

CMSC 330: Organization of Programming Languages

OCaml Higher Order Functions

Anonymous Functions

- ▶ Recall code blocks in Ruby

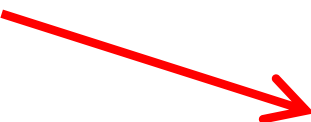
```
(1..10).each { |x| print x }
```

- Here, we can think of `{ |x| print x }` as a function

- ▶ We can do this (and more) in OCaml

Anonymous Functions

- ▶ As with Ruby, passing around functions is common
 - So often we don't want to bother to give them names
- ▶ Use `fun` to make a function with no name

Parameter  Body
(in which parameter x is bound)

```
fun x -> x + 3
```

```
(fun x -> x + 3) 5 = 8
```

Anonymous Functions

▶ Syntax

- **fun** x_1 ... x_n \rightarrow e

▶ Evaluation

- An anonymous function is an expression
- In fact, *it is a value* – no further evaluation is possible
 - As such, it can be passed to other functions, returned from them, stored in a variable, etc.

▶ Type checking

- $(\text{fun } x_1 \dots x_n \rightarrow e) : (t_1 \rightarrow \dots \rightarrow t_n \rightarrow u)$
when $e : u$ under assumptions $x_1 : t_1, \dots, x_n : t_n$.
 - (Same rule as `let f x1 ... xn = e`)

Calling Functions, Generalized

Not just a variable f

▶ Syntax $e_0 e_1 \dots e_n$

▶ Evaluation

- Evaluate arguments $e_1 \dots e_n$ to values $v_1 \dots v_n$
 - Order is actually right to left, not left to right
 - But this doesn't matter if $e_1 \dots e_n$ don't have side effects
- Evaluate e_0 to a function $\text{fun } x_1 \dots x_n \rightarrow e$
- Substitute v_i for x_i in e , yielding new expression e'
- Evaluate e' to value v , which is the final result

▶ Example:

• $(\text{fun } x \rightarrow x+x) 1 \Rightarrow 1+1 \Rightarrow 2$

Calling Functions, Generalized

- ▶ Syntax $e0\ e1\ \dots\ en$
- ▶ Type checking (almost the same as before)
 - If $e0 : t1 \rightarrow \dots \rightarrow tn \rightarrow u$ and $e1 : t1, \dots, en : tn$ then $e0\ e1\ \dots\ en : u$
- ▶ Example:
 - `(fun x -> x+x) 1 : int`
 - since `(fun x -> x+x) : int -> int` and `1 : int`

Quiz 1: What does this evaluate to?

```
let y = (fun x -> x+1) 2 in  
(fun z -> z-2) y
```

- A. *Error*
- B. 2
- C. 1
- D. 0

Quiz 1: What does this evaluate to?

```
let y = (fun x -> x+1) 2 in  
(fun z -> z-2) y
```

- A. *Error*
- B. 2
- C. 1**
- D. 0

Quiz 2: What is this expression's type ?

`(fun x y -> x) 2 3`

- A. *Type error*
- B. `int`
- C. `int -> int -> int`
- D. `'a -> 'b -> 'a`

Quiz 2: What is this expression's type ?

`(fun x y -> x) 2 3`

- A. *Type error*
- B. `int`**
- C. `int -> int -> int`
- D. `'a -> 'b -> 'a`

Functions and Binding

- ▶ Functions are **first-class**, so you can bind them to other names as you like

```
let f x = x + 3;;
```

```
let g = f;;
```

```
g 5      = 8
```

- ▶ In fact, **let** for functions is syntactic **shorthand**

```
let f x = body
```

↓ is semantically equivalent to

```
let f = fun x -> body
```

Example Shorthands

- ▶ `let next x = x + 1`
 - Short for `let next = fun x -> x + 1`
- ▶ `let plus x y = x + y`
 - Short for `let plus = fun x y -> x + y`
- ▶ `let rec fact n =
 if n = 0 then 1 else n * fact (n-1)`
 - Short for `let rec fact = fun n ->
 (if n = 0 then 1 else n * fact (n-1))`

