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# Distance Vector Routing

- Distributed Algorithm
  - Routing Protocol
  - Forwarding
- 

what are we trying  
to compute?

# Model

- nodes have unique addresses
- can send out incident links
- can detect neighbors
- links are not reliable

Notation

at node  $x$

$$D^x(y, z)$$

Distance known at  $x$   
to dest.  $y$  via  
neighbor  $z$ .

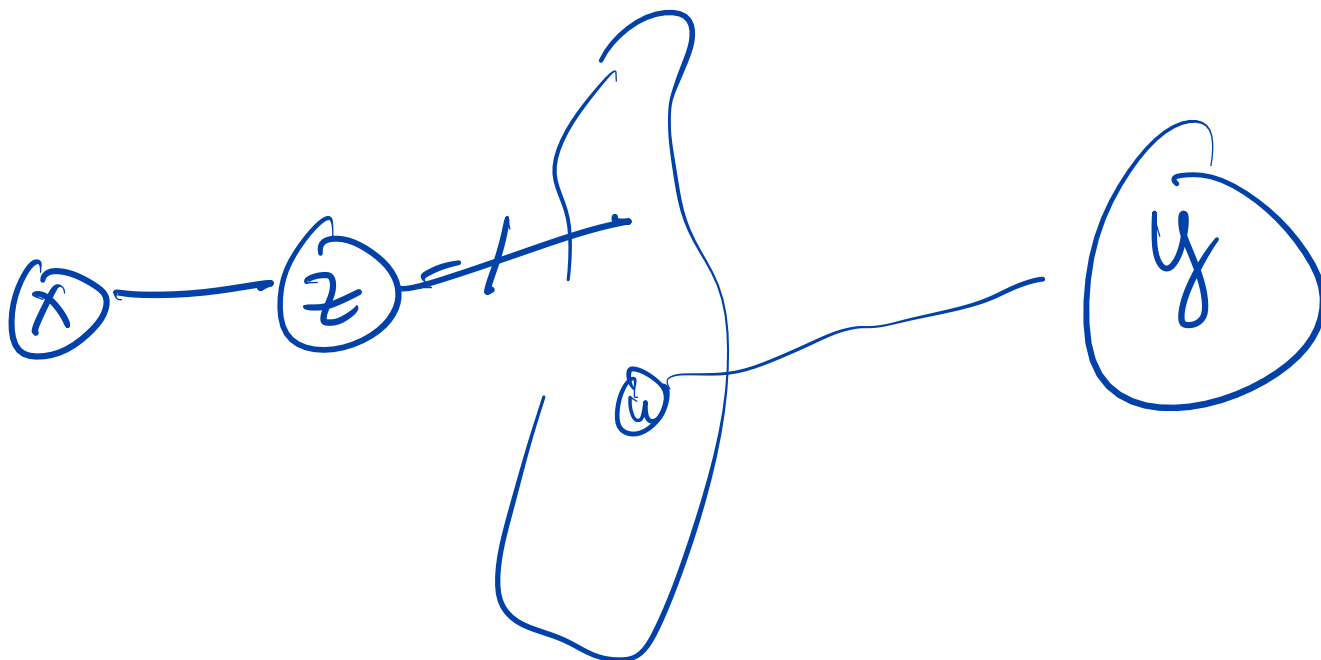
$\text{Cost}(x, y)$ : link cost  
between  $x, y$   
(non-negative)

