1 General description

The primary motivation of this project is of course to get you to build and make use of data structures, and to take a look at benefits and shortcomings for each one.

The TA’s motivation though is to be able to make pretty pictures with some rudimentary drawing primitives.

In project one we’ll have you build a dictionary to keep track of named dots (of various sizes and colors) using java’s api. You’ll also get a warmup implementing an Adjacency List along with Dijkstra’s shortest path algorithm.

In future parts of the project we will replace some data structures to examine performance tradeoffs. In particular you can probably expect to see the prof’s old time favorite the PM3 quadtree, and the TA’s favorite the B+ tree. We’ll probably start off project 2 with a skiplist and a kdtree for working with spacial data.

Expect a plot to evolve by project 2, possibly about a pinball planet. Or not. :)

1.1 Part 1: Crash Course in Java, Graphs, Shortest Path

This is the introductory project meant to acquaint you with Java, the language required for all projects this semester. You will hopefully learn the meaning of packages, interfaces, Iterators, the uses of Comparator and Comparable, and how to work with the very robust java API. You’ll also see that this will be a fast paced class, with this project due only 1 week after your first day of class. Future projects will usually have between 2 and 3 weeks to allow you scheduling flexibility, but will also be much larger and more technically challenging. It is expected that most of you are familiar with Graphs and Dijkstra’s shortest path algorithm, even if you don’t quite remember them perfectly.

1.1.1 CMSC420: Introduction to Command Parsing, where nothing could possibly go wrong.

This section is left in for completeness from previous semesters in case you want to ’roll your own’ Command Parser. However, this semester you will be provided with the TA’s ’adequate’ command parser. It’s not optimal(either in speed or coding style), but it is more than enough for what we will be doing.

Begin original commentary:

You are all blessed with a professor with a Ph.D. in fault tolerance. Because of this you will be expected to develop fault tolerant programs. This means bounds checking for input numbers and checking a number of possible error conditions for each command. It also means that your parser should never fail or crash because of a malformed command.

* Participation in this project may prove HAZARDOUS to your health. Unfortunately, failure to participate early and often will definitely have an adverse effect upon your GPA. Take my advice. Start now, because you’re already behind. If you don’t believe me, ask someone who took this course last semester.
A really useful command interpreter would give useful error messages about commands, such as “wrong number of arguments”, or “invalid argument type”, or even better “second argument was int, expected string”. For our purposes it will be sufficient to print a single error message regardless of the error:

*****
Error: Invalid Command.

Note the standard asterisks are printed. Do not echo the erroneous command.

Your parser must completely ignore blank lines. For all other lines which are not fully formed and correct commands you must print the above error. So if you can’t get a fully error checking parser you will not be hurt on the other parts of the project. You may assume that for commands other than error checking all commands will be upper case and there will be no spaces within a command. There will be no blank lines and all commands will be valid.

A working parser cannot make any of those assumptions. Commands should be correctly interpreted regardless of case (i.e. CREATE_DOT and create_DOT should both be interpreted as the CREATE_DOT command). Blank lines should be completely ignored (in particular do not print extraneous “*****”s). Whitespace in general should be ignored when parsing with one exception- any string (command name, dot name, or string argument like the colors “RED” “BLUE”, etc) cannot contain internal spaces. So a command like create dot (foo bar, 5, 6, BLUE) should be flagged as invalid.

1.1.2 Comments on java

This semester all projects are to be written in java. The version on the detective cluster is 1.4 and can be downloaded at:

http://java.sun.com/j2se/1.4/download.html

The online version of the documentation is at

http://java.sun.com/j2se/1.4/docs/api/index.html

I highly recommend you download the java sdk and do most of your work from home, if nothing else this will lighten the load on the now overworked dc machines ;) They can be very slow around lower level project due dates ;) Of course as in most cs classes your projects will be compiled and run on the detective cluster machines (dc.umd.edu), so you should check to be sure your projects work there. However, there should not be any portability issues as long as you develop with the correct java version. We’ll eventually be working with graphics, and if you work over ssh from home you will have to figure out how to set up an X-server to do this on your own machine, something I haven’t figured out in the last 2 years :) (largely due to just using java locally on my own machine).

I’ve heard good things about Borland’s free jbuilder (registration required); You may wish to look into that.

