

# CMSC 330: Organization of Programming Languages

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## Logic Programming with Prolog

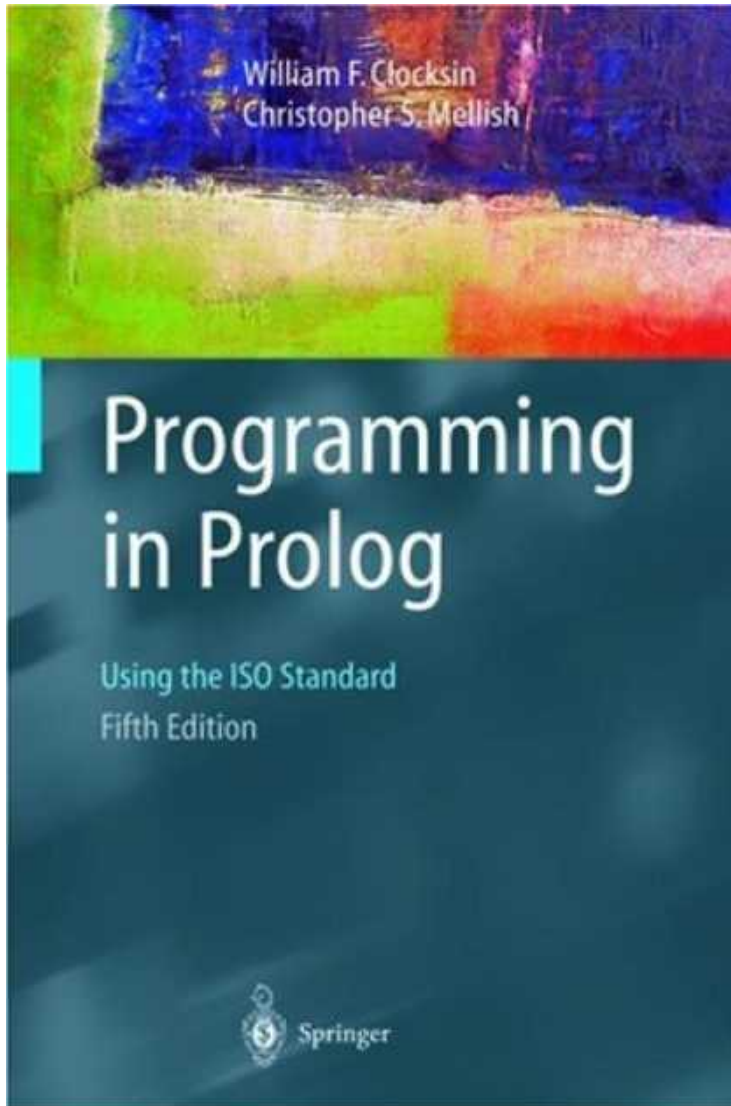
# Background

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- ▶ 1972, University of Aix-Marseille
- ▶ Original goal: Natural language processing
- ▶ At first, just an interpreter written in Algol
  - Compiler created at Univ. of Edinburgh

# More Information On Prolog

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- ▶ Various tutorials available online
- ▶ Links on webpage

# Logic Programming

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- ▶ At a high level, logic programs model the relationship between “objects”
  1. Programmer specifies relationships at a high level
  2. Language builds a database
  3. Programmer then queries this database
  4. Language searches for answers

# Features of Prolog

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- ▶ Declarative
  - Specify what goals you want to prove, not how to prove them (mostly)
- ▶ Rule based
- ▶ Dynamically typed
- ▶ Several built-in datatypes
  - Lists, numbers, records, ... but no functions
- ▶ Several other logic programming languages
  - Datalog is simpler; CLP and  $\lambda$ Prolog more feature-ful
  - Erlang borrows some features from Prolog

# A Small Prolog Program – Things to Notice

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Use `/* */` for comments, or `%` for 1-liners

Period ends statements

Lowercase logically terminates

Program consists of facts and rules

Uppercase denotes variables

```
/* A small Prolog program */  
female(alice).  
male(bob).  
male(charlie).  
father(bob, charlie).  
mother(alice, charlie).  
  
% "X is a son of Y"  
son(X, Y) :- father(Y, X), male(X).  
son(X, Y) :- mother(Y, X), male(X).
```

# Running Prolog (Interactive Mode)

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Navigating location and loading program at top level

```
?- working_directory(C,C). ← Find current directory  
C = 'c:/windows/system32/' .
```

```
?- working_directory(C,'c:/Users/me/desktop/p6') . ← Set directory  
C = 'c:/Users/me/desktop/' .
```

```
?- ['01-basics.pl'] . ← Load file 01-basics.pl  
% 01-basics.pl compiled 0.00 sec, 17 clauses  
true.
```

```
?- make. ← Reload modified files; replace rules  
true.
```

