

# CMSC 132: Object-Oriented Programming II

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## **Synchronization in Java**

**Department of Computer Science**  
**University of Maryland, College Park**

# Multithreading Overview

## ■ Motivation & background

## ■ Threads

### ■ Creating Java threads

### ■ Thread states

### ■ Scheduling

## ■ Synchronization

### ■ Data races ←

### ■ Locks

### ■ Deadlock



# Data Race

## ■ Definition

- Concurrent accesses to same shared variable, where at least one access is a write

## ■ Properties

- Order of accesses may change result of program
- May cause intermittent errors, very hard to debug

## ■ Example

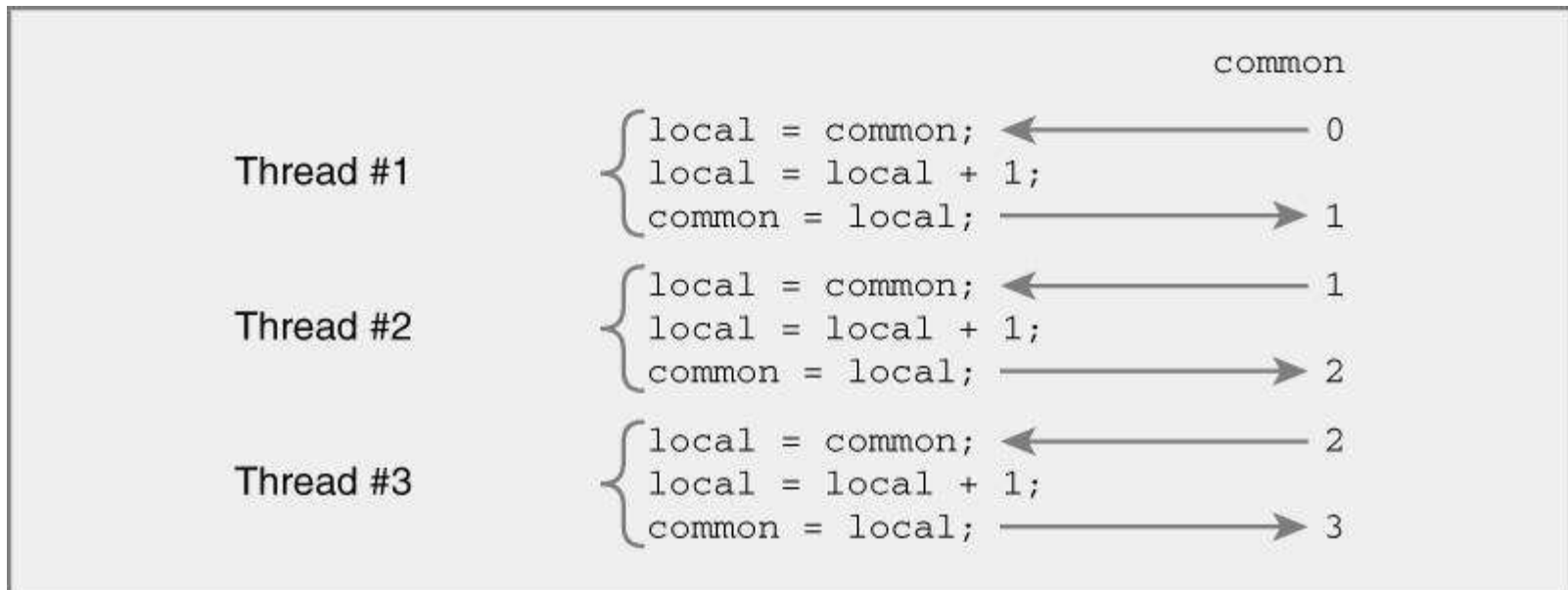
```
public class DataRace extends Thread {  
    static int x;    // shared variable x causing data race  
    public void run() { x = x + 1; } // access to x  
}
```

# Data Race Example

```
public class DataRace extends Thread {
    static int common = 0;
    public void run() {
        int local = common; // data race
        local = local + 1;
        common = local; // data race
    }
    public static void main(String[] args) throws InterruptedException {
        int max = 3;
        DataRace[] allThreads = new DataRace[max];
        for (int i = 0; i < allThreads.length; i++)
            allThreads[i] = new DataRace();
        for (DataRace t : allThreads)
            t.start();
        for (DataRace t : allThreads)
            t.join();
        System.out.println(common); // may not be 3
    }
}
```

# Data Race Example

## ■ Sequential execution output



# Data Race Example

## ■ Concurrent execution output (possible case)

Thread #1: `local = common;` ← 0  
Thread #2: `local = common;` ← 0  
Thread #3: `local = common;` ← 0  
Thread #1: `local = local + 1;`  
Thread #2: `local = local + 1;`  
Thread #3: `local = local + 1;`  
Thread #1: `common = local;` → 1  
Thread #2: `common = local;` → 1  
Thread #3: `common = local;` → 1

■ Result depends on thread execution order!