While you are permitted to use any java drawing facility you are comfortable with, a simple drawing package is available on the class web page. It is this package that will be most readily supported by the TA’s should any problems arise. The package ’Canvas.java’ provides a simple class which allows drawing of circles, squares, lines, captions, and other simple primitives in a java jframe. While this isn’t being used in project one, it will show in the not to distant future.

1.1.3 Sets vs. Maps

For your work in this class Maps will generally be preferred over Sets. Maps are like sets, except that each key entry is associated with some value which can be anything. So, for instance, if you are using a class with a name like our ”Dot” class, you can use the String name as a key and the whole Dot as the value. Maps are good because you need only know the key to get back the value object you want. (i.e., read in a string name from the user, and use it to look up the whole corresponding dot).
In a Set in java there is no easy way to pull out an object with a specific key other than to iterate over all the elements of the set. There is an add(Object) method which lets you add to the set, but there is no corresponding get(Object x) method. The assumption is that if you already have x, why do you need the set to return you another copy? They’re the same thing! Sets in java instead have a contains(Object) method, which won’t be of much use as a dictionary.

1.1.4 More on Java- pass by reference, but not really

Every semester a new group of students get caught up by the same thing in Java. They start out hearing ”Java is always pass by reference” and they do silly looking things like the following:

```java
void foo(String t)
{
    t = new String("World");
}

String s = new String("Hello");
foo(s);

System.out.print(s); //prints "Hello". Why didn’t it change?
```

In true pass by reference C++ this would have worked. But what is happening is not really pass by reference, it is pass by value, except what is being passed is a pointer. If you were to transfer the above to C++ it would look like:

```java
void foo(String *t)
{
    t = new String("World");
}

String *s = new String("Hello");
foo(s);

cout<<*s<<endl; //prints "Hello". Hopefully obvious why
```

You can see in the second example that t is only a local copy of s. If you alter the value t is pointing at then s will see the change. However, if you point t at something else s will never know. In this example there is actually no way for foo to change s, since java Strings are immutable after creation. An error less obvious than the above is:

```java
void foo(String t)
{
    t = t+"World";
}
```

This looks like concatenation, not reallocation, but that ‘+’ operator actually allocates a new String. The above is actually just a shortcut in java for:

```java
void foo(String t)
{
    StringBuffer temp = new StringBuffer();
    temp.append(t);
    temp.append("World");
    t = temp.toString();
}
```
It's important to realize what's going on in the background! Of course in the above example, foo still doesn't change t, but what you could do instead is:

```java
void foo(StringBuffer t)
{
    t.append("World");
}
```

This time, since t always points to the same place the original value really is modified. In java "pass by reference" as C++ programmers tend to think of it always requires some kind of wrapper. In the last example, StringBuffer is a wrapper for a dynamically sized character array. There is a quick and dirty hack to get a similar effect without building and entire class wrapper, pass a 1 element array instead:

```java
void foo(String[] t)
{
    t[0] = new String("World");
}
```

String s[]=new String[1];
s[0] = new String("Hello");
foo(s); //s[0]= "World"

This works for a similar reason. t points to the same array in memory that s does. When an element of the array is updated by t, s will see the change as well. This ends my FYI on pass by reference, try not to get caught up by this common error :)

### 1.1.5 Java last part, comparators

Java has two basic tools for doing comparison- the compareTo() method of Comparable objects, and the compare() method of the Comparator class. I will try to explain this with a quick example. Suppose I wanted to sort a collection of strings in alphabetical (ignoring case for the moment) order. I might use a TreeSet to do this:

```java
SortedSet sorter = new TreeSet();
sorter.add("hello");
sorter.add("world");
sorter.add("cat");
sorter.add("dog");

Iterator i = sorter.iterator();
while(i.hasNext())
    System.out.print(i.next()+ " ");//prints "cat dog hello world"
```

