# 15-213 "The Class That Gives CMU Its Zip!"

# Introduction to Computer Systems

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#### **Topics:**

- **Theme**
- **■** Five great realities of computer systems
- How this fits within CS curriculum

### **Course Theme**

Abstraction is good, but don't forget reality!

#### Courses to date emphasize abstraction

- Abstract data types
- Asymptotic analysis

#### These abstractions have limits

- Especially in the presence of bugs
- Need to understand underlying implementations

#### **Useful outcomes**

- Become more effective programmers
  - Able to find and eliminate bugs efficiently
  - Able to tune program performance
- Prepare for later "systems" classes in CS & ECE
  - Compilers, Operating Systems, Networks, Computer Architecture, Embedded Systems

#### Int's are not Integers, Float's are not Reals

#### **Examples**

- Is  $x^2 \ge 0$ ?
  - Float's: Yes!
  - Int's:
    - » 40000 \* 40000 --> 1600000000
    - » 50000 \* 50000 --> ??
- ls(x + y) + z = x + (y + z)?
  - Unsigned & Signed Int's: Yes!
  - Float's:
    - » (1e20 + -1e20) + 3.14 --> 3.14
    - » 1e20 + (-1e20 + 3.14) --> ??

# **Computer Arithmetic**

#### Does not generate random values

Arithmetic operations have important mathematical properties

#### Cannot assume "usual" properties

- Due to finiteness of representations
- Integer operations satisfy "ring" properties
  - Commutativity, associativity, distributivity
- Floating point operations satisfy "ordering" properties
  - Monotonicity, values of signs

#### **Observation**

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- Need to understand which abstractions apply in which contexts
- Important issues for compiler writers and serious application programmers
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#### You've got to know assembly

#### Chances are, you'll never write program in assembly

Compilers are much better & more patient than you are

# Understanding assembly key to machine-level execution model

- Behavior of programs in presence of bugs
  - High-level language model breaks down
- **Tuning program performance** 
  - Understanding sources of program inefficiency
- Implementing system software
  - Compiler has machine code as target
  - Operating systems must manage process state

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# **Assembly Code Example**

#### **Time Stamp Counter**

- Special 64-bit register in Intel-compatible machines
- Incremented every clock cycle
- Read with rdtsc instruction

#### **Application**

- Measure time required by procedure
  - In units of clock cycles

```
double t;
start_counter();
P();
t = get_counter();
printf("P required %f clock cycles\n", t);
```

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### **Code to Read Counter**

- Write small amount of assembly code using GCC's asm facility
- Inserts assembly code into machine code generated by compiler

```
static unsigned cyc_hi = 0;
static unsigned cyc_lo = 0;

/* Set *hi and *lo to the high and low order bits
  of the cycle counter.

*/
void access_counter(unsigned *hi, unsigned *lo)
{
    asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
        : "=r" (*hi), "=r" (*lo)
        :
        : "%edx", "%eax");
}
```

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## **Code to Read Counter**

```
/* Record the current value of the cycle counter. */
void start counter()
    access counter(&cyc hi, &cyc lo);
/* Number of cycles since the last call to start counter. */
double get counter()
    unsigned ncyc hi, ncyc lo;
    unsigned hi, lo, borrow;
    /* Get cycle counter */
    access counter(&ncyc hi, &ncyc lo);
    /* Do double precision subtraction */
    lo = ncyc lo - cyc lo;
    borrow = lo > ncyc lo;
    hi = ncyc hi - cyc hi - borrow;
    return (double) hi * (1 << 30) * 4 + lo;
```

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# **Measuring Time**

#### **Trickier than it Might Look**

Many sources of variation

#### **Example**

Sum integers from 1 to n

n	Cycles	Cycles/n
100	961	9.61
1,000	8,407	8.41
1,000	8,426	8.43
10,000	82,861	8.29
10,000	82,876	8.29
1,000,000	8,419,907	8.42
1,000,000	8,425,181	8.43
1,000,000,000	8,371,2305,591	8.37