# Synchronization

## ■ Definition

- Coordination of events with respect to time

## ■ Properties

- May be needed in multithreaded programs to eliminate **data races**
- Incurs runtime overhead
- Excessive use can reduce performance

# Lock

## ■ Definition

- Entity can be held by only one thread at a time

## ■ Properties

- A type of synchronization
- Used to enforce **mutual exclusion**
  - Thread can acquire / release locks
  - Only 1 thread can acquire lock at a time
- Thread will wait to acquire lock (stop execution)
  - If lock held by another thread
- Used to implement **monitors**
  - Only 1 thread can execute (locked) code at a time



# Synchronized Objects in Java

## ■ Java objects provide locks

### ■ Apply **synchronized** keyword to object

■ Will acquire / release lock associated with object

### ■ Mutual exclusion for code in synchronization **block**

## ■ Example

Object x = new Object();

**block** {  
    synchronized(x) {   // acquire lock on x on entry  
        ...           // hold lock on x in block  
    }                   // release lock on x on exit  
}

# Synchronized Methods In Java

- Apply synchronized keyword to method
  - Mutual exclusion for entire body of method
  - Synchronizes on object invoking method

- Example

**block**  


```
synchronized foo() { ...code... }
```

**// shorthand notation for**

```
foo() {  
    ↓  
    synchronized (this) { ...code... }  
}
```

# Synchronized Methods In Java

```
public synchronized void enqueue( Object item ) {  
    // Body of method goes here  
}
```



*Shorthand notation for*

```
public void enqueue( Object item ) {  
    synchronized ( this ) {  
        // Body of method goes here  
    }  
}
```