In the above example the string constants are automatically cast to String, which implements the compareTo() method (just like c/c++ strcmp). The sorted map assumes its elements are Comparable and uses compareTo() to sort them. Note that this is a unique case where next() can be used without a cast, since in java ALL objects implement the toString() method which is automatically called here. Back on topic, what if I wanted to sort the words backwards? One way is to wrap the strings in another class that implements compareTo() backwards like:

```java
class MyString
{
    String s;
    MyString(String s){this.s=s;}
    public int compareTo(Object other){return -1*(s.compareTo(other));}
}
```
This would work, but it is a bit of a mess. We can’t extend String directly, since Sun has made it a final
class. In any case it will be non obvious what it is you are doing. However, there is a better way that you
can use with all of Java’s sorted classes (and which you will implement in your own sorted maps later this
semester). You can use a comparator:

```java
class ReverseCompare implements Comparator
{
    public int compare(Object a, Object b)
    {
        return (-1*(((Comparable)a).compareTo(b)));
    }
}
```

```java
SortedSet sorter = new TreeSet(new ReverseCompare());
sorter.add("hello");
sorter.add("world");
sorter.add("cat");
sorter.add("dog");
```

```java
Iterator i = sorter.iterator();
while(i.hasNext())
    System.out.print(i.next()+" "); //prints "world hello god cat 
```

Because the TreeSet was given a Comparator in its constructor it will no longer assume its elements are
Comparable, and will use the Comparator for sorting instead. Comparators allow you to easily have sets
with different types of objects which aren’t natively comparable with each other, or to impose your own
sorting rules on other people’s classes (like String) with ease. Cool stuff. You’ll hopefully find this useful
for doing your coordinate checking in project 1, as well as implementing a priority queue for your adjacency
lists.

1.1.6 Adjacency List

The only data structure actually have to implement for project 1 is and adjacency list. You’ll recall that an
adjacency is conceptually a ‘list of lists’, ie:

```
A->C->D
C
D
F->K->J
H->K
J
K
```

You should implement this with TreeSets, TreeMaps, Comparators, and any other java classes you find
handy. (You of course may not use an existing graph class if one already exists. You are also admonished
to not use HashMaps/Sets over TreeMap/Sets.)

To implement dijkstra’s algorithm you will also need a priority queue which the Java api sadly does not
provide. This can be implemented through clever use of a TreeSet/Map and Comparators (the problem with
a plain set is that you can’t have two elements with the same priority, so you need a way to always break
ties).

Note that insertion/deletion from this structure is \(O(\log(n)\log(m))\), where \(n\) is the number of nodes in
the graph and \(m\) is the degree of the graph (the max number of edges incident to a single vertex). This
represents a binary search to find the correct row of the list to find the starting vertex, followed by a a binary
search to check for existance of the ending vertex.
You'll use the graph to implement shortest path using dijkstra's algorithm. If you have a fibonacci heap handy the run time of this algorithm is \( O(V^*\lg(V)+E) \) (where \( E \) is the number of edges and \( V \) is the number of vertices). And a fibonacci heap is what you ask? Well, I have no idea, but those of you who go on to take algorithms classes will no doubt hear about fibonacci heaps again, and I just wanted to be the first to share. Let's just say that they are magical, and that you will probably *never* learn how they work. If you've printed this spec on paper, feel free to X this paragraph to avoid reading it ever again. End digression.

So anyway, you will be required to do shortest path in \( O(E\lg V) \). It would be nice if you understand where this bound comes from, so I will explain it below, but it suffices that if you implement the algorithm correctly, that is the runtime you will have.

Allow me to sketch the algorithm to explain the running time (I'm not trying to teach the algorithm here- see your book/class/newsgroup/google). Every iteration of this algorithm you are guaranteed to find the correct shortest path to exactly one node- so we know right away there will be \( V \) iterations. At each state your graph is split in two sections- the 'solved' section, for which you know the correct distances, and the rest, which have some distance values associated with them which may or may not be accurate- this set is stored in some kind of priority queue. An iteration begins by selecting the node (call it 'N') from this queue with the best distance value, adding it to the 'solved' set, and 'relaxing' all it's edges. Relaxing is when for each node adjacent to \( N \) we see if it is faster to get to that node through \( N \) than it's current best known path. We update distances and back-pointers appropriately (so we know what the shortest path actually is when we finish), and that ends the round. Note that if a node's distance value is changed, its position in the priority queue has to be fixed up somehow. (this is where the magical fibonacci heap would come into play, it's got an advantage in this 'fix up' step). One way to do this is just to have multiple copies of the same node in the queue and ignore them when they come back up (this is the approach I will use in the explanation), or else to remove the old value before reinserting it. Either works.