# Convergence

$$D^x(y, z) \leftarrow \text{cost}(x, z)$$

$$+ \min_w D^z(y, w)$$

neighbors of  
 $z$



DV

$$G = (N, E)$$

initially

$$D^x(*, v) \leftarrow \infty$$

$$D^x(v, v) \leftarrow \text{cost}(x, v)$$

// direct  
links

$v$  : n'bor

$\forall$  dest  $y$   
send min  
 $w \in N$

$$D^x(y, w)$$

to

each neighbor

forever

get rt. update from  $v$   
- can be new link cost  
or update for some dest

if  $\text{cost}(x, v)$  changed by  $\delta$

$\forall$  dest  $y$

0.  $D^x(y, v) \leftarrow D^x(y, v) + \delta$

else if cost to  $y$  via  $v$  changed

1.  $D^x(y, v) \leftarrow \text{cost}(x, v) + \underline{\text{new cost}}$

forever

get rt. update from  $v$   
- can be new link cost  
or update for some dest

if  $\text{cost}(x, v)$  changed by  $\delta$   
& dest  $y$

update  $D^x(y, v) \leftarrow D^x(y, v) + \delta$

else if cost to  $y$  via  $v$  changed

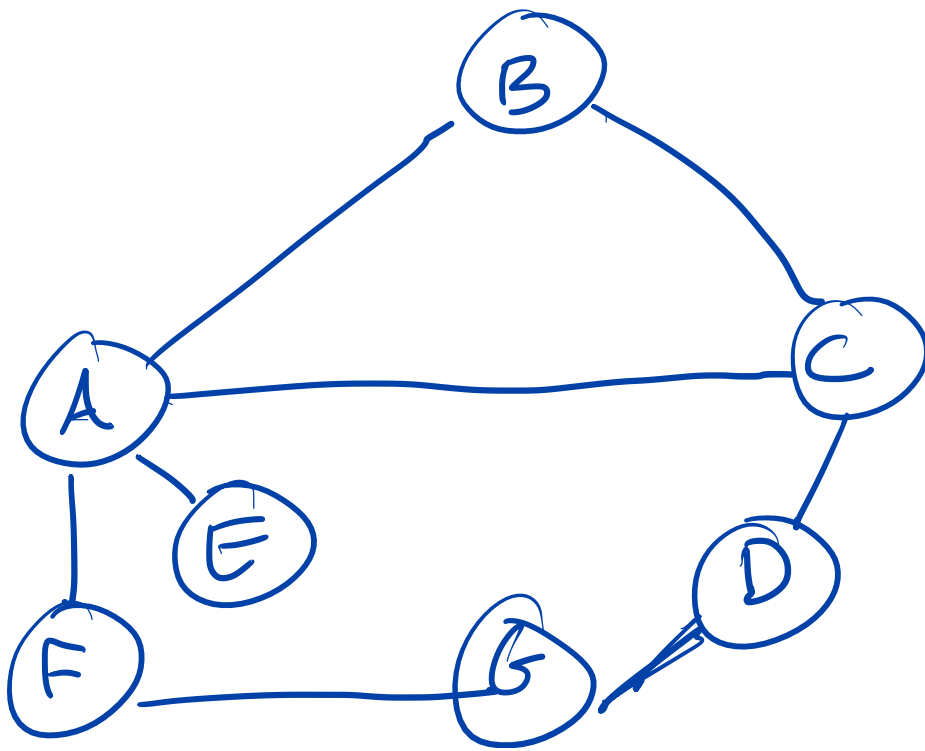
$D^x(y, v) \leftarrow \text{cost}(x, v)$   
+ new cost

propagate

if  $\min_w D^x(y, w)$  changed

Send  $D^x(y, w)$  to  
all neighbors.