# Locks in Java

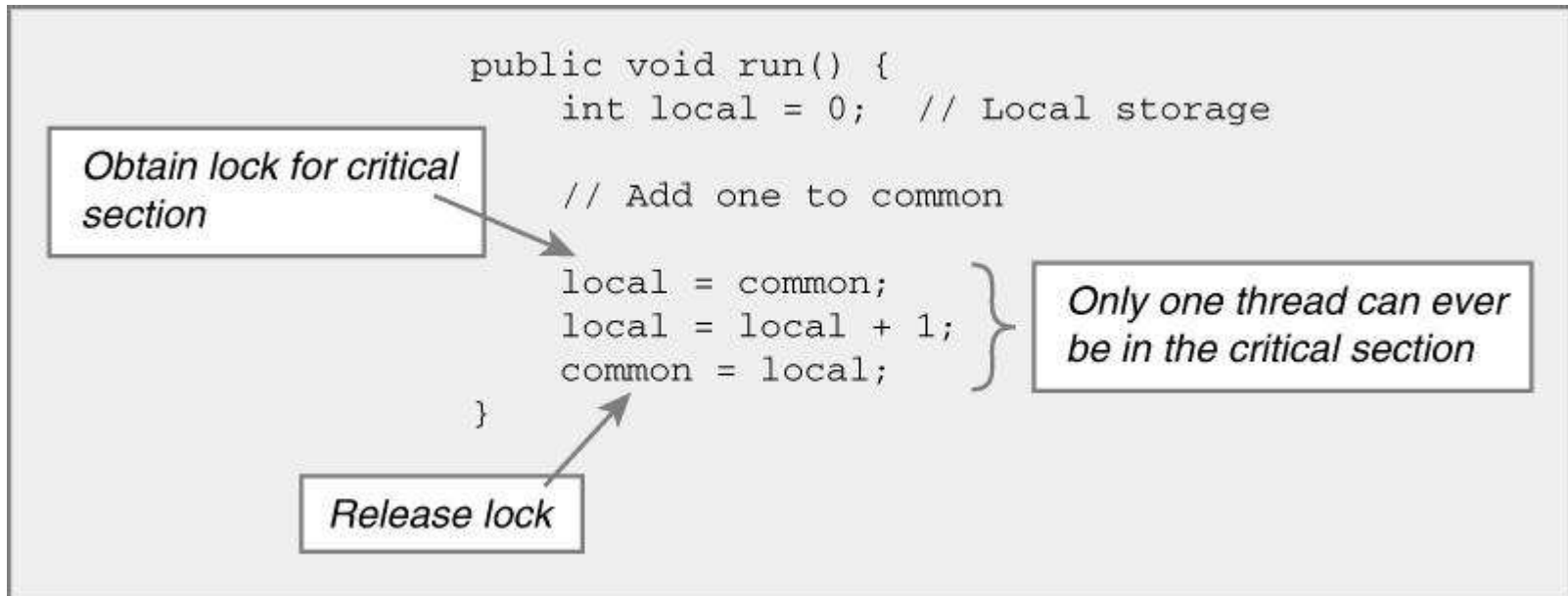
## ■ Properties

- No other thread can get lock on x while in block
- Other threads can still access/modify x!
- Locked block of code  $\Rightarrow$  **critical section**

## ■ Lock is released when block terminates

- End of block reached
- Exit block due to return, continue, break
- Exception thrown

# Synchronization Example



# Lock Example

```
public class DataRace extends Thread {  
    static int common = 0;  
    static Object o;           // all threads use o's lock  
    public void run() {  
        synchronized(o) {    // single thread at once  
            int local = common; // data race eliminated  
            local = local + 1;  
            common = local;  
        }  
    }  
    public static void main(String[] args) {  
        o = new Object();  
        ...  
    }  
}
```

# Synchronization Issues

1. Use same lock to provide mutual exclusion
2. Ensure atomic transactions
3. Avoiding deadlock

# Issue 1) Using Same Lock

## ■ Potential problem

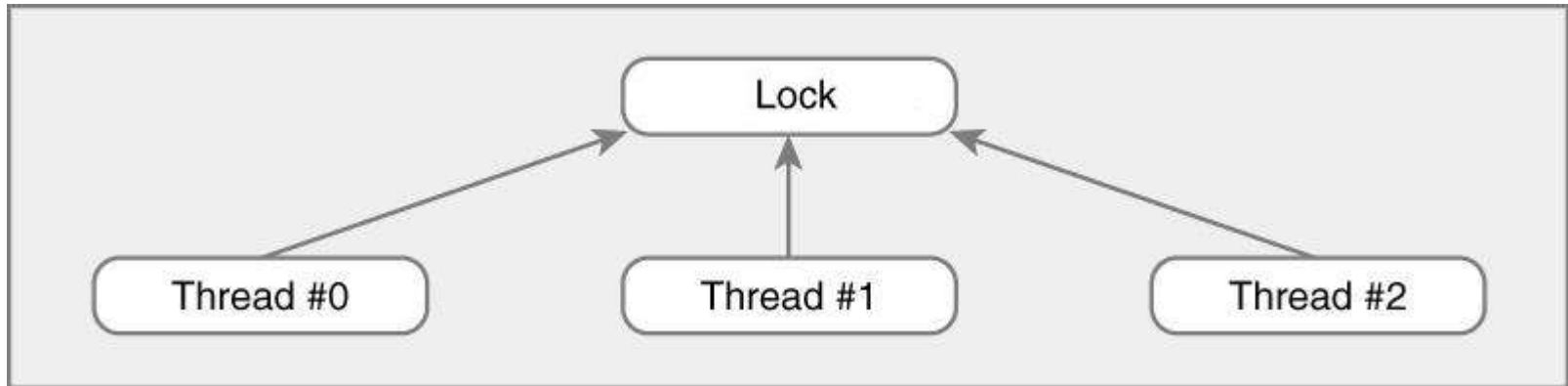
- Mutual exclusion depends on threads acquiring same lock
- No synchronization if threads have different locks