Now, how long does all this take. There are \( V \) rounds, and in every round we have to pull something out of the front of the priority queue- which with a TreeMap is an \( O(\lg E) \) operation, so each round takes \( O(V\lg E) \) time. Rather than try and deal with how many elements are added to and removed from the queue in any single round, it is easier to think about how many such operations can occur in the life of the algorithm. Every single edge in the graph has exactly one opportunity to add a vertex to the queue (during a relax operation), so there are \( O(E) \) possible insertions. If we allow duplicates, the size of the queue can grow to \( O(E) \), so we may treat each queue operation as \( O(\log(E)) \). That gives \( O(E\log E) \) running time for all queue operations during the life of the algorithm. Together that gives a running time of \( O(V\lg E+E\lg E) \), which since \( E \) is bounded by \( V^2 \), is \( O(E\lg V) \). And that's the required running time of your search;

*Please be careful with your implementation details.* In particular, do not use a linear time priority queue. This nailed a lot of people last semester. I will time your project on *very* large inputs. Check out:
http://www.cs.umd.edu/class/spring2003/cmse420/p2graphs/

If you scroll down you'll runtimes for the large shortest path test from last semester. A negative time indicates that the test was failed. The bottom third of students timed out. They may have found the right answers eventually, but all the same they got no credit for that test. The middle third did not find the correct paths. Only one third of the class got the correct output in the 60 second time limit, and you can see that most of them were much slower than the best solutions. Lots of people like to blame java for the slowness of their own code, but before you start to pick on java you need to make sure you implement your code intelligently :) That is all, you've been warned.

1.1.7 Part 1 Command Specification

No longer relevant.

1.2 Part 2: SkipList and Drawing

No longer relevant.
1.3 Part 3E: PM3/PM1 Quadtree, B+ Tree, and More Drawing

1.3.1 Part 3E: Command Specification

The following is a list of commands you should support for part 3E and a description of the output you should give for each one. Note that for all functions, you should print '****
' followed by a ' ' and an echo of the command given. For instance, the entire valid output to EXIT() is

****
=> EXIT()
Have a nice day!

The sample output should make this clear. This is done to negate the effects of input redirection and to assist in grading. Note that although it is done in the samples that will appear later, you are not required to reformat the original command (fixing spacing, for instance) in any way.

The definitions below will use the following standard BNF definitions.

<dotlist>:=<dot><nl><dotlist>|<dot><nl>
<dot>:=<name> at (<int>,<int>) color:<color>
<color>:= RED|GREEN|BLUE|BLACK|WHITE
<DNE>:=Error: The specified dot does not exist.<nl>

<framelist>:<frame>|<frame><framelist>
<frame>:Frame <int><nl><seglist>
<seglist>:<segment>|<segment><seglist>
<segment>:= (<name1>,<name2>)

The <int> in a frame is the 0 based frame number for the current drawing command. For instance, if you are drawing the 300th frame of ANIMATE_HORIZONTAL_PATH, then the int should be 299. The segment list should be printed as described in the PRINT_QUADTREE command. That is, the endpoints of each segment should be sorted in increasing coordinate order (even though it is the names that are printed) and segments in the list should be sorted based on the first endpoint (ties broken by second endpoint). This will require you to do some sorting for every frame displayed in text- this overhead would not be needed for the graphic version since order of drawing is not really relevant.

Coordinate order will be defined as in p2, with: P1 < P2 if P1.x < P2.x or (P1.x==P2.x and P1.y>P2.y) P1 > P2 if P1.x > P2.x or (P1.x==P2.x and P1.y<P2.y) P1 == P2 otherwise

Whenever a <double> appears (not in this part, but it will come up later), it means a floating point decimal number printed with exactly three digits after the decimal place (including trailing zeros as necessary).

Also, when looking at the list of errors, eg.

<error>:= <DNE>|<NR>|<AI>|<DC>|<NZ>

the leftmost applicable error should always be the one printed.

SET_BPTREE_ORDER(btree_order) will indicate the order of the B+ tree used in the data set. It will always be the first command. Don’t bother checking to see if some OTHER command is the first. We won’t do that. However, you should detect that the B+ tree order has already been set if the command appears again.

The order will never be less than 3. You should check for this condition to avoid crashing horribly on a bad input, but I won’t add a specific error for this part.