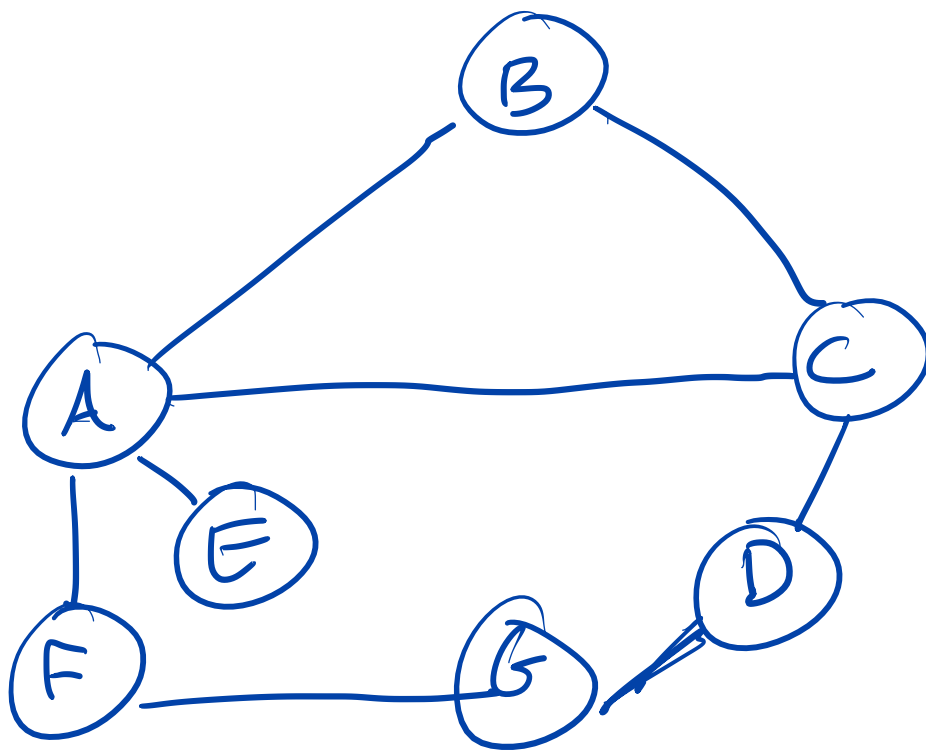




all links  
unit  
cost

at A

Dest	NH	Cost
B	B	1
C	C	1
D	-	$\infty$
E	E	1
F	F	1
G	-	$\infty$



all links  
unit  
cost

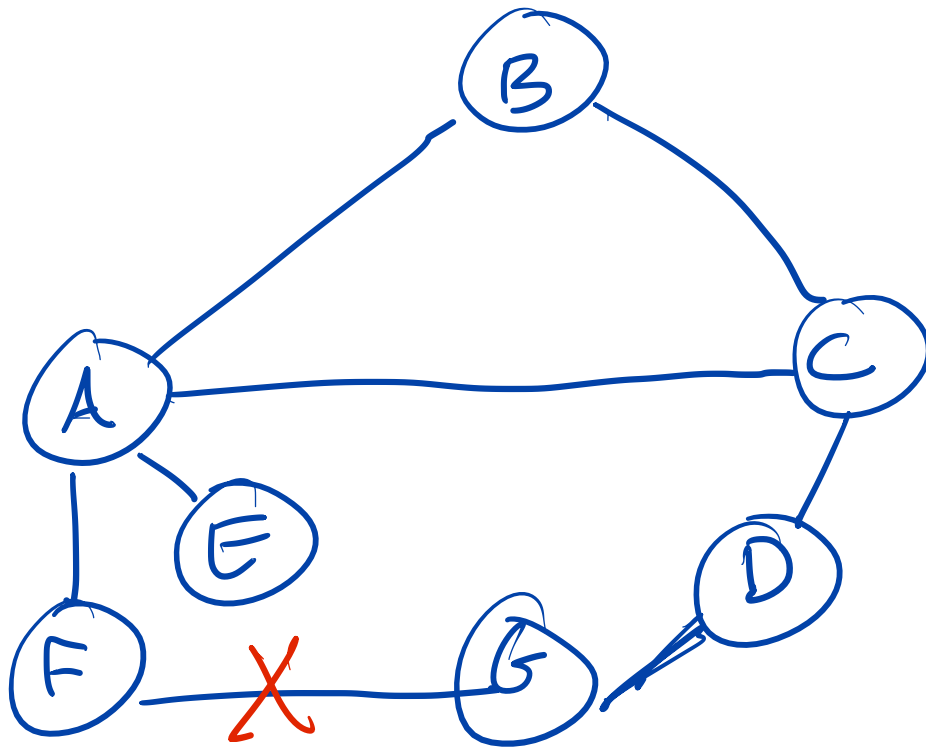
at A

Dest	NH	Cost
B	B	1
C	C	1
D	<u>C</u>	<del>2</del> 1
E	E	1
F	F	1
G	<u>F</u>	<del>2</del> 2

