Fall 2016 Kruskal

CMSC 351: Practice Questions for Final Exam

These are practice problems for the upcoming final exam. You will be given a sheet of notes for the exam. Also, go over your homework assignments. **Warning:** This does not necessarily reflect the length, difficulty, or coverage of the actual exam.

Problem 1. Let $A[1,..,n]$ be an array of n numbers (some positive and some negative).

(a) Give an algorithm to find which *three* numbers have sum closest to zero. Make your algorithm as efficient as possible. Write it in pseudo-code.

(b) Analyze its running time.

Problem 2. Assume that you developed an algorithm to find the (index of the) $n/3$ smallest element of a list of n elements in $2n$ comparisons.

(a) Using the algorithm (as a black box), give an algorithm, efficient in the worst case, to find the kth smallest element of a list.

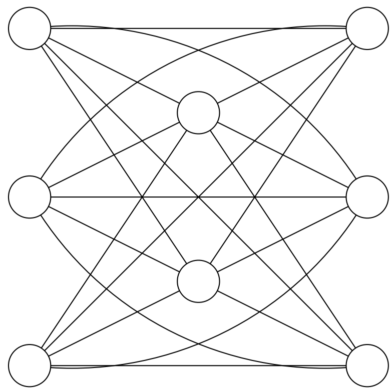
(b) Write down a recurrence for (a bound on) the number of comparisons it executes in the worst case.

(c) Solve the recurrence (using constructive induction). Find the high order term exactly (but you do not need any low order terms).

(d) Using the (black box) algorithm for finding the $n/3$ smallest element and using the ideas and results of Parts (a), (b), and (c), give an efficient algorithm to find (the index of) two elements, the k_1th smallest and the k_2 smallest (for inputs k_1 and k_2). The algorithm description can be very high level and brief.

(e) How many comparisons does it use? Find the high order term exactly (but you do not need any low order terms). Give a brief justification.

Problem 3. A graph is tripartite if the vertices can be partitioned into three sets so that there are no edges internal to any set. The *complete* tripartite graph, $K(a, b, c)$, has three sets of vertices with sizes a, b, and c and all possible edges between each pair of sets of vertices. $K(3, 2, 3)$ is pictured below. A *Hamiltonian* cycle in a graph is a cycle that traverses every vertex exactly once.



(a) For which values of n does $K(1, 1, n)$ have a Hamiltonian cycle. Justify your answer.

(b) For which values of n does $K(1, n, n)$ have a Hamiltonian cycle. Justify your answer.

(c) For which values of n does $K(n, n, n)$ have a Hamiltonian cycle. Justify your answer.

Problem 4. Let $G = (V, E)$ be an undirected graph. A *triangle* is a set of three vertices such that each pair has an edge.

(a) Give an efficient algorithm to find all of the triangles in a graph.

(b) How fast is your algorithm?

Problem 5. In a graph $G = (V, E)$ edge (x, y) *touches* vertices x and y.

A *newtonian cluster* in a graph $G = (V, E)$ is a subset of the edges such that every vertex is touched by at least one edge. The *size of a newtonian cluster* is the number of edges in the subset.

(a) Give an example of a graph that has a newtonian cluster of size four but not of size three.

(b) Let C be a newtonian cluster of G. What can you say about the minimum size of C as a function of the number of vertices n? Justify.

(c) A *minimal newtonian cluster* is a newtonian cluster such that if any edge is removed it is no longer a newtonian cluster. Let C be a minimal newtonian cluster of G. What can you say about the maximum size of C as a function of the number of vertices n? Justify.

(d) The (decision version of) *newtonian cluster problem* is given a graph $G = (V, E)$ and an integer k does G have a newtonian cluster of size (at most) k. Show that the newtonian cluster problem is in **NP**. What is the certificate?

Problem 6. A *vertex cover* in a graph $G = (V, E)$ is a subset of vertices such every edge is incident on at least one vertex of the subset. The *Weighted Vertex Cover Problem (WVCP)* is, given a graph $G = (V, E)$ with integer weights on the vertices, find a vertex cover whose sum of weights is as small as possible. You can assume that the weights are between 1 and n (inclusive).

(a) WVCP is an optimization problem. Define a decision version of WVCP.

(b) Show that the decision version is in **NP**. Make sure to state the certificate and give the pseudo code.

(c) Show that if you could solve the optimization version in polynomial time that you could also solve the decision version in polynomial time.

(d) Show that if you could solve the decision version in polynomial time that you could also solve the optimization version in polynomial time. HINT: First find the weight of an optimal weighted vertex cover.

Problem 7. This problem is more open-ended than you would see on an exam: If you do not know how to play Sudoku, look it up. Normally, Sudoku is played on a 9×9 grid.

(a) Generalize Sudoku to larger grids.

(b) State the (generalized) Sudoku game as a decision problem.

(c) Show that the decision version of (generalized) Sudoku is in NP.

(d) Show that if you can solve the decision version of (generalized) Sudoku in polynomial time, you can solve a (generalized) Sudoku puzzle in polynomial time.