Note that the following rules apply to B+ tree nodes:

Internal: must always contain between floor((btree_order−1)/2) and btree_order−1 keys, with exactly one more child than the number of keys at all times. (this implies between ceiling(btree_order/2) and btree_order children per node, inclusive).
Leaf: must always contain between \(\text{ceiling}(\text{btree\_order} - 1)/2\) and \(\text{btree\_order} - 1\) keys, inclusive. Must not contain \(\text{btree\_order}\) keys! This is to force consistency between our projects so that I have some chance of grading. Remember the root is an exception, in that it never has a lower bound on the number of keys it contains.

Also, whenever a value is equal to a key it must go to that key’s RIGHT child. This is mandatory, meaning no credit will be given for the B+ tree if this rule is not observed.

Even if you do not implement the b+ tree you must still implement this function! Default to printing the correct <success> message. Because this is always the first command and diff is used in grading, if you skip this function your project will fail every test!

Output summary:
<output>:=<success> | <error>

<success>:= Order set to <btree\_order>.<nl>
<error>:= Error: B+ tree already initialized.<nl>

EXIT() ends the program. Your program should also naturally terminate (with no exit message) when end-of-file is reached. Print a goodbye message (as specified below) and exit the command interpreter.

Output summary:
<success>:= Have a nice day!

CREATE\_DOT(name, x, y, radius, color) creates a 'dot' object with the appropriate name, coordinates, radius, and color. The dot will then be added to a B\+Tree sorted based on the name (the data dictionary) and to one based on the coordinates. The latter structure is used to check for duplicate coordinates.

Coordinates will be non-negative. This should be an O(logn) operation where n is the number of dots already in the dictionary. The dots should be stored in an ascibetically sorted B\+tree. If a dot with the same name already exists print an error. If a dot already exists at the specified coordinates print an error.

Output summary:
<output>:=<success> | <error>
<success>:= Created dot <dot>.<nl>
<error>:= <AE> | <DC>
<AE>:= Error: Dot <name> already exists.<nl>
<DC>:= Error: Dot <other\_dot\_name> already exists at the specified coordinates.

DELETE\_DOT(name) Removes the dot with the given name from the data dictionary (and the coordinate checking dictionary). If the dot does not exist print an error. Delete should be O(logn).

Output summary:
<output>:=<success> | <error>
<success>:= Deleted dot <name>.<nl>
<error>:= <DNE>

RANGE\_DOTS(name1, name2) Lists all dots with names between name1 and name2. If name1 \& name2 the dots must be listed in increasing strcmp order (endpoints included, neither name1 nor name2 need actually be in the dictionary). If name1 \& name2 The dots must be listed in reverse strcmp (String.compareTo()) order. Show off that B\+ range search ;)
Output summary:
<output>::<success>|<error>

<success>::<dotlist>
<error>:: No matching dots found.<nl>

PRINT_BPTREE() requires you to list the B+ in a breadth first search order. If you used links between internal nodes this will be easier, BFS is more complicated. Every level of the tree is enclosed in braces {}, every node is enclosed in parenthesis, every key within a node is separated by commas. Each level of the tree should appear on its own line and in order. A sample tree of order 3 is printed below.

{(bar)}
{(DOT3),(foo)}
{(DOT1,DOT2),(DOT3),(bar),(foo)}

Note the leaf DOT3 is to the RIGHT of the key DOT3.
Even at the leaves print only the key (the dotname). If the tree is empty, print "Tree is empty." Your tree is not expected to match mine exactly. Your grade will be based on your tree displaying the properties described above in the SET_BPTREE.ORDER command.
See SET_BPTREE.ORDER for rules on the number of keys in a node. Remember that there must be one fewer guides than children (no ‘extra’ key on the far left should be printed, even if you used one in your implementation). Your tree, of course, must also contain the correct data at the leaves!

Output summary:
<output>::<success>|<error>

<success>::<b+rows><nl>
<b+rows>::<b+row><nl><b+row>::<nodes>
<nodes>::<node>,<nodes>|<node>
<node>::<keys>
<keys>::<key>,<keys>|<key>
<key>::<dotname>

<error>:: Tree is empty.<nl>

INIT_QUADTREE(size) Sets the size for the quadtree. After initialization the quadtree should hold points/lines that fall between [0,0] at the southwest/lower left hand corner, and \([2^{size},2^{size}]\) in the northeast upper right hand corner. If size is less than 2 or more than 30 print an error. For simplicity your quadtrees should have a minimum partition size of 1x1. In particular, this means that your trees cannot hold adjacent coordinates.
This command will always precede the first command which requires the quadtree. This command is only valid the first time it is successfully called (out of range is not a successful call). If an attempt is made to reinitialize the tree, print an error.