Quiz 3: What does this evaluate to?

```
let f = fun x -> 0 in
let g = f in
g 1
```

- A. *Error*
- B. 2
- C. 1
- D. 0

Quiz 3: What does this evaluate to?

```
let f = fun x -> 0 in
let g = f in
g 1
```

A. *Error*

B. 2

C. 1

D. 0

Defining Functions Everywhere

```
let move l x =  
  let left x = x - 1 in (* locally defined fun *)  
  let right x = x + 1 in (* locally defined fun *)  
  if l then left x  
  else      right x  
;;
```

```
let move' l x = (* equivalent to the above *)  
  if l then (fun y -> y - 1) x  
  else      (fun y -> y + 1) x
```

Pattern Matching With Fun

- ▶ `match` can be used within `fun`

```
(fun l -> match l with (h::_) -> h) [1; 2]  
= 1
```

- ▶ But use named functions for complicated matches
- ▶ May use standard pattern matching abbreviations

```
(fun (x, y) -> x+y) (1, 2)  
= 3
```


Passing Functions as Arguments

- ▶ In OCaml you can pass functions as arguments (akin to Ruby code blocks)

```
let plus_three x = x + 3 (* int -> int *)
```

```
let twice f z = f (f z) (* ('a->'a) -> 'a -> 'a *)  
twice plus_three 5 = 11
```

- ▶ Ruby's `collect` is called `map` in OCaml
 - `map f l` applies function `f` to each element of `l`, and puts the results in a new list (preserving order)

```
map plus_three [1; 2; 3] = [4; 5; 6]
```

```
map (fun x -> (-x)) [1; 2; 3] = [-1; -2; -3]
```

map function

What is Map?

Map generates a new list by applying a function to every item in the given list

$\text{map } f [n1;n2;n3] == > [f \ n1; f \ n2; f \ n3]$

map cook [🐮, 🍷, 🐔, 🌽]
== > [🍔, 🍟, 🍗, 🍿]

Why do we need Map?

```
let rec double lst =  
  match lst with  
  []->[]  
  |h::t-> h * 2 :: double t
```

```
let rec neg lst =  
  match lst with  
  []->[]  
  |h::t-> h * (-1) :: neg t
```

```
double [1; 2; 3; 4];;  
- : int list = [2; 4; 6; 8]
```

```
neg [1;2;3;4];;  
- : int list = [-1; -2; -3; -4]
```

Why do we need Map?

```
let rec double lst =  
  match lst with  
  []->[]  
  |h::t-> h * 2 :: double t
```

```
let rec neg lst =  
  match lst with  
  []->[]  
  |h::t-> h * (-1) :: neg t
```

```
let rec map f lst =  
  match lst with  
  []->[]  
  |h::t-> (f h):: map f t
```

How to implement Map?

```
let rec map f lst =  
  match lst with  
  | [] -> []  
  | h :: t -> (f h) :: (map f t)
```

Type of Map

```
let map f lst =  
  match lst with  
  | [] -> []  
  | h::t -> (f h):: map f t
```

```
('a -> 'b) -> 'a list -> 'b list
```

How to use Map?

```
let double x = x * 2 ;;
```

```
let lst = [1; 2; 3; 4; 5] ;;
```

```
let t = map double lst ;;
```

```
t : int list = [2; 4; 6; 8; 10]
```


Example 1

Subtract 1 from every item in an int list

```
let t = [1; 2; 3; 4];;  
map (fun x-> x-1) t;;
```

```
let t = [1; 2; 3; 4];;  
let sub1 x = x - 1;;  
map sub1 t;;
```

```
int list = [0; 1; 2; 3]
```

