

CMSC 430 - 8 Sept 2026

Our first compilers

- Quick Racket and a86 review
- Abscond - integer literals
- Blackmail - unary primitives
- Con - conditional expressions
- Dupe - booleans, revised conditional

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Announcements

- Assignment 1 due Thurs, midnight
- New quiz released today, due by start of class Thursday
- TA Office Hours published on course web page

Abseond

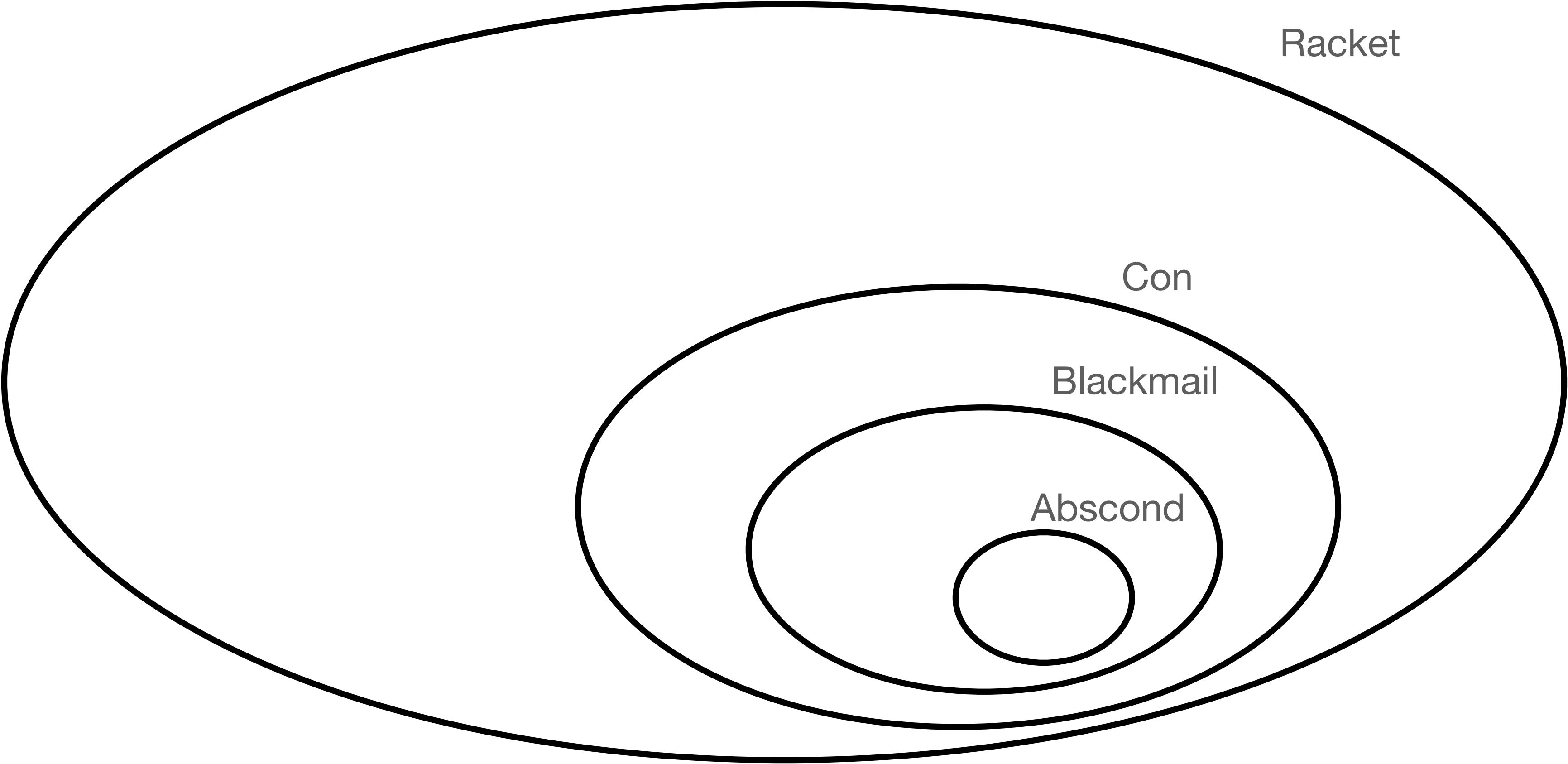
Learning objectives

Abscond

You should be able to answer:

- What is the structure of the Abscond compiler?
- How are Abscond programs represented abstractly?
- What invariant relates source programs and compiled code?