Output summary:
<output>::<success>|<error>
<success>::Quadtree initialized.<nl>

<error>::<00R>|<INIT>
<00R>::Error: size out of range.<nl>
<INIT>::Error: The Quadtree has already been initialized.<nl>
MAP SEGMENT(name1,name2) Adds a segment to the PM3. It is possible that name1==name2, this should not be an error. The PM3 should be able to know whether any given point has a path to itself or not. There is one further error to detect. If the segment intersects a segment already in the tree (except for shared endpoints), print an error and leave the PM3 unchanged. If a subdivision smaller than 1x1 would be forced then this should also be treated as an intersection. You may use this error even if the segment intersected is the same as the one you are trying to insert (ie. trying to remap a segment that was already in the tree).

Output summary:
<output>::=<success>|<error>
<success>::=Mapped segment (name1,name2).<nl>
<error>::=<DNE>|<ID>

<ID>::=Error: Intersection detected.

UNMAP SEGMENT(name1,name2) Deletes a segment from the PM3. If the points do not exist, or the segment does not exist, print an error. If after the operation an endpoint would be left with no adjacent segments then it should also be completely removed from the tree. The tree should be collapsed so that it is minimal- which basically means that the tree should look as if the segment had never been in the tree to begin with. (Some collapsing will need to be done). Remember that if (A,A) is explicitly added to the tree via MAP SEGMENT, then (A,A) must be unmapped before A is removed. Likewise A should never be an isolated point in the tree unless (A,A) was explicitly added.

Output summary:
<output>::=<success>|<error>
<success>::=Unmapped Segment (name1,name2).<nl>
<error>::=<DNE>|<SNF>
<SNF>::=Error: The specified segment was not found on the map.

NEAREST SEG TO POINT(x,y) Finds the Segment in the PM3 closest to the specified point, along with the distance to the segment. If multiple segments are the same distance from the point pick your favorite. If the quadtree is empty print an error. Print the endpoints of the segment in coordinate order.

Output summary:
<output>::=<success>|<error>

<success>::= Nearest segment: <segment>. Distance: <double>.<nl>
<segment>::= (name1,name2)
<error>::= Tree is empty.<nl>

PRINT QUADTREE() Prints out the PM3 Quadtree map. When printing you must print in the following order: Northwest, Northeast, Southwest, Southeast. If you reach a leaf with a dot in it you should print the dot, followed by all the segments in the quadrant in coordinate order (more detail below). If a black leaf does not contain an endpoint dot, just print the segments.

Output summary:
<output>::=<success>|<error>

<success>::= <pmtree><nl>
<pmtree>::=<black_node><nl>|<white_node><nl>|<nl><grey_node>
<grey_node>::= NW <pmtree> NE <pmtree> SW <pmtree> SE <pmtree>
<black_node>:=<dot>:<seglist>|<seglist>
<white_node>:=
<error>:= Tree is empty.<nl>

Clarifying the description above, for a black node you should first print the endpoint in that segment if there is one, followed by a list of all the segments in the quadrant. For each segment the dots corresponding to name1 and name2 should be sorted in the coordinate order used in p2. The edges themselves in the seglist should be sorted in coordinate order based on the first endpoint (ties broken by the second endpoint), even though it is the String name that is printed.

EXIT() ends the program. Your program should also naturally terminate (with no exit message) when end-of-file is reached. Print a goodbye message (as specified below) and exit the command interpreter.

Output summary:
<success>:=Have a nice day!

COLOR_SEGMENTS((lx,ly,ux,uy,color)) This function will first locate all segments in PM3 which OVERLAP (inclusive) the specified rectangular region (determined by (lx,ly),(lx,uy),(ux,ly),(ux,uy)). Then it will change the color of the endpoints of those segments to the color specified. The segment endpoints will not necessarily be within the rectangle. Print the number of unique segments matched. Success can also be determined with the DRAW_MAP or LIST_DOTS commands.