Example 2

Negate every item in an int list

```
let t = [1; 2; 3; 4];;  
let neg x = x * (-1);;  
map neg t;
```

```
int list = [-1; -2; -3; -4]
```

Example 3

Apply a list functions to an int list

```
let lst = [1;2;3];;  
let neg x = x * (-1);;  
let sub1 x = x-1;;  
let double x = x + x;;
```

```
let fs = [neg; sub1; double];;  
map (fun x -> map x lst) fs;;
```

```
int list list = [[-1; -2; -3]; [0; 1; 2]; [2; 4; 6]]
```

Example 4: Permute a list

```
let permute lst =
  let rec rm x l = List.filter ((<>) x) l
  and insertToPermute lst x =
    let t = rm x lst in
    List.map ((fun a b->a::b) x )(permuteall t)
  and permuteall lst =
    match lst with
    | []->[]
    | [x]->[[x]]
    | _->List.flatten(List.map (insertToPermute lst) lst)
  in permuteall lst
;;

# permute [1;2;3];;
- : int list list =
[[1; 2; 3]; [1; 3; 2]; [2; 1; 3]; [2; 3; 1]; [3; 1; 2];
[3; 2; 1]]
```

Example 5: Power Set

```
let populate a b =
  if b=[] then [[a]]
  else let t = List.map (fun x->a::x) b in
        [a]::t@b
;;

let powerset lst = List.fold_right populate lst []
;;

# powerset [1;2;3];;
- : int list list = [[1]; [1; 2]; [1; 2; 3]; [1; 3];
[2]; [2; 3]; [3]]

# populate 1 [[2];[3]];
- : int list list =
[[1]; [1; 2]; [1; 3]; [2];
[3]]
```

What we learned?

▶ Map:

- A higher order function.
- List module
- Takes a function and a list as arguments, applies the function to each member of the list, generates a new list
- It is powerful.

fold function

What is Fold

- Fold generally
 - takes a **function of two arguments**, a **list**, and an **initial value** (accumulator)
 - **combines the list** by applying the function to the accumulator and one element from the list and the result of recursively folding the function over the rest of the list.

Accumulator: (i.e. 0 for addition, 1 for multiplication, false for boolean OR, negative infinity for maximum, etc.)

What is Fold

```
fold (fun x y-> x+y) 0 [1;2;3;4;5];;
```

```
- : int = 15
```

Why do we need Fold?

sum a list of integers

```
let rec sum l =  
  match l with  
  [] -> 0  
  |h::t -> h + (sum t)
```

```
sum [1;2;3;4];;  
- : int = 10
```

Concatenate a list of strings:

```
let rec concat l =  
  match l with  
  [] -> ""  
  |h::t -> h ^ (concat t)
```

```
concat ["a";"b";"c"];;  
- : string = "abc"
```

Why do we need Fold?

sum a list of integers

```
let rec sum l =  
  match l with  
  [] -> 0  
  |h::t -> h + (sum t)
```

Concatenate a list of strings:

```
let rec concat l =  
  match l with  
  [] -> ""  
  |h::t -> h ^ (concat t)
```

```
let rec fold f acc lst =  
  match l with  
  [] -> acc  
  |h::t -> fold f (f acc h) t
```

How to implement Fold

```
let rec fold f acc lst =  
  match l with  
  [] -> acc  
  |h::t -> fold f (f acc h) t
```

Type of Fold

```
let rec fold f acc lst =  
  match l with  
  [] -> acc  
  |h::t -> fold f (f acc h) t
```

f acc lst -> return type

('a -> 'b -> 'a) -> 'a -> 'b list -> 'a

How to use Fold?

```
let add x y = x + y ;;
```

```
let lst = [2; 3; 4] ;;
```

```
let t = fold add 0 lst ;;
```

```
t : int = 9
```