Language subsets



Example Abscond program

Concrete syntax

```
#lang racket
```

```
42
```


Example Abscond program

Abstract syntax

(Lit 42)

```
:: type Expr =  
;; | (Lit Integer)  
(struct Lit (n))
```

Example Abscond program

Parser constructs AST from string input

#lang racket
42

parse

→ (Lit 42)

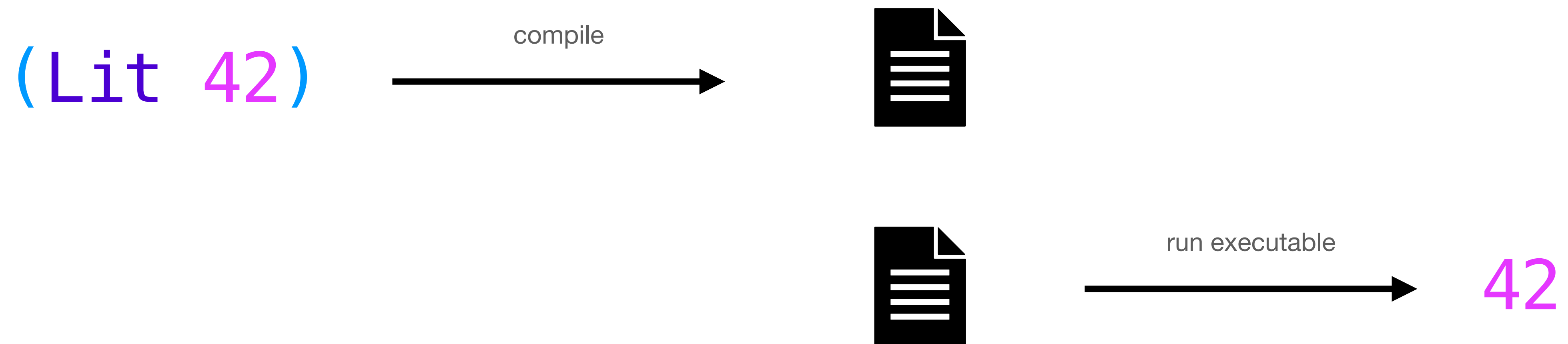
Example Abscond program

Interpreter computes meaning from AST



