## Conventions

- $1(P)$, for a predicate $P$, denotes the indicator function of $P$, i.e., 1 if $P$ holds and 0 otherwise.
- $X_{i, j}$ denotes $\left(X_{i}\right.$ and $\left.X_{j}\right)$
- $X_{i . . j}$ denotes $\left(X_{i}\right.$ and $X_{i+1}$ and $\cdots$ and $\left.X_{j}\right)$
- $X$ satisfies invariance rule means
- initialization establishes $X$
- every atomic rule unconditionally preserves $X$
- $X$ satisfies invariance rule given $Y$ means
- Inv $Y$ holds
- initialization establishes $Y$
- every atomic rule unconditionally establishes $Y \Rightarrow X$ starting from $Y$ and $X$


## Part a

Define the following predicates:

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\(B_{0}: \alpha 01+\alpha 10+1\) (done) \(=1\)
\(B_{1}\) : union(s0.bg, s1.bg, \(\alpha 01, \alpha 10\) ) \(=\) union( \(\mathrm{B} 0, \mathrm{~B} 1\) )
\(B_{2}\) : s1.bg.size \(=\mathrm{B} 1\). size
\(B_{3}\) : s0.bg.size \(+\alpha 01+\alpha 10=\) B0.size
\(B_{4}: \alpha 10 \neq[] \Rightarrow \alpha 01\).head \(\leq \min (s 1 . \mathrm{bg}) \quad / /\) is it ok to write \(B_{4}\) as: \(\alpha 10\). head \(=\mathrm{q} \Rightarrow \mathrm{q} \leq \min (\mathrm{s} 1 . \mathrm{bg})\)
\(B_{5}\) : (done \(\left.\Rightarrow \max (\mathrm{s} 0 . \mathrm{bg})\right) \leq \min (\mathrm{s} 1 . \mathrm{bg})\)
```

$B_{0 . .5}$ satisfies the invariance rule and implies $A_{0}$

## More detail on how $B_{0 . .5}$ satisfies invariance rule

Each of $B_{0}, B_{1}, B_{2}$ and $B_{3}$ (individually) satisfies the invariance rule.
$B_{4}$ satisfies the invariance rule given Inv $B_{0}$. (Where is $B_{0}$ used?)
$B_{5}$ satisfies the invariance rule given $\operatorname{Inv} B_{4}$.

More detail on how $B_{0 . .5}$ implies $A_{0}$
$B_{0 . .3}$ and done imply s0.bg.size $=\mathrm{B} 0$. size, s1.bg.size $=\mathrm{B} 1$.size, and $\alpha 01$ and $\alpha 10$ are empty.
This and $B_{5}$ imply $A_{0}$.

## Part b

Consider the following function:
$G: \operatorname{sum}(u n i o n(s 0 . b g, \alpha 01, \alpha 10)$ )
$/ /$ sum of integers in $\mathrm{s} 0 . \mathrm{bg}, \alpha 01$ and $\alpha 10$
The following hold:

| $L_{1}:(G=k$ and $\alpha 01 \neq[])$ | leads-to $(\alpha 10 \neq[]$ and $(G<k$ or $\alpha 10$.head $\geq \max (\mathrm{s} 0 . \mathrm{bg})))$ | // via s1.receive |
| :--- | :--- | :--- |
| $L_{2}:(G=k$ and $\alpha 10 \neq[])$ | leads-to $((G=k$ and $\alpha 01 \neq[])$ or done) | // via s0.receive |
| $L_{3}:(\alpha 10 \neq[]$ and $\alpha 10$. head $\geq \max (\mathrm{s} 0 . \mathrm{g}))$ | leads-to done | // via s0.receive |
| $C_{0}: \operatorname{Inv} G \geq 0$ |  |  |

So $G$ is almost there; just need to augment it to handle $L_{2}$ and $L_{3}$ For brevity, define

$$
H: \alpha 10 \neq[] \text { and } \alpha 10 . \text { head } \geq \max (\mathrm{sog} \mathrm{~g})
$$

The following function works

$$
F:[G-1(H), \alpha 10 . \text { size }]
$$