Output summary:
<output>:=<success>
<success>:=Update complete. Found <int> segments.

SET_DRAW_MODE(mode, xsze, ysize, step) Changes the mode for all graphical output commands. The mode will be either "TEXT", "DRAW", or "BOTIF". If BOTH, then you should print text before drawing to the screen. "TEXT" will be the mode used when grading, the others will mostly be for your benefit in debugging and for impressing your friends. If the input is invalid (I will not test this) stay in the current mode. To be safe I will always call this function before using a drawing command, so you may use whatever default you find handiest.

The parameters xsze and ysize will specify the size of the view window used on the map. When drawing a frame the current coordinate should be treated as the center of the view window, with ysize units above and below the center and xsze units to the left and right so that the actual window is 2*xsze by 2*ysize. The step is only used for animation commands and says how far to move each frame in the current direction of travel (always east/right now). If you are following a particular line and the next step will cause you to pass the end of that line then you should have a frame exactly at that endpoint instead.

DRAW_FRAME(xpos,ypos) Processes the first frame only of the ANIMATE_HORIZONTAL_PATH command, centered at (xpos,ypos). Equivalent to ANIMATE_HORIZONTAL_PATH(xpos,xpos,ypos).

Output summary:
<output>:=<optional textoutput><nl><success>
<success>:=Draw Frame Complete.<nl>
<optional textoutput>:=<frame>

ANIMATE_HORIZONTAL_PATH(xmin,xmax,ypos) The first frame should be centered at (xmin,ypos), and then you should proceed right moving 'step' units each frame (as defined SET_DRAW_MODE) till reaching a frame centered at (xmax, ypos).

There is no such thing as failure, and there are no bounding requirements for xmin,xmax, and ypos (well, they will be non negative). If the current mode is BOTH or TEXT then each frame should also have a textual version printed according to the BNF.
For P3 the frame content will change to now be a list of segments rather than a list of points, which overlap the rectangle corresponding to each frame (boundary edges inclusive) and you will have to use the PM3 rather than the SkipList to select the points. (You may need a skiplist to do a secondary sort to print the segments in the correct order).

Essentially, for each frame you will have to select a rectangle of lines from the PM3, sort it appropriately, and print the lines. You will have to perform a similar rectangle search for the COLOR SEGMENTS command. For this reason I recommend you write a single function that can get rectangles of lines and return them in a set of some kind to be operated on :) Your search is expected to be as efficient as possible for a PM3.

Output summary:
<output>::<optional textoutput><nl><success>
<success>::=Animation Complete.<nl>
<optional textoutput>::=<framelist>

1.4 Pick Your Poison Commands

Your Part 3E score will come from the points for the commands listed in Section 1.3.1, and from whatever functions you choose to implement in this section. nical scores to bolster your project grades because there are some ridiculously easy points to be had in this section. Credit is given for taking risks and screwing up, not for avoiding the challenge all together.

1.5 Submission Instructions

To make your submission file, make a directory and copy all required files into it. Change to that directory and type:

```
tar -cvf part#.tar *
gzip part#.tar
```

To submit type(submit will usually be working at least a few days before the due date ;):

```
~/ms20001/Bin/submit # part#.tar.gz
```

In all cases ’#’ represents the number of the project part you are submitting(1,2,3 or 4). The filename is not really important; It is important that the file is in .tar.gz format and that your Main.java and other required files are not in a subfolder of the tar file. Subfolders are allowed, such as the cmisc420 package folder.

You must include the following with every submission: All necessary source files (*.java etc.) to compile your program. A file called README, all upper case, which contains your name, login id, and any information you would like to add.

 If you leave out the README your project will fail to submit!
 There may also be other per project required files that will be checked by the submit program.

 You are welcome to use a makefile for development (javac doesn’t track dependencies very well) but I should be able to run your project with the following two commands:

```
javac Main.java
java Main
```

No promises are made that I will read your READMEs, but they are useful when problems come up with a project.

There is a 100K filesize limit. Please don’t include .class files- I will probably strip them out before testing your projects anyway.