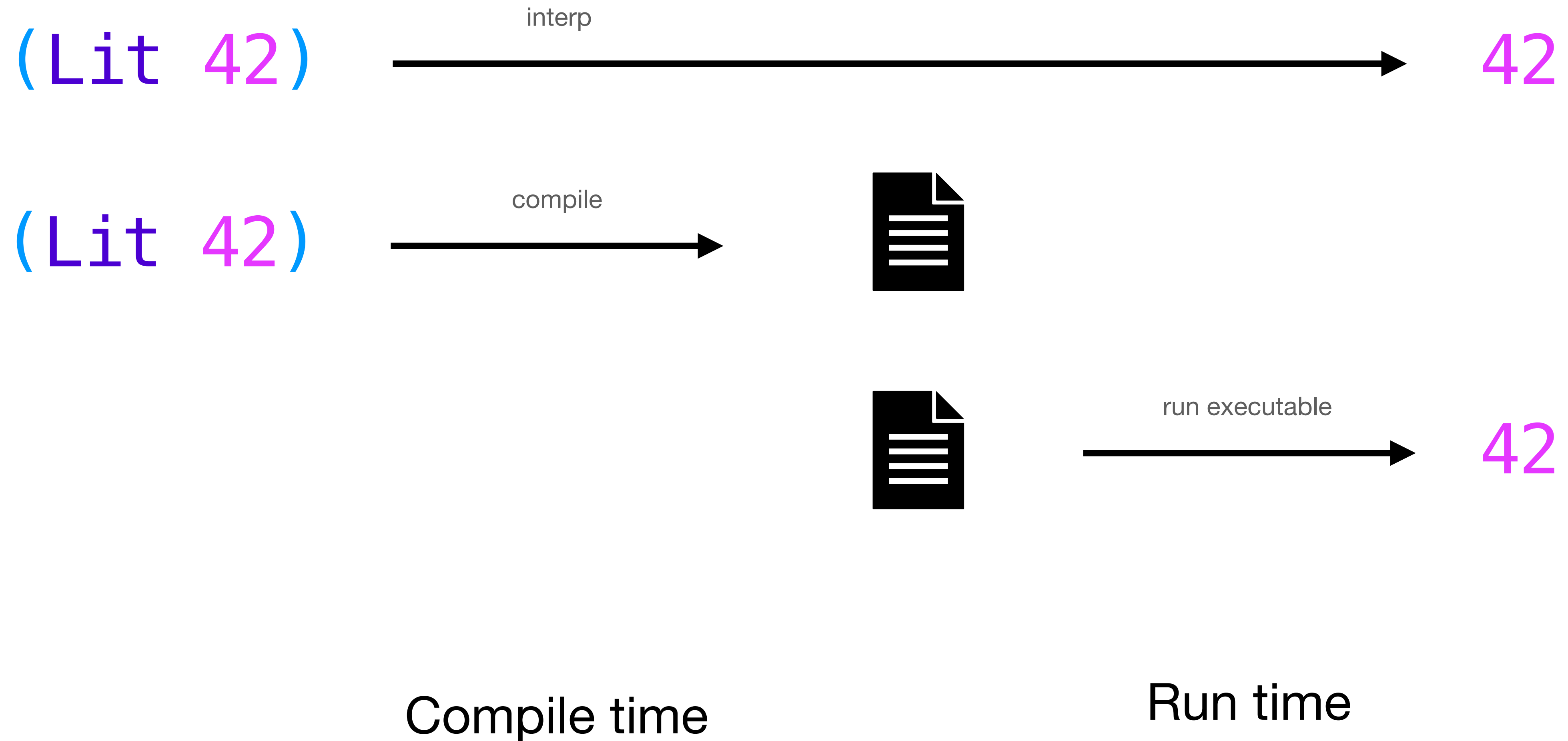
Example Abscond program

Compiler computes a program that computes meaning



Example Abscond program

Compiler specification



Example Abscond program

Our approach to parsing: reader + parser

#lang racket
42

read



42

parse



(Lit 42)

String

S-Expr

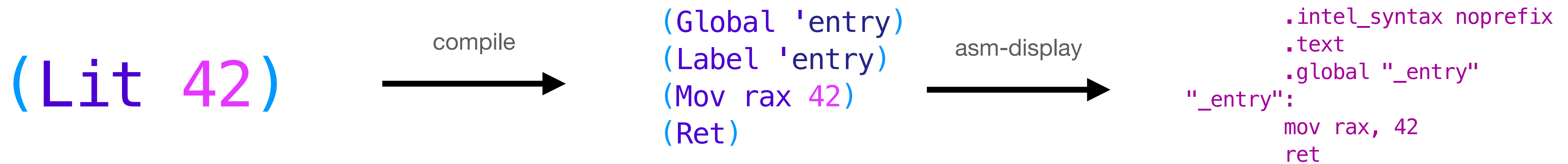
Expr

Semi-structured data;
ensures parens balanced, etc.

