

## WORKSHEET ON PROPORTIONAL ALLOCATIONS SOLUTIONS

### PROBLEM ONE

	$I_1$	$I_2$	$I_3$	$I_4$	$I_5$	$I_6$
$A$						
$B$						

Fill in this table such with natural numbers such that:

- The numbers in a row add up to 12.
- If  $A$  knows  $B$ 's valuation then, in cut and choose, she can cut the cake so that she will do much better than 6 (which is half).

### PROBLEM TWO

	$I_1$	$I_2$	$I_3$	$I_4$	$I_5$	$I_6$
$A$						
$B$						
$C$						

Fill in this table such with natural numbers such that:

- The numbers in a row add up to 12.
- There is a way that if  $A$  cheats in the lone divider protocol, she will do worse than if she didn't.

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PROBLEM THREE

	$I_1$	$I_2$	$I_3$	$I_4$	$I_5$	$I_6$
$A$						
$B$						
$C$						

Fill in this table such with natural numbers such that:

- The numbers in a row add up to 12.
- There is a way that if  $A$  cheats in the Trim protocol, she will do worse than if she didn't.

PROBLEM FOUR

Generalize the Lone Chooser Algorithm to 4 players

PROBLEM FIVE

Do the last case of the Divide and Conquer Algorithm that I omitted.

PROBLEM SIX Recall that for the DC algorithm we have:

Let  $C(n)$  be the number of cuts used for DC with  $n$  players.

$$C(2) = 1$$

$$C(3) = 3$$

If  $n \geq 4$  and  $n$  is even then  $C(n) = 2C(\frac{n}{2}) + n - 1$ .

If  $n \geq 5$  and  $n$  is odd then  $C(n) = C((n-1)/2) + C((n+1)/2) + n - 1$ .

Hence

$$C(2^m) = 2C(2^{m-1}) + 2^m - 1.$$

For each of math we take a slightly worse bound:

$$C(2^m) = 2C(2^{m-1}) + 2^m.$$

Derive what  $C(2^m)$  is. Letting  $n = 2^m$  derive a bound on rough  $C(n)$  is for  $n$  a power of 2.