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#### **Memory Matters**

#### Memory is not unbounded

- It must be allocated and managed
- Many applications are memory dominated

#### Memory referencing bugs especially pernicious

■ Effects are distant in both time and space

#### Memory performance is not uniform

- Cache and virtual memory effects can greatly affect program performance
- Adapting program to characteristics of memory system can lead to major speed improvements

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# Memory Referencing Bug Example

```
main ()
{
  long int a[2];
  double d = 3.14;
  a[2] = 1073741824; /* Out of bounds reference */
  printf("d = %.15g\n", d);
  exit(0);
}
```

	Alpha	MIPS	Linux
-g	5.30498947741318e-315	3.1399998664856	3.14
-0	3.14	3.14	3.14

(Linux version gives correct result, but implementing as separate function gives segmentation fault.)

# **Memory Referencing Errors**

#### C and C++ do not provide any memory protection

- Out of bounds array references
- Invalid pointer values
- Abuses of malloc/free

#### Can lead to nasty bugs

- Whether or not bug has any effect depends on system and compiler
- Action at a distance
  - Corrupted object logically unrelated to one being accessed
  - Effect of bug may be first observed long after it is generated

#### How can I deal with this?

- Program in Java, Lisp, or ML
- Understand what possible interactions may occur
- Use or develop tools to detect referencing errors

# **Memory Performance Example**

#### Implementations of Matrix Multiplication

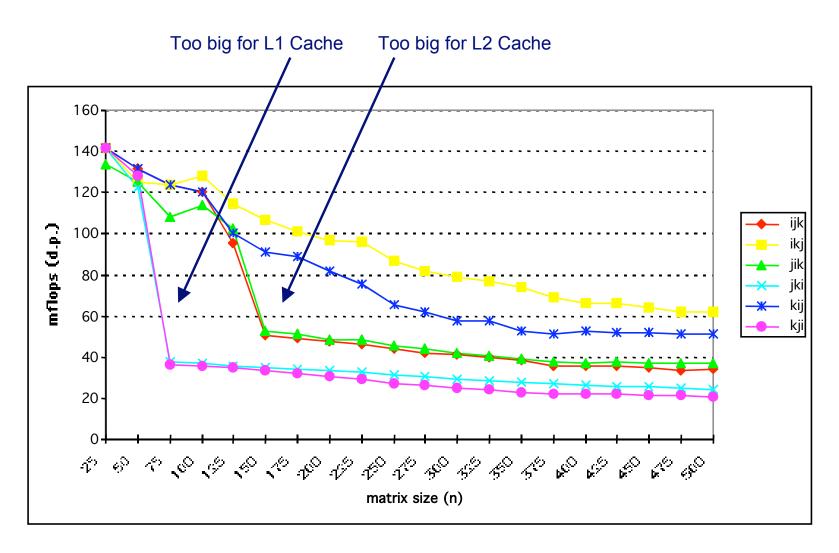
Multiple ways to nest loops

```
/* ijk */
for (i=0; i<n; i++) {
  for (j=0; j<n; j++) {
    sum = 0.0;
    for (k=0; k<n; k++)
      sum += a[i][k] * b[k][j];
    c[i][j] = sum;
}
}</pre>
```

```
/* jik */
for (j=0; j<n; j++) {
  for (i=0; i<n; i++) {
    sum = 0.0;
    for (k=0; k<n; k++)
      sum += a[i][k] * b[k][j];
    c[i][j] = sum
  }
}</pre>
```