Structured data; ensures
grammatically well-formed

Example Abscond program

Our approach to compiling: compiler + stand-alone runtime



Expr

Construct AST of assembly
program

Asm

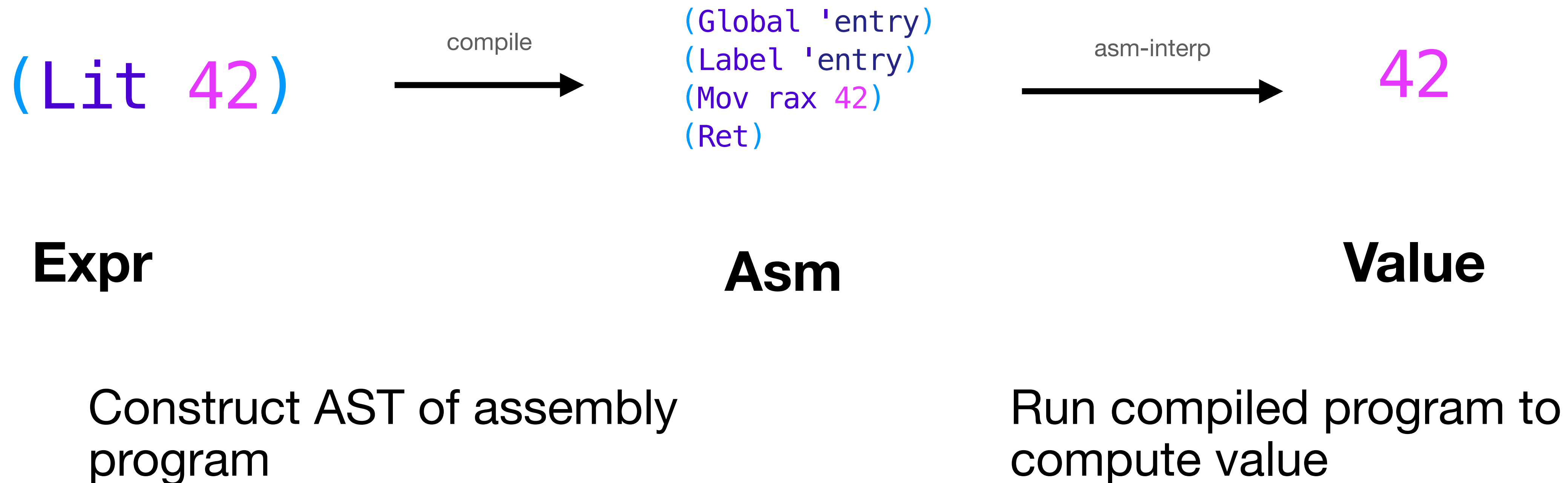
.S

Print to standard assembly
format; call assembler

Link w/ stand-alone runtime
system

Example Abscond program

Our approach to testing: compiler + asm-interp



Abscond

At the heart of it

```
;; Expr -> Asm  
(define (compile e) ...)  
;; Asm -> Value  
(define (run a) ...)
```

```
;; Expr -> Value  
(define (interp e) ...)
```

```
;; compile ◦ run ≡ interp
```

Abscond interpreter

At the heart of it

```
:: Expr -> Integer  
(define (interp e)  
  (match e  
    [(Lit i) i]))
```


Abscond compiler

At the heart of it

```
;; Expr -> Asm  
(define (compile e)  
  (list (Global 'entry)  
        (Label 'entry)  
        (match e  
          [(Lit i) (Mov rax i)]  
          (Ret))))
```

Abscond correctness

At the heart of it

```
;; Expr -> Void  
(define (check-compiler e)  
  (check-equal? (interp e)  
                (asm-interp (compile e))))
```

Hosted vs Standalone Runtime System

The runtime system is code that is needed at runtime to assist execution of compiled code. We will use two different runtime systems:

- Hosted: written in Racket w/ asm-interp; requires Racket at runtime
- Stand-alone: written in C; doesn't require Racket at runtime

Mostly we used hosted for development and testing, but the standalone is useful for when you want to make standalone executables.

Abscond directories

- `syntax/` - AST definition, parser
- `interpreter/` - interpreter, command line interpreter
- `compiler/` - compiler, command line compiler
- `executor/` - hosted runtime, command line executor
- `runtime/` - standalone runtime
- `test/` - unit tests