From F  
(F → G), 1

at A \*

From C  
C → D, 1



$F \rightarrow G$  link fails

at  $F$ ,  $(F \rightarrow G) \rightarrow \infty$

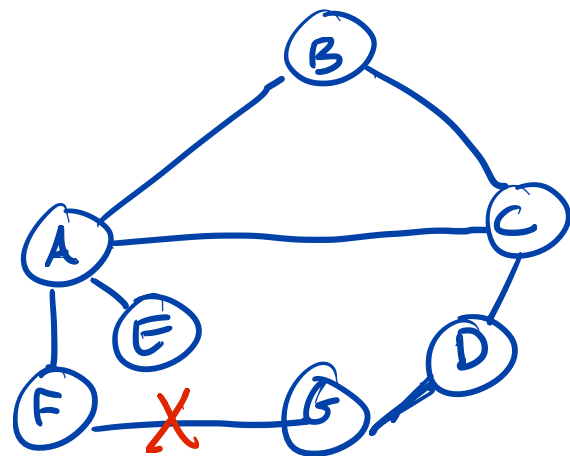
update

$F \rightarrow A : F \rightarrow G, \infty$

at  $A : A \rightarrow G, \underline{\underline{\infty}}$  ?

periodic updates  
vs. Full table

periodic update



later

$C \rightarrow A : C \rightarrow G, 2$

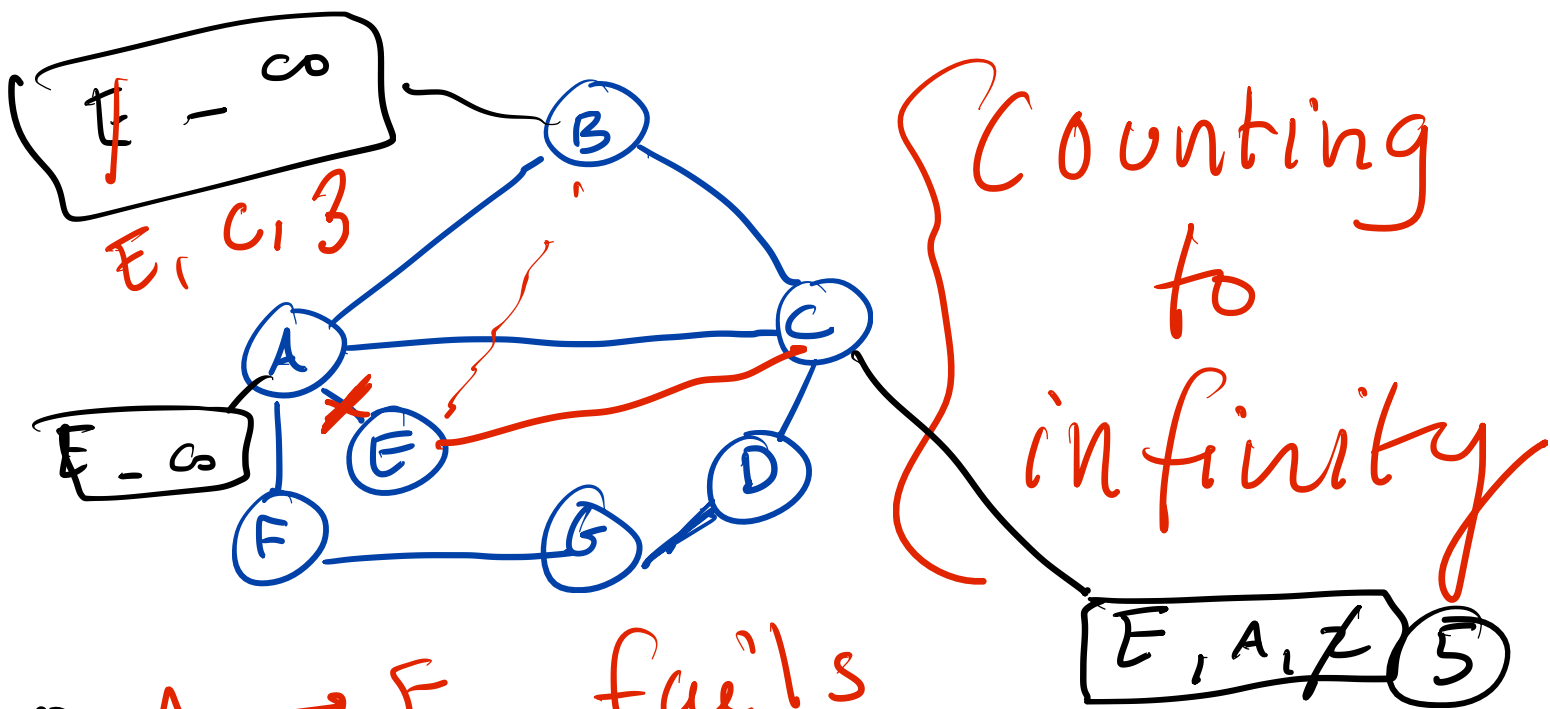
at A :  $A \rightarrow G, c, 3$

update to F

$A \rightarrow F : A \rightarrow G, 3$

at F

$F \rightarrow G : A, 4$



- 0  $A \rightarrow E$  fails
- 1 at A :  $(A \rightarrow E) \rightarrow \infty$  dropped
- 2 sends updates to B, ~~C~~ ...
- 3 B hears for C :  
 $C \rightarrow E, 2$
- 4 at B,  $B \rightarrow E$  :  $C, 3$  sent to A
- 5 at A,  $A \rightarrow E$  :  $B, 4$  sent to C