# Running Prolog (Interactive Mode)

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Listing rules and entering queries at top level

```
?- listing(son). ← List rules for son
```

```
son(X, Y) :-  
    father(Y, X),  
    male(X).
```

```
son(X, Y) :-  
    mother(Y, X),  
    male(X).
```

```
true.
```

```
?- son(X,Y).
```

```
X = charlie,
```

```
Y = bob;
```

```
X = charlie,
```

```
Y = alice.
```

User types ; to request additional answer

Multiple answers

User types return to complete request



# Style

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One predicate per line

```
blond(X) :-  
  father(Father, X),  
  blond(Father),      % father is blond  
  mother(Mother, X),  
  blond(Mother).     % and mother is blond
```

Descriptive variable names

Inline comments with % can be useful

# Outline

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- ▶ Syntax, terms, examples
- ▶ Unification
- ▶ Arithmetic / evaluation
- ▶ Programming conventions
- ▶ Goal evaluation
  - Search tree, clause tree
- ▶ Lists
- ▶ Built-in operators
- ▶ Cut, negation

# Prolog Syntax and Terminology

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## ► Terms

- **Atoms:** begin with a lowercase letter

horse    underscores\_ok    numbers2

- **Numbers**

123    -234    -12e-4

- **Variables:** begin with uppercase or `_` “don't care” variables

X    Biggest\_Animal    `_the_biggest1`    \_

- **Compound terms:** functor(arguments)

bigger(horse, duck)

bigger(X, duck)

f(a, g(X, \_), Y, \_)

No blank spaces between functor and (arguments)

# Prolog Syntax and Terminology (cont.)

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## ► Clauses

- **Facts:** define predicates, terminated by a period  
    `bigger(horse, duck).`  
    `bigger(duck, gnat).`

Intuitively: “this particular relationship is true”

- **Rules:** Head :- Body  
    `is_bigger(X,Y) :- bigger(X,Y).`  
    `is_bigger(X,Y) :- bigger(X,Z), is_bigger(Z,Y).`

Intuitively: “Head if Body”, or “Head is true if each of the **subgoals** can be shown to be true”

## ► A **program** is a sequence of clauses

# Prolog Syntax and Terminology (cont.)

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## ► Queries

- To “run a program” is to submit queries to the interpreter
- Same structure as the body of a rule
  - Predicates separated by commas, ended with a period
- Prolog tries to determine whether or not the predicates are true

?- is\_bigger(horse, duck).

?- is\_bigger(horse, X).

“Does there exist a substitution for  $X$  such that `is_bigger(horse,X)?`”

*Without which, nothing*

# Unification – The Sine Qua Non of Prolog

▶ Two terms unify **if and only if**

- They are identical

?- gnat = gnat.  
true.

- They can be made identical by **substituting** variables

?- is\_bigger(X, gnat) = is\_bigger(horse, gnat).

X = horse.

} This is the substitution: what X must be for the two terms to be identical.

?- pred(X, 2, 2) = pred(1, Y, X)

false

Sometimes there are multiple possible substitutions; Prolog can be asked to enumerate them all

?- pred(X, 2, 2) = pred(1, Y, \_)

X = 1,

Y = 2.

# The = Operator

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- ▶ For unification (matching)
- ▶ `?- 9 = 9.`  
`true.`
- ▶ `?- 7 + 2 = 9.`  
`false.`
- ▶ Why? Because these terms do not match
  - `7+2` is a compound term (e.g., `+(7,2)`)
- ▶ Prolog does not evaluate either side of =
  - Before trying to match

# The **is** Operator

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- ▶ For arithmetic operations
- ▶ “LHS **is** RHS”
  - First **evaluate** the RHS (and RHS only!) to value  $V$
  - Then match:  $LHS = V$

## ▶ Examples

?- 9 is 7+2.  
true.

?- 7+2 is 9.  
false.

?- X = 7+2.  
X = 7+2.

?- X is 7+2.  
X = 9.



# No Assignment

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▶ = and is operators do not perform assignment

▶ Example

- `foo(...,X) :- ... X = 1,...` % true only if X = 1
- `foo(...,X) :- ... X = 1, ..., X = 2, ...` % always fails
- `foo(...,X) :- ... X is 1,...` % true only if X = 1
- `foo(...,X) :- ... X is 1, ..., X is 2, ...` % always fails

X can't be unified with 1 & 2 at the same time

# Function Parameter & Return Value

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## ► Code example

increment(X,Y) :-  
    Y is X+1.