Blackmail

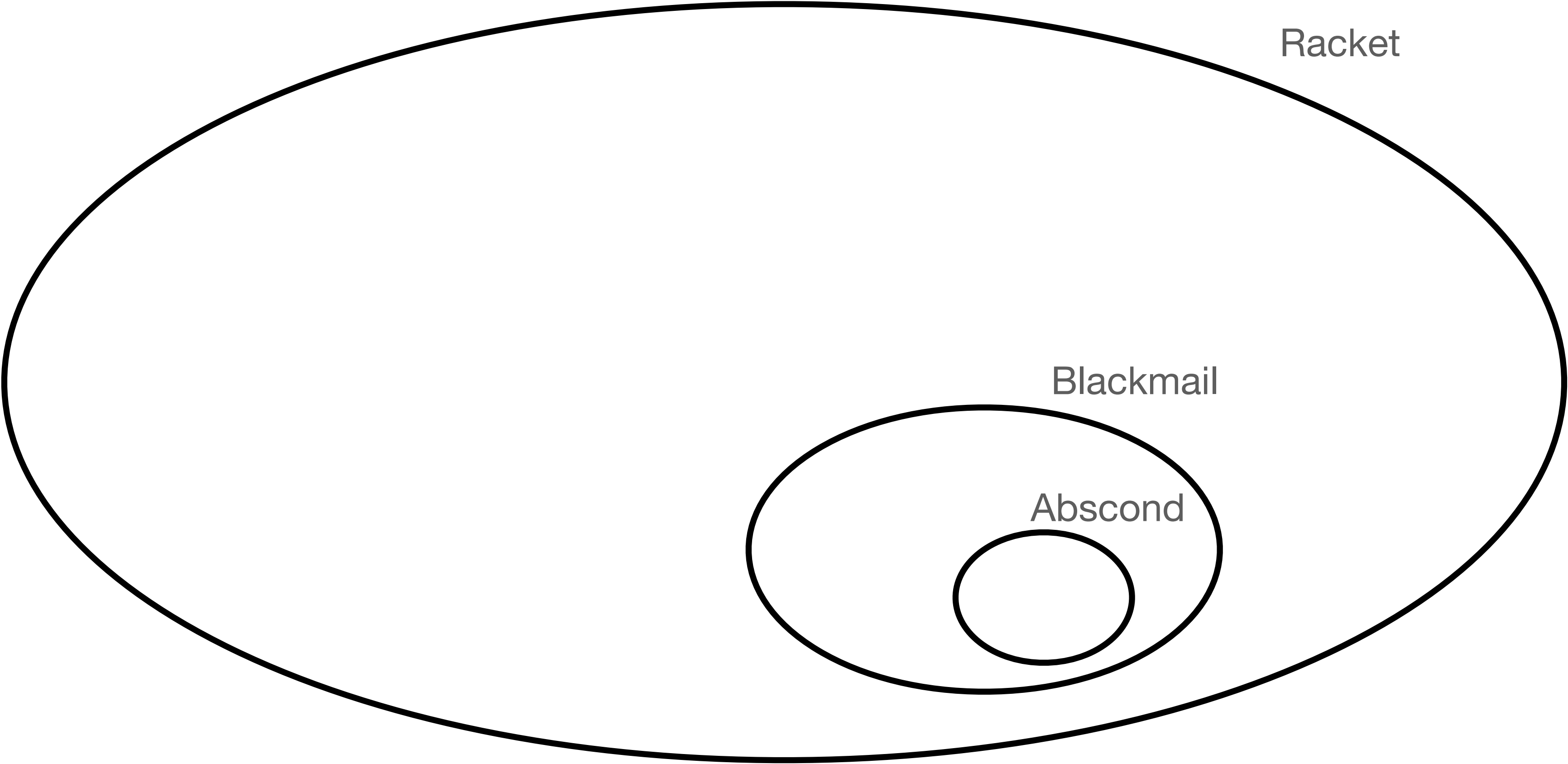
Learning objectives

Blackmail

You should be able to answer:

- How are Blackmail programs represented abstractly?
- What's the process for systematically growing our language?
- How do we structure an interpreter and compiler for an inductively defined AST data type?
- How do we both use and ensure the compiler invariants?

Language subsets



Recipe for growing a language

- Write examples
- Extend concrete syntax
- Extend abstract syntax
- Extend parser
- Revise interpreter to specify semantics
- Revise compiler & run-time system to implement semantics
- Test against examples
- *Rinse and repeat*

Example Blackmail program

Concrete syntax

```
#lang racket  
(add1 (add1 40))
```

Example Blackmail program

Abstract syntax

```
(Prim1 'add1 (Prim1 'add1 (Lit 40)))
```

```
:: type Expr = (Lit Integer)
::           | (Prim1 Op1 Expr)
:: type Op1 = 'add1 | 'sub1
(struct Lit (i))
(struct Prim1 (p e))
```

Example Blackmail program

Interpreter computes meaning from AST

`(Prim1 'add1 (Prim1 'add1 (Lit 40)))` $\xrightarrow{\text{interp}}$ 42

Example Blackmail program

Interpreter computes meaning from AST

```
(Prim1 'add1 (Prim1 'add1 (Lit 40)))
```

compile



```
(Global 'entry)
(Label 'entry)
(Mov rax 40)
(Add rax 1)
(Add rax 1)
(Ret)
```

