

**Note**

- This homework is more like a take-home exam.  
Your solution should meet the requirements of exam 1 solution.  
It will be graded in the same way.  
It will have much higher weight than other homeworks.
- You cannot ask questions about how to proceed, whether you are on the right track, etc.  
You can “translate” such questions to exam 1 solution, and we will answer those.
- You can do this homework individually or with one partner. It’s entirely up to you.  
Can’t find a partner? Not happy with your partner? Want to leave your partner? Not my concern.
- Your solution should be neat and readable.

**Problem 1 [30 points]**

This program below is a variation of the Otway-Reese protocol. It has an attacker, kdc Z, client A and server B. The attacker can read-write the channel and get A's old password (only when A is between sessions).

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Protocol(Z, A, B) { // kdc, client, server
  chan ← [];
  hst ← [];          // connect history
  kAZ ← random();   // initial A-Z key
  kBZ ← random();   // initial B-Z key
  startSystem(Attacker());
  startSystem(Kdc(Z,A,B,kAZ,kBZ));
  startSystem(Client(A,Z,B,kAZ));
  startSystem(Server(B,Z,A,kBZ));
}

Attacker() {
  α; // initially has A, B, Z, all programs
  // functions executable by attacker
  function rChan {α ← chan;} // read chan
  function wChan(x) {chan ← x;} // write chan
  function getPwdA() { // get A.key iff A.t at 1
    if (A.t at 1) {
      α.append(A.key);
      A.key ← Z.keyA ← random();
    }
  }
}

Kdc(Z, A, B, kAZ, kBZ) { // atomicity points: 1
  keyA ← kAZ;
  keyB ← kBZ;
  t ← startThread(client());
  return;

  function kdc() {
    while (true) {
      1: msg ← rx([B,Z,.]);
         x ← dec(keyB, msg[2]);
         if (x.size = 4 and x[0,1] = [A,B]) {
           nB ← x[2];
           y ← dec(keyA, x[3]);
           if (y.size = 3 and y[0,1] = [A,B]) {
             nA ← y[2];
             kAB ← random();
             rA ← enc(keyA, [nA,kAB]);
             rB ← enc(keyB, [nB,kAB]);
             tx([Z,B,rA,rB]);
           }
         }
    }
  }
}

```

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Client(A, Z, B, kAZ) { // atomicity points: 1, 2
  key ← kAZ;
  t ← startThread(client());
  return;

  function client() {
    while (true) {
      1: nL ← random();
         tx([A,B,1, enc(key, [A,B,nL])]);
      2: msg ← rx([B,A,.]);
         x ← dec(key, msg[2]);
         if (x.size = 2 and x[0] = nL) {
           kAB ← x[1];
           hst.append([A,kAB]);
           tx([A,B,2, enc(kAB, 'HELLO')]);
         }
    }
  }
}

Server(B, Z, A, kBZ) { // atomicity points: 1,2,3
  key ← kBZ;
  t ← startThread(server());
  return;

  function server() {
    while (true) {
      1: msg ← rx([A,B,1,.]);
         nL ← random();
         tx([B,Z, enc(key, [A,B,nL,msg[3]])]);
      2: msg ← rx([Z,B,..]);
         x ← dec(key, msg[3]);
         if (x.size = 2 and x[0] = nL) {
           kAB ← x[1];
           hst.append([B,1,kAB]);
           tx([B,A,msg[2]]);
        }
      3: msg ← rx([A,B,2,.]);
         if (dec(kAB, msg[3]) = 'HELLO')
           hst.append([B,2,kAB]);
    }
  }
}

```

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**Problem 1 (cont)****Part a.**

Does  $Inv A_1$  hold, where

$$A_1 : ((j \text{ in } \text{hst.keys}) \text{ and } j > 0 \text{ and } \text{hst}[j] = [A,p]) \Rightarrow \text{hst}[j-1] = [B,1,p]$$

If yes, assume that A appends  $[A,p]$  to hst at time  $t_0$  and prove that  $[B,1,p]$  is the last entry in hst just before  $t_0$ .

If no, come up with a counter-example evolution, i.e., ending in a state where  $A_1$  does not hold.

**Part b.**

Does  $Inv A_2$  hold, where

$$A_2 : ((j \text{ in } \text{hst.keys}) \text{ and } j > 0 \text{ and } \text{hst}[j] = [B,2,p]) \Rightarrow \text{hst}[j-1] = [A,p]$$

If yes, assume that B appends  $[B,2,p]$  to hst at time  $t_0$  and prove that  $[A,p]$  is the last entry in hst just before  $t_0$ .

If no, come up with a counter-example evolution, i.e., ending in a state where  $A_1$  does not hold.

**(Hint:**  $Inv \psi(A.\text{key})$  may hold.  $Inv \psi(B.\text{key})$  may not hold.)

**Problem 2 [30 points]**

Repeat problem 1 after changing the kdc-to-server message to include the response to A inside the response to B.

The change can be made as follows:

- In function kdc

<pre> ... rA ← enc(keyA, [nA, kAB]); rB ← enc(keyB, [nB, kAB]); tx([Z, B, rA, rB]); ... </pre>	becomes	<pre> ... rA ← enc(keyA, [nA, kAB]); rB ← enc(keyB, [nB, kAB, rA]); tx([Z, B, rB]); ... </pre>
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- In function server

<pre> ... 2: msg ← rx([Z, B, ...]); x ← dec(key, msg[3]); if (x.size = 2 and x[0] = nL) { ... tx([B, A, msg[2]]); ... </pre>	becomes	<pre> ... 2: msg ← rx([Z, B, ...]); x ← dec(key, msg[2]); if (x.size = 3 and x[0] = nL) { ... tx([B, A, x[2]]); ... </pre>
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