

CMSC433, Fall 2001 Software Components and Programming by Contract

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November 20, 2001

Administrivia

- Exams returned next Tuesday
 - answers posted
 - questions next Tuesday, after return
- Project 3
 - Public and private drivers posted
 - questions to TA
- Project 5
 - new due date Nov. 28
- Component readings posted on Readings page

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“Reusable code”

- Say I develop a String class in C++
- It is used in 5 different applications
- Each of which is statically linked
 - a self-contained executable

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Problems with static linking

- Redundant executable code
 - on disk and
 - in memory
- Can't update defective class without replacing application
- Can't add extensions of String class

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Dynamic linking

- Only one copy of DLL on disk
 - and in memory
- Can upgrade/replace library
 - all applications that use class get upgrade

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Dynamic linking problems

- All applications must use same link format and object layout
 - member alignment, virtual method table, virtual base classes, RTTI
- Class frozen in C++
 - adding variables or methods breaks everything
- Upgrade “breaks” your software
 - upgrade changed things it depended on

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Versioning

- Every time DLL changes, give it a new version ID
- Application only uses DLL version ID it was compiled against
- Problems:
 - multiple version bloat
 - inability to upgrade/enhance

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Compatible DLL's

- Binary compatibility
 - old code is able to invoke methods in new library
- Semantic compatibility
 - code that expected the old functionality won't be broken by the new functionality

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Binary compatibility

- Need to agree on link and runtime standards
- Can't make changes that would violate compatibility
 - e.g., add/delete variables/methods in C++

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Semantic compatibility

- Restrictions on changes in pre and post conditions
- Which can be “upgraded” to the other?
 - `int search1(int[] a, int x)`
 - pre: exists i s.t. $a[i] = x$
 - `int search2(int[] a, int x)`
 - pre: a sorted **and** exists i s.t. $a[i] = x$

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Which upgrades are compatible?

- Which can be “upgraded” to the other:
 - `int search1(int[] a, int x)`
 - pre: a sorted **and** exists i s.t. $a[i] = x$
 - post: $a[\text{retVal}] = x$
 - `int search2(int[] a, int x)`
 - pre: a sorted **and** exists i s.t. $a[i] = x$
 - post: $a[\text{retVal}] = x$
and $(\text{retVal} = 0 \text{ or } a[\text{retVal}-1] != x)$

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Rules for semantic compatibility

- f has pre: P and post: Q
- f' has pre: P' and post: Q'
- f can be upgraded to f' iff
 - $P \Rightarrow P'$
 - makes precondition looser
 - $Q' \Rightarrow Q$
 - makes postcondition stricter

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Problems with semantic compatibility

- Pre and post conditions are rarely specified
 - when they are, they are rarely complete and correct
- Programmers might depend upon conditions of your implementation rather than the interface
 - e.g., “I know that this function always returns a sorted array”

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Programming to an interface

- An interface is a contractual binding of two (or more) parties
- Should not depend on examining implementations
- Documentation is essential
 - e.g., for MiniServlet, the fact that the servlet is responsible for closing the output stream

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Indirect interfaces

- Function pointers, virtual method dispatch, callbacks
- A word processing component might invoke a grammar checker it didn't know existed

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What belongs in a contract?

- Pre and post conditions
- Resource requirements
 - execution time
 - memory requirements
- Leads-to
 - e.g., “invokes all registered notifiers”

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Callbacks

- In purely procedural word, library calls run to completion
- But allowing them to invoke methods in the caller code complicates matters
- A library function is often thought of as transitioning from one consistent state to another consistent state
 - possibly passing through inconsistent states

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Directory Service w/Callback

```
Module Directory
Type Notifier
  = Procedure (name: String);
Procedure AddNotifier(n : Notifier);
// pre: n != null
// post: n is registered, will be called on
//       AddEntry and RemoveEntry
Procedure AddEntry(n : String, f: File);
// pre: n != "" && f != null
// post: ThisFile(n) == f
Function ThisFile(n : name) : File
// pre: n != ""
// post: result = file named n
//       or result = null and no such file

Module DirectoryDisplay
Procedure myNotifier(n : String) {
  if Directory.ThisFile(n) == null
    // delete n in display
  else
    // n has been added, display it
  }
..
Directory.AddNotifier(myNotifier);
...
```

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What is missing?

- Precondition for notifier
 - that state of directory reflects change being notified about
 - or that it doesn't
- Postcondition for notifier
 - that it doesn't change the state of the directory

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When is the notifier called?

- Before or after change?
- When file is renamed, how many times?
- When file is replaced, how many times?

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Why is changing state bad?