- 6 at C,  $C \rightarrow E$  :  $A, 5$ , sent to B
- 7 at B,  $B \rightarrow E$  :  $C, 6$ , sent to A
- ...

at A

Dest	NH	Cost
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B		
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→ C

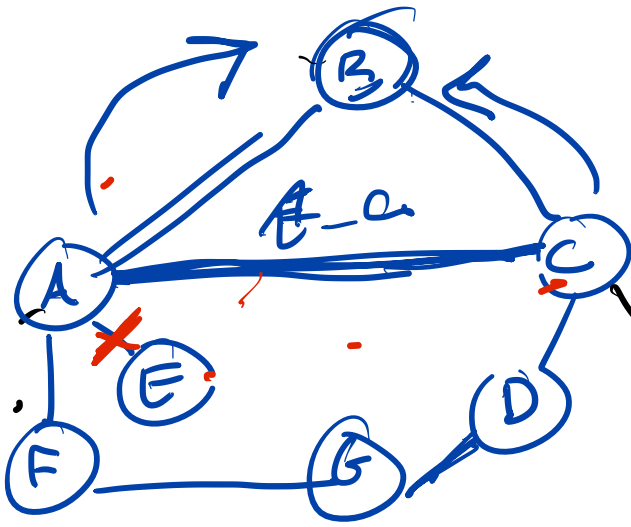
⋮

Z		
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Full Table



~~E, C, 3~~

Counting  
to  
infinity

①  $A \rightarrow E$  fails

at  $A : (A \rightarrow E) \rightarrow \infty$

2 sends updates to B, ~~C~~ ...

3 b hearc for C  
C → E, 2

4 at B,  $B \rightarrow E$ : 2, 3 sent to

at A,  $A \rightarrow E$  : B, 4 sent to

at  $C \rightarrow E: A, 5$ , send to

6. At  $C$ ,  $C \Rightarrow E$ .  $A, C, E$  sent to

at B,  $B \rightarrow E$ ; C, G, sent to

7 at

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