asm-interp



42

Expr is an inductive data type

Abstract syntax

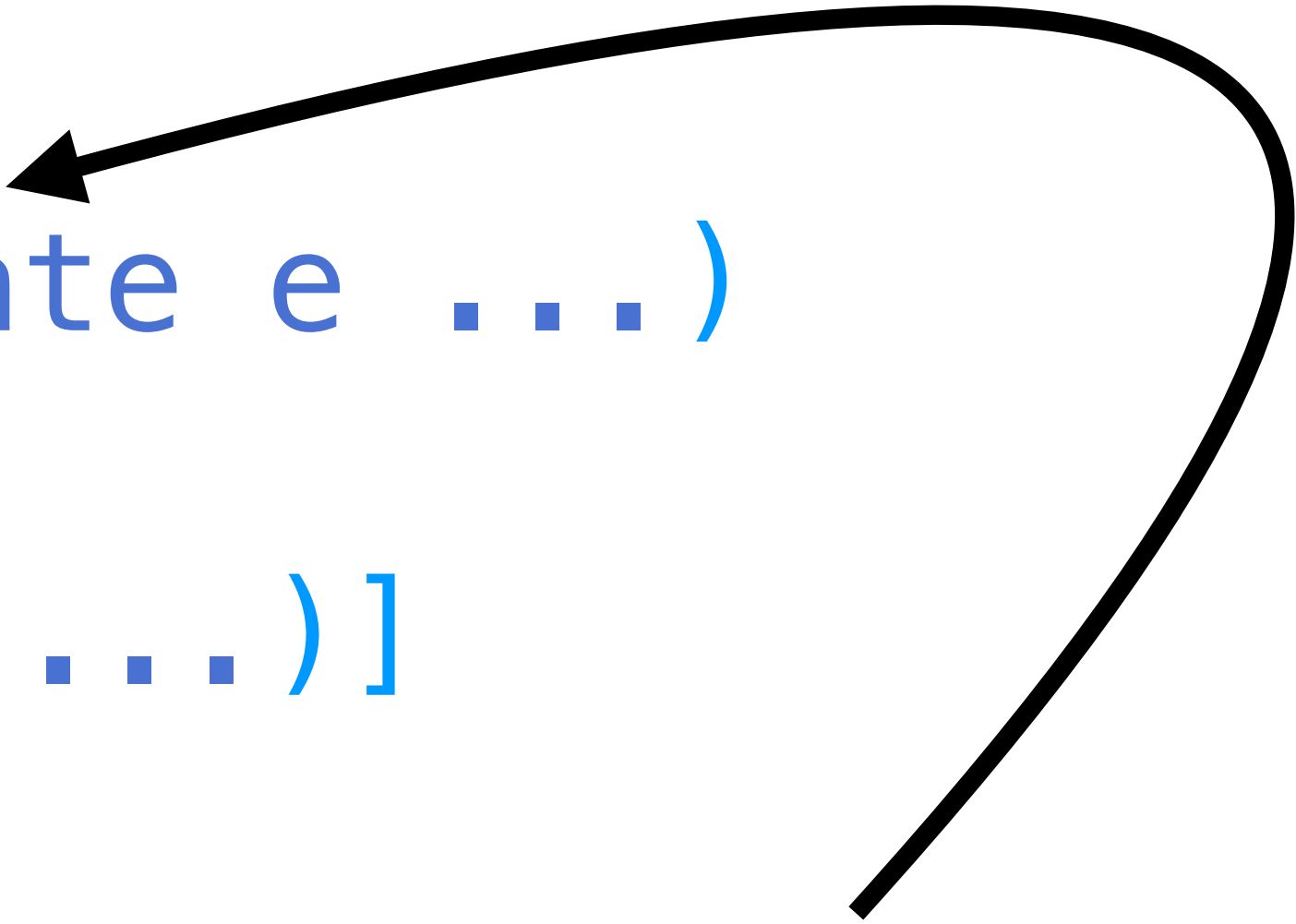
```
;; type Expr = (Lit Integer)
;;           | (Prim1 Op1 Expr)
;; type Op1 = 'add1 | 'sub1
(struct Lit (i))
(struct Prim1 (p e))
```