Parameter  
Return value

?- increment(1,Z). ← Query

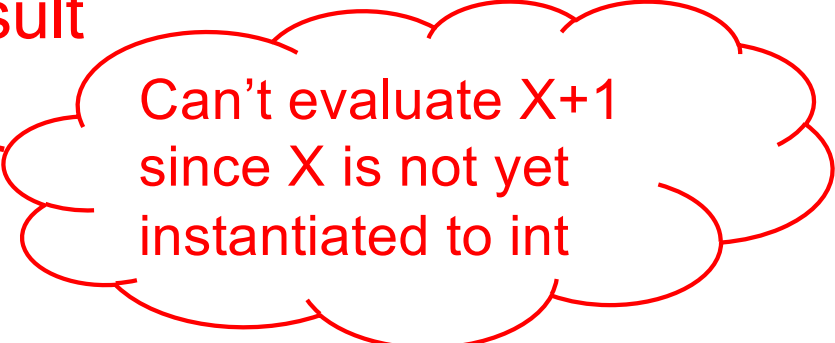
Z = 2. ← Result

?- increment(1,2).

true.

?- increment(Z,2).

Can't evaluate X+1  
since X is not yet  
instantiated to int



**ERROR: incr/2: Arguments are not sufficiently instantiated**

# Function Parameter & Return Value

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## ▶ Code example

`addN(X,N,Y) :-`  
`Y is X+N.`

Parameters  
Return value

The diagram shows two red dashed arrows pointing from the labels 'Parameters' and 'Return value' to the arguments 'X', 'N', and 'Y' in the function definition. The 'Parameters' label is connected to 'X' and 'N', while the 'Return value' label is connected to 'Y'.

`?- addN(1,2,Z).` ← Query  
`Z = 3.` ← Result

The diagram shows two red solid arrows pointing from the labels 'Query' and 'Result' to the query and result lines respectively.

# Recursion

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► Code example

`addN(X,0,X).` ← Base case  
`addN(X,N,Y) :-` ← Inductive step  
    `X1 is X+1,`  
    `N1 is N-1,`  
    `addN(X1,N1,Y).` ← Recursive call

?- `addN(1,2,Z).`  
`Z = 3.`

# Factorial

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▶ Code

factorial(0,1).

factorial(N,F) :-

    N > 0,

    N1 is N-1,

    factorial(N1,F1),

    F is N\*F1.

# Tail Recursive Factorial w/ Accumulator

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- ▶ Code

```
tail_factorial(0,F,F).  
tail_factorial(N,A,F) :-  
    N > 0,  
    A1 is N*A,  
    N1 is N - 1,  
    tail_factorial(N1,A1,F).
```

# And and OR

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## ► And

- To implement  $X \ \&\& \ Y$  (use  $,$  in body of clause)
- Example  
 $Z \text{ :- } X, Y.$

## ► OR

- To implement  $X \ || \ Y$  (use two clauses)
- Example  
 $Z \text{ :- } X.$   
 $Z \text{ :- } Y.$

# Goal Execution

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- ▶ When submitting a query, we ask Prolog to substitute variables as necessary to make it true
- ▶ Prolog performs **goal execution** to find a solution
  - Start with the goal
  - Try to unify the head of a rule with the current goal
  - The rule hypotheses become subgoals
    - Substitutions from one subgoal constrain solutions to the next
  - If it reaches a dead end, it **backtracks**
    - Tries a different rule
  - When it can backtrack no further, it reports **false**
- ▶ More advanced topics later – cuts, negation, etc.



# Goal Execution (cont.)

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- ▶ Consider the following:
  - “All men are mortal”  
`mortal(X) :- man(X).`
  - “Socrates is a man”  
`man(socrates).`
  - “Is Socrates mortal?”  
`?- mortal(socrates).`  
`true.`
- ▶ How did Prolog infer this?
  1. Sets `mortal(socrates)` as the initial goal
  2. Sees if it unifies with the head of any clause:  
`mortal(socrates) = mortal(X).`
  3. `man(socrates)` becomes the new goal (since `X=socrates`)
  4. Recursively scans through all clauses, **backtracking** if needed ...

# Clause Tree

- ▶ Clause tree
  - Shows (recursive) evaluation of all clauses
  - Shows value (instance) of variable for each clause
  - Clause tree is true if all leaves are true

- ▶ Factorial example

factorial(0,1).

factorial(N,F) :-

N > 0,

N1 is N-1,

factorial(N1,F1),

F is N\*F1.

