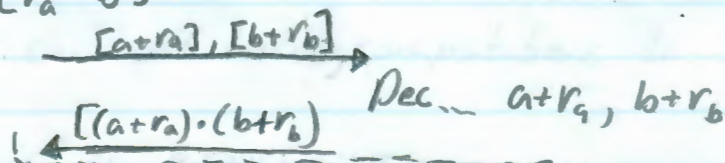
Paillier encryption scheme is homomorphic
over \mathbb{Z}_N



choose random r_a, r_b

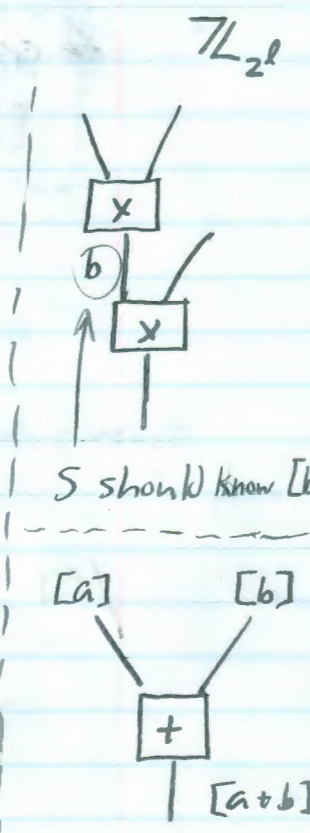
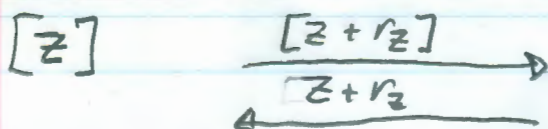
$$[a+r_a], [b+r_b]$$

$$[r_b \cdot a], [r_a \cdot b]$$



$$[(a+r_a) \cdot (b+r_b)] \cdot [-r_a \cdot b] \cdot [-r_b \cdot a] \cdot [-r_a r_b]$$

$$= [a \cdot b]$$



S

C

\tilde{x} : garbled output value

$$\tilde{x} = [(k_1, \lambda_1), \dots, (k_\ell, \lambda_\ell)]$$

Encpk (λ_i), ...

for all i {

if $\pi_i = 0$

randomize C_i

if $\pi_i = 1$

$$C_i = C_i^{-1}$$

}

$$\text{output } [x] = \prod_{i=1}^{\ell} C_i^{z_i}$$

* conversion between modes is expensive