Every ontime submission will overwrite any previous ontime submissions, every 1day late submission will overwrite any previous 1day late submission and so on. So I will have one submission for every valid
submission period. (Late submissions will be dealt with by Dr Hugue on a case-by-case basis). I will grade every submission that is saved (including applicable bonuses and penalties) and you will get the highest grade among those that

If there are any errors in my IO you are still responsible for them- the spec is what is in charge. So if you submit all my current IO and submit early and then someone points out that I printed the wrong error message for some function, you have to fix your project and resubmit ;) You should be coding to match the command specification, not my sample IO.

Here is a (c++) makefile that you might use as a hint for how to set up dependencies for make. Note that that’s a TAB before $(CC), and make does care. (You’ll get an ‘invalid separator’ error, or something like that if you use spaces).

CC = cxx
FLAGS = -lm

proj4: bnnode.o btree.o cell.o main.o pmedge.o pmpoint.o pmquadtree.o util.o
$(CC) $(FLAGS) *.o -o proj4

bnnode.o: bnnode.cpp bnnode.h bdata.h
$(CC) -c $(FLAGS) bnnode.cpp

btree.o: btree.cpp btree.h bdata.h
$(CC) -c $(FLAGS) btree.cpp

cell.o: cell.cpp cell.h bdata.h pmpoint.h pmedge.h celledge.h
$(CC) -c $(FLAGS) cell.cpp

main.o: main.cpp btree.h cell.h celledge.h pmedge.h pmpoint.h util.h psdraw.h $
$(CC) -c $(FLAGS) main.cpp

pmedge.o: pmedge.cpp pmedge.h pmpoint.h util.h
$(CC) -c $(FLAGS) pmedge.cpp

pmpoint.o: pmpoint.cpp pmpoint.h pmedge.h
$(CC) -c $(FLAGS) pmpoint.cpp

pmquadtree.o: pmquadtree.cpp pmquadtree.h pmpoint.h pmedge.h util.h psdraw.h
$(CC) -c $(FLAGS) pmquadtree.cpp

util.o: util.h util.cpp
$(CC) -c $(FLAGS) util.cpp

clean:
rm -f *.o
rm -f proj4
rm -f core

1.6 Grading

There will be a few parts to grading your projects. Your projects will be graded running them on a number of test files for which I have already created correct (we hope) output. Your output will have all punctuation, blank lines, and non-newline whitespace stripped before differ similarly cleaned files.

Some of your data structures may be included with the TAs own testing code to test their efficiency and correctness.
Some text output cannot always be graded by simply differencing because there is no guarantee that we will have the same output. In these cases your project’s output will be pre-processed. In the case of the B+ tree, for instance, this program will verify that each node has the correct number of keys, that they are correctly ordered, and that all the correct data is at the leaves (and any other rules I may have left out).

Thanks to the miracle of automation you should expect your projects to be run on very very large inputs.

Typically each test file will be worth 10 points, and you will be eligible for either 10 or 0 points depending on whether you pass for fail that test. There is no partial credit for an individual test. I may give points projects that fail a test because of ‘small’ errors after initial grading at my own discretion. The tests will try to test mutually exclusive components of your projects independently. However, if you don’t have a dictionary which at least correctly stores all points so that some ‘get lost’, you may still end up failing other tests since they all require a working dictionary.

1.7 Standard Disclaimer: Right to Fail (twice for emphasis!)

As with most programming courses, the instructor reserves the right to fail any student who does not make a good faith effort to complete the project.

If you have problems with completing any given part of the project please talk to Dr. Hugue immediately – do not put it off! While the TA enjoys failing students, Dr. Hugue does not, so please be kind and do the project. A submission that gets only 20 or 30 points is considerably better for you than no submission at all.

1.8 Integrity Policy

From Dr. Hugue:

Your work is expected to be your own or to be labeled with its source, whether book or human or web page. Discussion of all parts of the project is permitted and encouraged, including diagrams and flow charts. However, pseudocode writing together is discouraged because it’s too close to writing the code together for anyone to be able to tell the difference.

Since the projects are interrelated, and double jeopardy is not my goal, we have a very liberal code use and reuse policy. First and foremost, use of code produced by anyone who is or has ever taken 420 from me requires email from provider and user to be sent to the instructor.

The instructor is the sole arbiter of code use and reuse, and reserves the right to fail any student who does not make a good faith effort on the project, or who refuses to adhere to the policies stated herein.

Remember, it is better to ask and feel silly, than not to ask and receive a complimentary F or XF.