Boycotting and Extorting Nodes in an Internetwork

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#### **Consumer Choice**

- What products to buy
- What companies to buy from
- What companies to not buy from
- Driving force in many markets
- In the Internet
  - First-hop ISPs
  - Destinations
  - At the fringe of the network
- This talk: Consumer choice in the core of the Internet











- Boycotting as a routing policy
  - Do not allow packets to go through the following nodes

• Assurances that the policy is maintained

# Why Boycott?

- Questionable business practices
  - AOL published users' searches
  - Google complied with China's censorship laws
  - Verisign redirected mistyped DNS queries to their advertising

#### Questionable network policies

- Snooping
- Traffic shaping
- Conflict of interest
  - Pending lawsuits
  - Competing products

# Can Incentives Help?

- Incentives generally used to ensure something does happen
  - Perform the task or else you do not get paid
  - Must prove it gets done
  - Coupled with accountability
- Can we provide incentive to ensure something does not happen?
  - Accountability alone is not enough
  - Offer economic disincentive from deviating
  - "If you can't prove you're good, I'll make it expensive to be bad."



- Choose routes that are "far" away from boycotted nodes
- Far enough away to remove any economic incentive

#### Boycotted by s

S

d



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Source node *s* pays for the whole path

Nodes can lie about c(i, j)

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#### A mechanism that gives

- A short-cost path with no boycotted nodes
- Economic incentive to truthfully report c(i, j)
- Economic disincentive from going through boycotted nodes instead





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#### VCG

I. Gather c'(i, j) from each i for each  $(i, j) \in E$ 

2. Compute the shortest cost path (SCP) from s to d

3. For each node i on the path, s pays i

$$\frac{\mathcal{C}(\operatorname{SCP}(s, d, G_{-i})) - \mathcal{C}(\operatorname{SCP}(s, d, G))}{||} + \frac{c'(i, j)}{||}$$
Marginal benefit of *i i*'s stated cost









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Insight: the shortest path between **t** and **u** went through a boycotted node

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Insight: the shortest path between **t** and **u** went through a boycotted node

I. Gather c'(i, j) from everyone

2. Compute all-pairs shortest paths

3. Remove all edges (*i*, *j*) where the shortest *path* from *i* to *j* goes through a boycotted node



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# Extorting Nodes

- Problem: B can threaten A
  - "I will understate unless you pay me 99% of your profit"
- Credible threat?
  - $\checkmark\,$  Both  $\boldsymbol{A}$  and  $\boldsymbol{B}$  profit
  - X B risks global understatement
    - Only if profit from extortion > loss from understating



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• Solution: Ignore everything **B** says

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  - If boycotting i, ignore it: c'(i, j) = 0

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# Quality of the Mechanism

- Truthfulness
- Relative cost of paths
- Ease of deployment

## Truthfulness

- Truthfulness is not always a dominant strategy
- A can overstate to ensure that
   A-C will not get trimmed
- Only if gain > loss in profit from overstating
- How do we resolve this?
  - Other ways to infer costs?



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#### Relative Costs of Paths

• "Cost of boycotting" set  $\mathscr{B}$  =

Total \$ paid when boycotting  $\mathscr{B}$ Total \$ paid without boycotting  $\mathscr{B}$ 



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#### Total Payment Made by the Source



#### Sometimes costs less to boycott

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#### $| + 5^*(24+1) = 126$

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# Can We Deploy This in the Internet?

- Not any time soon
- We currently lack infrastructure
  - Micropayments
  - Accountability
- Many remaining technical challenges
  - User policy vs. network policy
- We can
  - Deploy in ad hoc networks
  - Use overlays as an approximation to source routing

# Should We Deploy This?

- Consumer choice drives most markets
  - Why not Internet routing?
- Pro: Places the tussle in-band
  - Customer vs. ISP companies
- Con: Introduces new tussles
  - User policy vs. Network admin policies
  - Boycotting at the cost of performance?

## Summary

- Accountability alone is not enough; use incentives
  - "If you can't prove you're good, I'll make it expensive to be bad."
- Variations on VCG
  - Anti-boycott, anti-extortion
- Open (theoretical):
  - Overstating to disassociate from boycotted nodes
  - Boycotting to use paths with lower marginal cost
  - Infrastructure to support this: overlays?
- Can / Should we deploy this?

# Supplemental

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#### Path Costs



Always longer paths; generally not much longer

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As long as you allow virtualized links (always), accountability is not enough

#### Extreme Case: A Sybil Attack

XI

X2

**X**3

S

B

C

- Overstate costs to infinity to appear disconnected
- One Sybil per each pair of neighbors



- Difficult for ISPs to perform
  - Need to lie about connectivity
  - Need multiple companies