How to use Fold?

```
let add x y = x + y ;;
let lst = [2; 3; 4] ;;
let t = fold add 0 lst ;;
t : int = 9
```

```
let rec fold f acc lst =
  match lst with
  [] -> acc
  |h :: t -> fold f (f acc h) t
```

```
fold add 0 lst
fold add (add 0 2) [3;4]
fold add 2 [3;4]
fold add (add 2 3) [4]
fold add 5 [4]
fold add (add 5 4) [ ]
fold add 9 [ ]
9
```

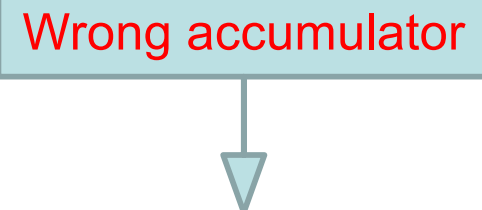
Example 1: Product of an int list

```
let mul x y = x * y;;
```

```
let lst = [1; 2; 3; 4; 5];;
```

```
fold mul 1 lst  
- : int = 120
```

Wrong accumulator



```
fold mul 0 lst;;  
- : int = 0
```


Example 2: Count elements of a list satisfying a condition

```
let countif p l =  
  fold (fun counter element -> if p element then counter+1  
                                else counter) 0 l ;;
```

```
countif (fun x -> x > 0) [30;-1;45;100;0];;
```

```
- : int = 3
```

Exaple 3: Collect even numbers in the list

```
let f acc y = if (y mod 2) = 0 then y::acc  
              else acc;;
```

```
fold f [] [1;2;3;4;5;6];;
```

```
- : int list = [6; 4; 2] ← Reversed
```

Example 4: Inner Product

first compute list of pair-wise products, then sum up

$$[x1;x2;x3]*[y1;y2;y3] = [x1*y1 + x2*y2 + x3*y3]$$

```
let rec map2 f a b =  
  match (a,b) with  
  |([],[])->([])  
  |(h1::t1,h2::t2)->(f h1 h2)::(map2 f t1 t2)  
  |_->invalid_arg "map2";;
```

```
let product v1 v2 =  
  fold (+) 0 (map2 ( * ) v1 v2);;  
# val product : int list -> int list -> int = <fun>  
product [2;4;6] [1;3;5];;  
#- : int = 44
```

Example 5: Find the maximum from a list

```
let maxList lst =  
  match lst with  
  []->failwith "empty list"  
 |h::t-> fold max h t ;;
```

```
maxList [3;10;5];;  
- : int = 10
```

```
(*  
maxList [3;10;5]  
fold max 3 [10;5]  
fold max (max 3 10) [5]  
fold max (max 10 5) []  
fold max 10 []  
10 *)
```

Quiz: Sum of sublists

Given a list of int lists, compute the sum of each int list, and return them as list.

For example:

```
sumList [[1;2;3];[4];[5;6;7]]  
- : int list = [6; 4; 18]
```

Solution: Sum of sublists

```
let sumList = map (fold (+) 0 );;
```

```
sumList [[1;2;3];[4;5;6];[10]];;
```

```
- : int list = [6; 15; 10]
```

Quiz: Maximum contiguous subarray

Given an int list, find the contiguous sublist, which has the largest sum and return its sum.

Example:

Input: [-2, 1, -3, 4, -1, 2, 1, -5, 4]

Output: 6

Explanation: [4, -1, 2, 1] has the largest sum = 6

Quiz: Maximum contiguous subarray

```
let f (m, acc) h =
  let m = max m (acc + h) in
  let x = if acc < 0 then 0 else acc in
  (m, x+h)
;;
let submax lst = let (max_so_far, max_current) =
  fold f (0,0) lst in
  max_so_far
;;

submax [-2; 1; -3; 4; -1; 2; 1; -5; 4];;
- : int = 6
```


