

CMSC421 Fall 2005 Homework 2

October 17, 2005

1 Problem 1: Constraint Satisfaction (25 pts)

- a. (i) See Figure 1.
(ii) 1=B, 2=G, 3=R, 4=R, 5=B.
(iii) Count number of nodes expanded: 8.
- b. (i) See Table 1.
(ii) 1=B, 2=G, 3=R, 4=R, 5=B
(iii) 5.
- c. (i) $3^5 = 243$.
(ii) Take the circle formed by region 1,2,4. As they are mutually connected to each other, each region has to take one unique color. The second circle formed by region 2,4,5 then determine the color of region 5. Region 2 and 5 determines the color of region 3. Therefore, $3! = 6$.

2 Problem 2: Game Playing (15 pts)

- a. See Figure 2.
b. (i) All leaf nodes are terminal nodes.
(ii) See Figure 2.
(iii) From left to right, number the root's children as node 1,2,3,4. Those nodes in the subtrees rooted at node 2,3,4 (including node 2,3,4) will be pruned.

Table 1: Forward Checking

	1	2	3	4	5
	BGR	BGR	BGR	BGR	BGR
1 = B	B	GR	BGR	GR	BGR
2 = G	B	G	BR	R	BR
4 = R	B	G	BR	R	B
5 = B	B	G	R	R	B
3 = R	B	G	R	R	B

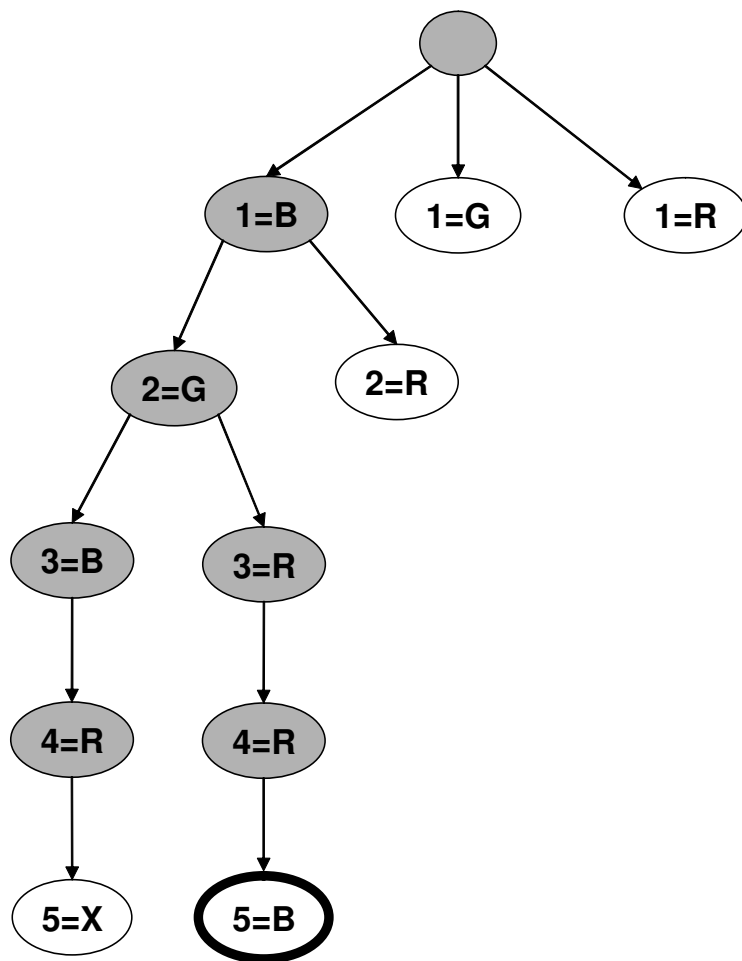


Figure 1: Backtracking Tree

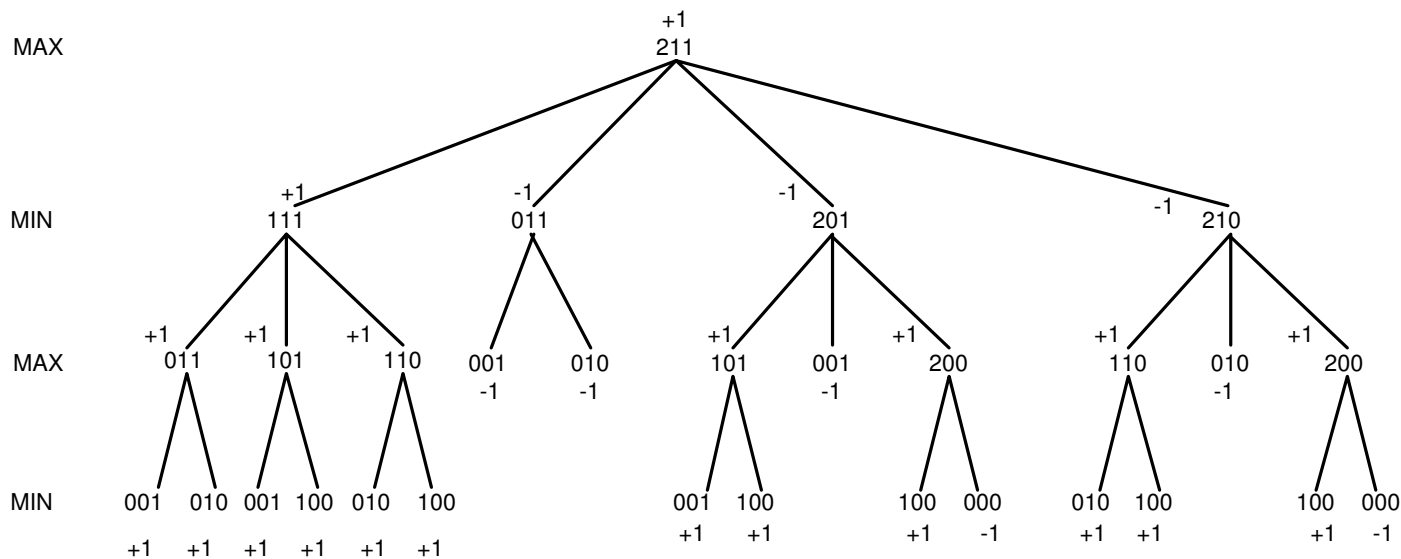


Figure 2: Game Tree of Nim211

3 Problem 3: Propositional Logic - Models (9 pts)

These can be computed by counting the rows in a truth table that come out true. Remember to count the propositions that are not mentioned; if a sentence mentions only A and B , then we multiply the number of models for $\{A, B\}$ by 2^2 to account for C and D .

- 6
- 12
8. The correct interpretation is $A \Leftrightarrow (B \Leftrightarrow C)$, therefore the number of models is 4×2 . However, if you interpreted it as $(A \Leftrightarrow B) \Leftrightarrow C$ or were explicit about this, you also got credit.

4 Problem 4: Propositional Logic - Validity(16 pts)