- Say we attach a notifier that, whenever an entry with the name X.htm is added, renames it to X.html
 - violates postcondition of AddEntry

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Locking doesn't help

- Could make Add/RemoveEntry lock directory
 - wouldn't help
 - notifier, and subsequent potential modification, are all done in same thread
- If locks weren't recursive, would just result in deadlock

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Thread spawning?

- One way to handle renaming notifier
- Use locks
- When notified that X.htm has been added, spawn thread to do renaming
 - will have to wait to obtain lock
- What is the pre/post condition of a method that spawns a thread?

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inProgress functions

- Add method that detects whether an operation is in progress
- Add precondition for operation that no other operation is in progress

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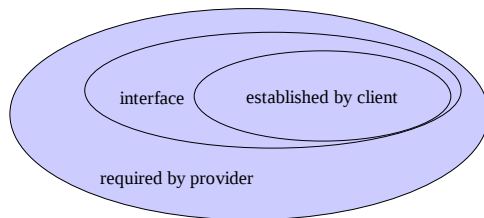
Substitutability

- Clients use an interface to invoke a method in a provider
 - pre: established by client
 - => demanded by interface
 - => required by provider
 - post: established by provider
 - => demanded by interface
 - => required by client

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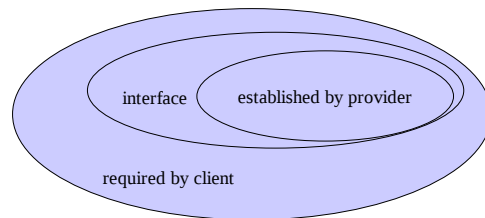
subtyping of preconditions



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subtyping of postconditions



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Co- and Contra- variance

- Co-variance = vary in same way
- Contra-variance = vary in opposite ways
- Co-variant return types are sound and useful
 - Supported in C++ but not in Java
- Contra-variant arguments are sound but rarely useful
 - Not supported in C++ or Java

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Co and Contra variance

```
interface View {  
    Model getModel();  
}  
interface TextView extends View {  
    TextModel getModel();  
}  
interface GraphicsView extends TextView {  
    GraphicsModel getModel();  
}
```

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Co and Contra variance

```
interface View {  
    Model getModel();  
    void setModel(Model);  
}  
interface TextView extends View {  
    TextModel getModel();  
    void setModel(TextModel);  
}
```

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Structural or Declared Subtyping

- Structural subtyping
 - B is a subtype of A iff B supports the interface of A
- Declared subtyping
 - B is a subtype of A iff B is declared as a subtype of A

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Problems

- With declared subtyping
 - Can't write supertype of third-party class
 - e.g., `TextView`, subtype of `TextObserver` and `View`
- With structural subtyping
 - doesn't check (implied) pre and post conditions
 - Same names doesn't imply same meaning

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Three Facets of Inheritance

- Implementation inheritance
 - reuse of code
- Interface inheritance
 - reuse of signatures
- Promise of substitutability
 - not in SmallTalk nor Eiffel
 - In Eiffel, can undefine methods

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Fragile base class problem

- Syntactic fragile base class problem
 - Changing a class requires that all subclasses be recompiled
 - offsets of fields and methods changed
- Semantic fragile base class problem
 - Recompiling may not fix things

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Semantic fragile base class problem

- Lots of calls between superclass code and subclass code
 - check those pre and post conditions
- Back and forth nature creates problems
 - code in superclass invokes code in subclass which invokes code in superclass
- Super/subclass setting/getting variables directly

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Solutions

- Use private variables and setter/getter methods
- Use protected members when appropriate
- Figure out groups of mutually dependent methods
 - if A invokes B and B invokes A, then A and B are mutually dependent
 - have to override an entire set of MD methods

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Multiple inheritance

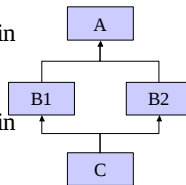
- multiple supertyping
 - inherit interface from multiple supertypes
 - promise substitutability for all of them
- inheritance from multiple supertypes is tricky

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Multiple inheritance of implementations

- What if a function is defined in both B1 and B2?
 - what version does C invoke?
- What if a function is defined in both A and B1?
 - what version does B2 invoke?
- Do B1 and B2 share an A or each have their own A?



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More issues?

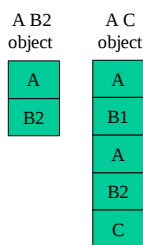
- Does it matter which is declared first, B1 or B2?
- In C++, B1 and B2 share an A if both declare A as a virtual base class
 - Don't know that B1 and B2 might need to share an A until class C is written
 - virtual base classes are less efficient

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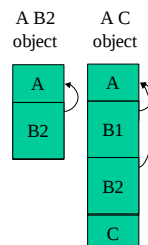
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Virtual base classes

A as a non-virtual base class



A as a virtual base class



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