## ■ Example

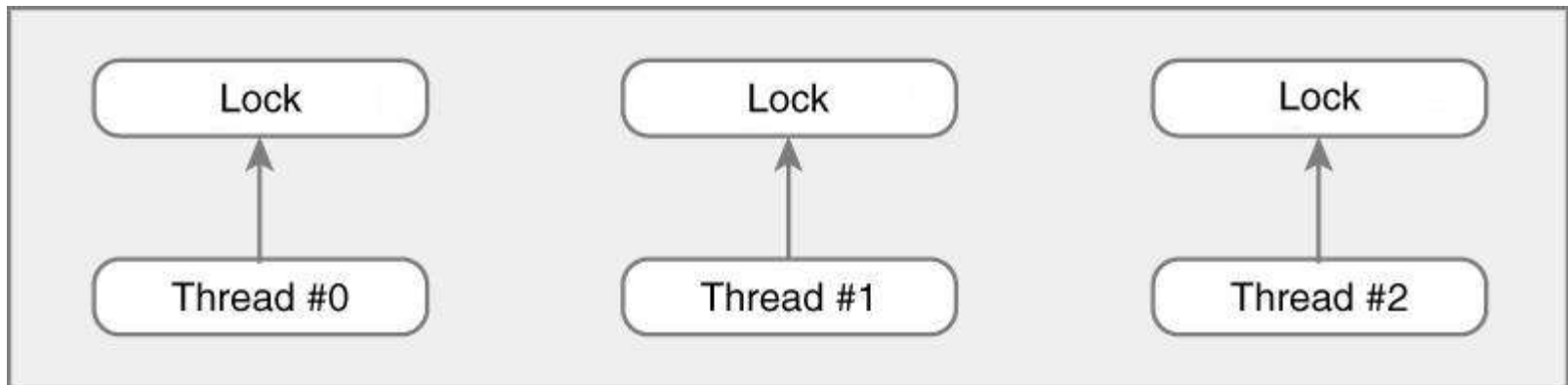
```
foo() {  
    Object o = new Object(); // different o per thread  
    synchronized(o) {  
        ... // potential data race  
    }  
}
```

# Locks in Java

- **Single lock for all threads (mutual exclusion)**



- **Separate locks for each thread (no synchronization)**



# Lock Example – Incorrect Version

```
public class DataRace extends Thread {  
    static int common = 0;  
    public void run() {  
        Object o = new Object(); // different o per thread  
        synchronized(o) {  
            int local = common;    // data race  
            local = local + 1;  
            common = local;        // data race  
        }  
    }  
    public static void main(String[] args) {  
        ...  
    }  
}
```

# Issue 2) Atomic Transactions

## ■ Potential problem

- Sequence of actions must be performed as single **atomic transaction** to avoid data race
- Ensure lock is held for duration of transaction

## ■ Example

```
synchronized(o) {  
    int local = common;  
    local = local + 1;  
    common = local;  
}
```

} // all 3 statements must  
// be executed together  
// by single thread

# Lock Example – Incorrect Version

```
public class DataRace extends Thread {  
    static int common = 0;  
    static Object o;           // all threads use o's lock  
    public void run() {  
        int local;  
        synchronized(o) {  
            local = common;  
        }  
        synchronized(o) {  
            local = local + 1;  
            common = local;  
        }  
    }  
}
```

} // transaction not atomic  
// data race may occur  
// even using locks

# Issue 3) Avoiding Deadlock

## ■ Potential problem

- Threads holding lock may be unable to obtain lock held by other thread, and vice versa
- Thread holding lock may be waiting for action performed by other thread waiting for lock
- Program is unable to continue execution (**deadlock**)

# Deadlock Example 1

Object **a**;

Object **b**;

```
Thread1() {  
    synchronized(a) {  
        synchronized(b) {  
            ...  
        }  
    }  
}
```

```
Thread2() {  
    synchronized(b) {  
        synchronized(a) {  
            ...  
        }  
    }  
}
```

// Thread1 holds lock for a, waits for b

// Thread2 holds lock for b, waits for a

## Deadlock Example 2

```
void swap(Object a, Object b) {  
    Object local;  
    synchronized(a) {  
        synchronized(b) {  
            local = a; a = b; b = local;  
        }  
    }  
}
```

```
Thread1() { swap(a, b); } // holds lock for a, waits for b  
Thread2() { swap(b, a); } // holds lock for b, waits for a
```

# Deadlock

## ■ Avoiding deadlock

- In general, avoid holding lock for a long time
- Especially avoid trying to hold two locks
  - May wait a long time trying to get 2<sup>nd</sup> lock

# Thread-safe

- **Thread-safe** – Code is considered thread-safe if it works correctly when executed by multiple threads simultaneously.
- **Example: ArrayList is not thread-safe**

**From Java API:** “Note that this implementation is not synchronized. If multiple threads access an ArrayList instance concurrently, and at least one of the threads modifies the list structurally, it *must* be synchronized externally.”

# Synchronization Summary

- Needed in multithreaded programs
- Can prevents data races
- Java objects support synchronization
- Many other tricky issues
  - To be discussed in future courses