We could use truth table or applying equivalence rules to determine validity of a statement.

- Valid since $\text{Smoke} \Rightarrow \text{Smoke} \equiv \sim \text{Smoke} \vee \text{Smoke}$ is a tautology.
- Neither (or Satisfiable) by examining the truth table 2.
- Neither (or Satisfiable).
- Valid.
- Valid.
- Valid.
- Valid.

Table 2: Truth Table of $Smoke \Rightarrow Fire$

$Smoke$	$Fire$	$Smoke \Rightarrow Fire$
T	T	T
T	F	F
F	T	T
F	F	T

h. Neither (or Satisfiable).

5 Problem 5: Propositional Logic - Resolution (30 pts)

From the first two statements, we see that if it is mythical, then it is immortal; otherwise it is a mammal. So it must be either immortal or a mammal, and thus horned. That means it is also magical. However, we can't deduce anything about whether it is mythical.

First, we introduce some symbols:

Y: Unicorn is mYthical

I: Unicorn is Immortal

M: Unicorn is a Mammal

H: Unicorn is Horned

G: Unicorn is maGical

Therefore, the propositional knowledgebase will be

$$\begin{aligned}
 Y &\Rightarrow I \\
 \sim Y &\Rightarrow M \\
 I \vee M &\Rightarrow H \\
 H &\Rightarrow G
 \end{aligned}$$

The final knowledgebase will be as following after applying Implication Elimination, DeMorgan's and Distributivity logical equivalences laws:

1. $\sim Y \vee I$
2. $Y \vee M$
3. $\sim I \vee H$
4. $\sim M \vee H$
5. $\sim H \vee G$

Now, we'd like to know if Y? G? H?

Using Resolution Refutation, we introduce $\sim Y$ into the knowledgebase and hope that we can get false in order to prove Y . However, it doesn't produce false (Note that we can't say Q is not entailed by the knowledgebase). Hence, we could infer nothing on mythical.

Include $\sim G$ into the knowledgebase.