Summary

- ▶ $\text{map } f \ [v1; v2; \dots; vn]$
= $[f \ v1; f \ v2; \dots; f \ vn]$
 - e.g., $\text{map } (\text{fun } x \rightarrow x+1) \ [1;2;3] = [2;3;4]$
- ▶ $\text{fold } f \quad v \quad [v1; v2; \dots; vn]$
= $\text{fold } f \quad (f \ v \ v1) \quad [v2; \dots; vn]$
= $\text{fold } f \ (f \ (f \ v \ v1) \ v2) \quad [\dots; vn]$
= ...
= $f \ (f \ (f \ (f \ v \ v1) \ v2) \ \dots) \ vn$
 - e.g., $\text{fold } \text{add } 0 \ [1;2;3;4] =$
 $\text{add } (\text{add } (\text{add } (\text{add } 0 \ 1) \ 2) \ 3) \ 4 = 10$

Quiz 4: What does this evaluate to?

```
map (fun x -> x *. 4) [1;2;3]
```

- A. [1.0; 2.0; 3.0]
- B. [4.0; 8.0; 12.0]
- C. Error
- D. [4; 8; 12]

Quiz 4: What does this evaluate to?

```
map (fun x -> x *. 4) [1;2;3]
```

A. [1.0; 2.0; 3.0]

B. [4.0; 8.0; 12.0]

C. Error -- the *. function takes floats, not ints

D. [4; 8; 12]

Quiz 5: What does this evaluate to?

```
fold (fun a y -> y::a) [] [3;4;2]
```

- A. [9]
- B. [3;4;2]
- C. [2;4;3]
- D. Error

Quiz 5: What does this evaluate to?

```
fold (fun a y -> y::a) [] [3;4;2]
```

- A. [9]
- B. [3;4;2]
- C. [2;4;3]
- D. Error

Quiz 6: What does this evaluate to?

```
let is_even x = (x mod 2 = 0) in  
map is_even [1;2;3;4;5]
```

- A. [false;true;false>true;false]
- B. [0;1;1;2;2]
- C. [0;0;0;0;0]
- D. false

Quiz 6: What does this evaluate to?

```
let is_even x = (x mod 2 = 0) in  
map is_even [1;2;3;4;5]
```

- A. **[false;true;false>true;false]**
- B. [0;1;1;2;2]
- C. [0;0;0;0;0]
- D. false

Combining map and fold

- ▶ Idea: map a list to another list, and then fold over it to compute the final result
 - Basis of the famous “map/reduce” framework from Google, since these operations can be parallelized

```
let countone l =  
  fold (fun a h -> if h=1 then a+1 else a) 0 l
```

```
let countones ss =  
  let counts = map countone ss in  
  fold (fun a c -> a+c) 0 counts
```

```
countones [[1;0;1]; [0;0]; [1;1]] = 4
```

```
countones [[1;0]; []; [0;0]; [1]] = 2
```


fold_right

- ▶ Right-to-left version of fold:

```
let rec fold_right f l a = match l with
  [] -> a
  | (h::t) -> f h (fold_right f t a)
```

- ▶ Left-to-right version used so far:

```
let rec fold f a l = match l with
  [] -> a
  | (h::t) -> fold f (f a h) t
```

Left-to-right vs. right-to-left

`fold f v [v1; v2; ...; vn] =`
 `f (f (f (f v v1) v2) ...) vn`

`fold_right f [v1; v2; ...; vn] v =`
 `f (f (f (f vn v) ...) v2) v1`

`fold (fun x y -> x - y) 0 [1;2;3] = -6`

since $((0-1)-2)-3 = -6$

`fold_right (fun x y -> x - y) [1;2;3] 0 = 2`

since $1-(2-(3-0)) = 2$

When to use one or the other?

- ▶ Many problems lend themselves to `fold_right`
- ▶ But it does present a performance disadvantage
 - The recursion builds of a deep stack: **One stack frame for each recursive call of `fold_right`**
- ▶ An optimization called **tail recursion** permits optimizing `fold` so that it **uses no stack at all**
 - We will see how this works in a later lecture!