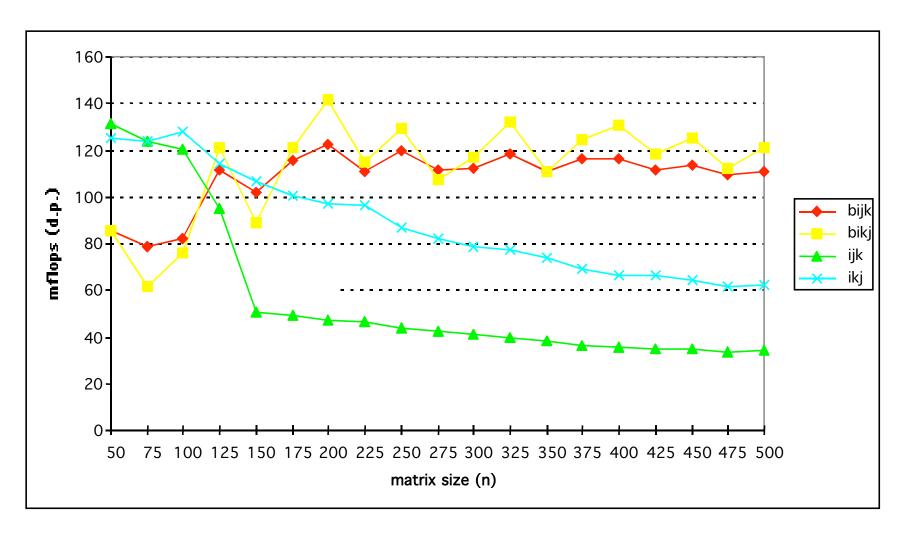
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# Matmult Performance (Alpha 21164)



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# Blocked matmult perf (Alpha 21164)



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# There's more to performance than asymptotic complexity

#### **Constant factors matter too!**

- Easily see 10:1 performance range depending on how code written
- Must optimize at multiple levels: algorithm, data representations, procedures, and loops

#### Must understand system to optimize performance

- How programs compiled and executed
- How to measure program performance and identify bottlenecks
- How to improve performance without destroying code modularity and generality

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#### Computers do more than execute programs

#### They need to get data in and out

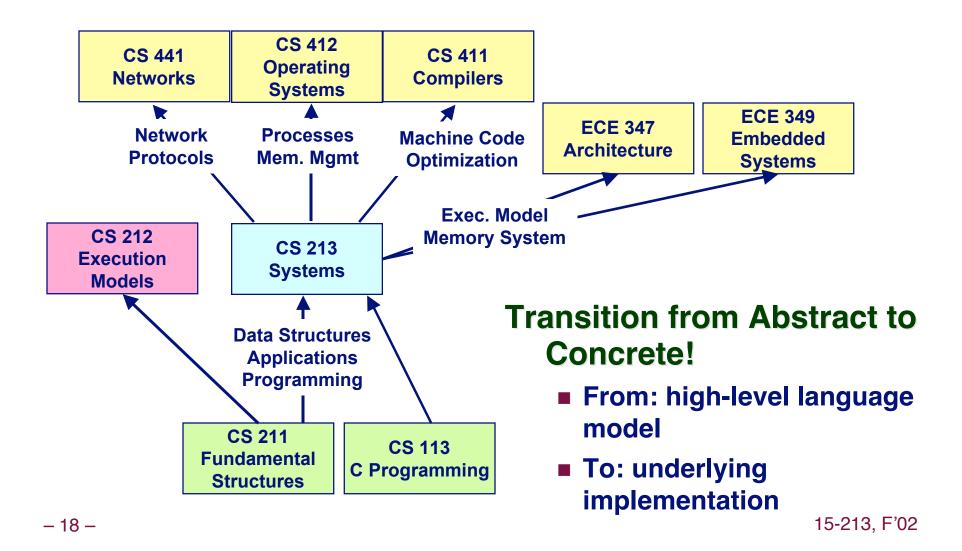
■ I/O system critical to program reliability and performance

#### They communicate with each other over networks

- Many system-level issues arise in presence of network
  - Concurrent operations by autonomous processes
  - Coping with unreliable media
  - Cross platform compatibility
  - Complex performance issues

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## Role within Curriculum



# **Course Perspective**

#### **Most Systems Courses are Builder-Centric**

- **Computer Architecture** 
  - Design pipelined processor in Verilog
- Operating Systems
  - Implement large portions of operating system
- Compilers
  - Write compiler for simple language
- Networking
  - Implement and simulate network protocols

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# **Course Perspective (Cont.)**

#### **Our Course is Programmer-Centric**

- Purpose is to show how by knowing more about the underlying system, one can be more effective as a programmer
- Enable you to
  - Write programs that are more reliable and efficient
  - Incorporate features that require hooks into OS
    - » E.g., concurrency, signal handlers
- Not just a course for dedicated hackers
  - We bring out the hidden hacker in everyone
- Cover material in this course that you won't see elsewhere

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