1. $\sim Y \vee I$
2. $Y \vee M$
3. $\sim I \vee H$
4. $\sim M \vee H$
5. $\sim H \vee G$
6. $\sim G$ *negated goal*
7. $\sim H$ 5, 6
8. $\sim I$ 3, 7
9. $\sim Y$ 1, 8
10. M 2, 9
11. H 4, 10
12. *false* 7, 11

We can prove that unicorn is magical. Similarly, we can prove unicorn is horned.

6 Problem 6: First Order Logic - Representation (25 pts)

There are different ways of doing this.

Let the basic vocabulary be as follows:

Takes(x, c, s): student x takes course c in semester s ;

Passes(x, c, s): student x passes course c in semester s ;

Score(x, c, s): the score obtained by student x in course c in semester s ;

$x > y$: x is greater than y .

F and G : specific French and Greek courses;

Buys(x, y, z): x buys y from z ;

Sells(x, y, z): x sells y to z ;

Shaves(x, y): person x shaves person y ;

Born(x, c): person x is born in country c ;

Parent(x, y): x is a parent of y ;

Citizen(x, c, r): x is a citizen of country c for reason r ;

Resident(x,c): x is a resident of country c ;
Fools(x,y,t): person x fools person y at time t ;
Student(x), Person(x), Man(x), Barber(x), Expensive(x), Agent(x), Insured(x), Smart(x), Politician(x):
 predicates satisfied by members of the corresponding categories.

- a. Some students took French in spring 2001.
 $\exists x \text{ Student}(x) \wedge \text{Takes}(x, F, \text{Spring2001})$.
- b. Every student who takes French passes it.
 $\forall x, s \text{ Student}(x) \wedge \text{Takes}(x, F, s) \Rightarrow \text{Passes}(x, F, s)$.
- c. Only one student took Greek in spring 2001.
 $\exists x \text{ Student}(x) \wedge \text{Takes}(x, G, \text{Spring2001}) \wedge \forall y y \neq x \Rightarrow \sim \text{Takes}(y, G, \text{Spring2001})$.
- d. The best score in Greek is always higher than the best score in French.
 $\forall s \exists x \forall y \text{ Score}(x, G, s) > \text{score}(y, F, s)$.
- e. Every person who buys a policy is smart.
 $\forall x \text{ Person}(x) \wedge (\exists y, z \text{ Policy}(y) \wedge \text{Buys}(x, y, z)) \Rightarrow \text{Smart}(x)$.
- f. No person buys an expensive policy.
 $\forall x, y, z \text{ Person}(x) \wedge \text{Policy}(y) \wedge \text{Expensive}(y) \Rightarrow \sim \text{Buys}(x, y, z)$.
- g. There is an agent who sells policies only to people who are not insured.
 $\exists x \text{ Agent}(x) \wedge \forall y, z \text{ Policy}(y) \wedge \text{Sells}(x, y, z) \Rightarrow (\text{Person}(z) \wedge \sim \text{Insured}(z))$.
- h. There is a barber who shaves all men in town who do not shave themselves.
 $\exists x \text{ Barber}(x) \wedge \forall y \text{ Man}(y) \wedge \sim \text{Shaves}(y, y) \Rightarrow \text{Shaves}(x, y)$.
- i. A person born in the UK, each of whose parents is a UK citizen or a UK resident, is a UK citizen by birth.
 $\forall x \text{ Person}(x) \wedge \text{Born}(x, UK) \wedge (\forall y \text{ Parent}(y, x) \Rightarrow ((\exists r \text{ Citizen}(y, UK, r)) \vee \text{Resident}(y, UK))) \Rightarrow \text{Citizen}(x, UK, \text{Birth})$.
- j. A person born outside the UK, one of whose parents is a UK citizen by birth, is a UK citizen by descent.
 $\forall x \text{ Person}(x) \wedge \sim \text{Born}(x, UK) \wedge (\exists y \text{ Parent}(y, x) \wedge \text{Citizen}(y, UK, \text{Birth})) \Rightarrow \text{Citizen}(x, UK, \text{Descent})$.
- k. Politicians can fool some of the people all of the time, and they can fool all of the people some of the time, but they can't fool all of the people all of the time.
 $\forall x \text{ Politician}(x) \Rightarrow$
 $(\exists y \forall t \text{ Person}(y) \wedge \text{Fools}(x, y, t)) \wedge$
 $(\exists t \forall y \text{ Person}(y) \Rightarrow \text{Fools}(x, y, t)) \wedge$

$$\sim (\forall t \forall y \textit{Person}(y) \Rightarrow \textit{Fools}(x, y, t)).$$