Expr is an inductive data type

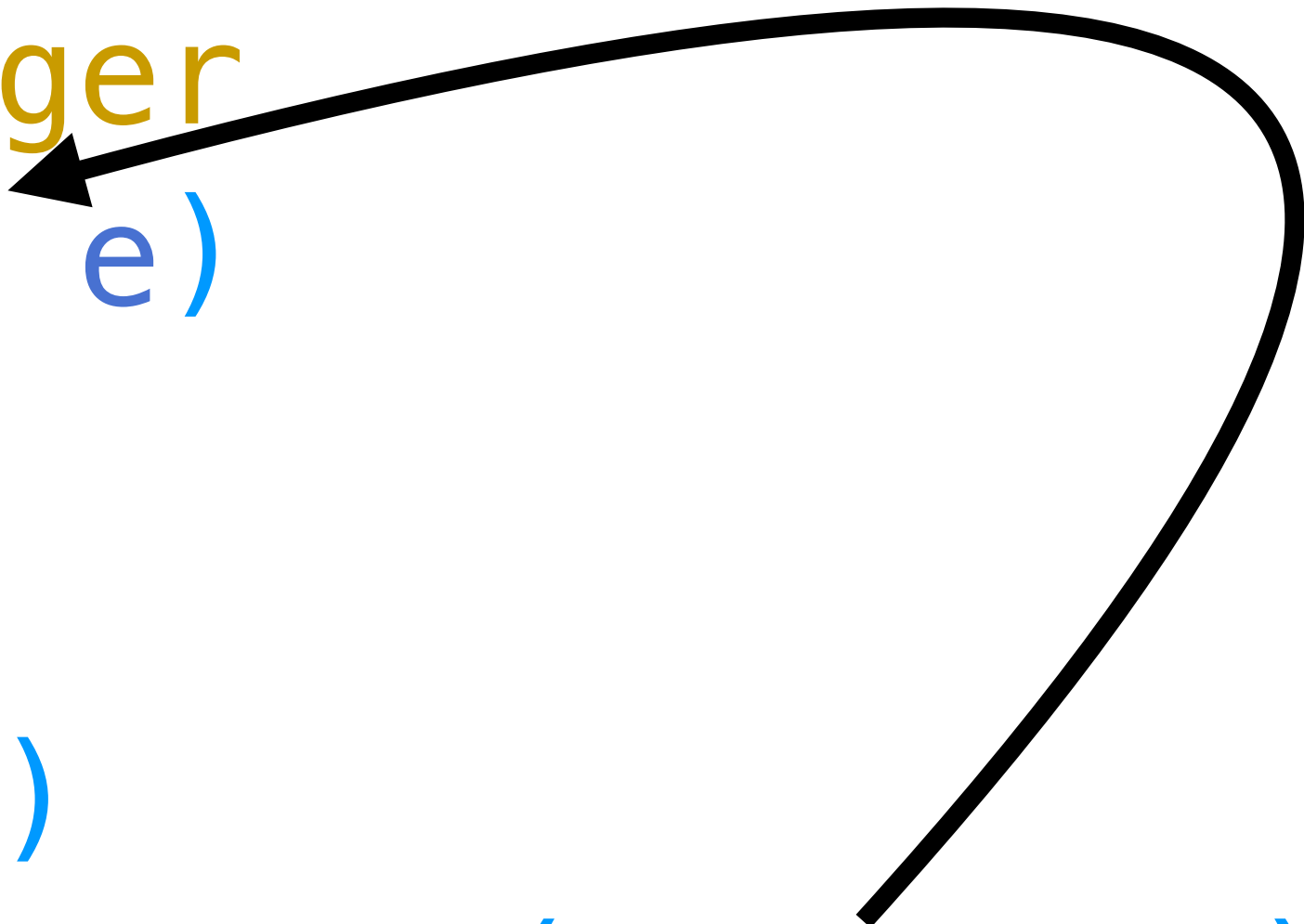
Abstract syntax

```
(define (expr-template e ...)  
  (match e  
    [(Lit i) (... i ...)]  
    [(Prim1 o e)  
     (... o ... (expr-template e ...) ...)]))
```



Blackmail interpreter

```
;; Expr -> Integer  
(define (interp e)  
  (match e  
    [(Lit i) i]  
    [(Prim1 p e)  
     (interp-prim1 p (interp e))]))
```



Blackmail interpreter

```
:: Op1 Integer -> Integer  
(define (interp-prim1 op i)  
  (match op  
    [ 'add1 (add1 i)]  
    [ 'sub1 (sub1 i)]))
```


Blackmail compiler

Compiling programs

```
;; Expr -> Asm  
(define (compile e)  
  (prog (Global 'entry)  
        (Label 'entry)  
        (compile-e e)  
        (Ret)))
```

Blackmail compiler

Key invariant

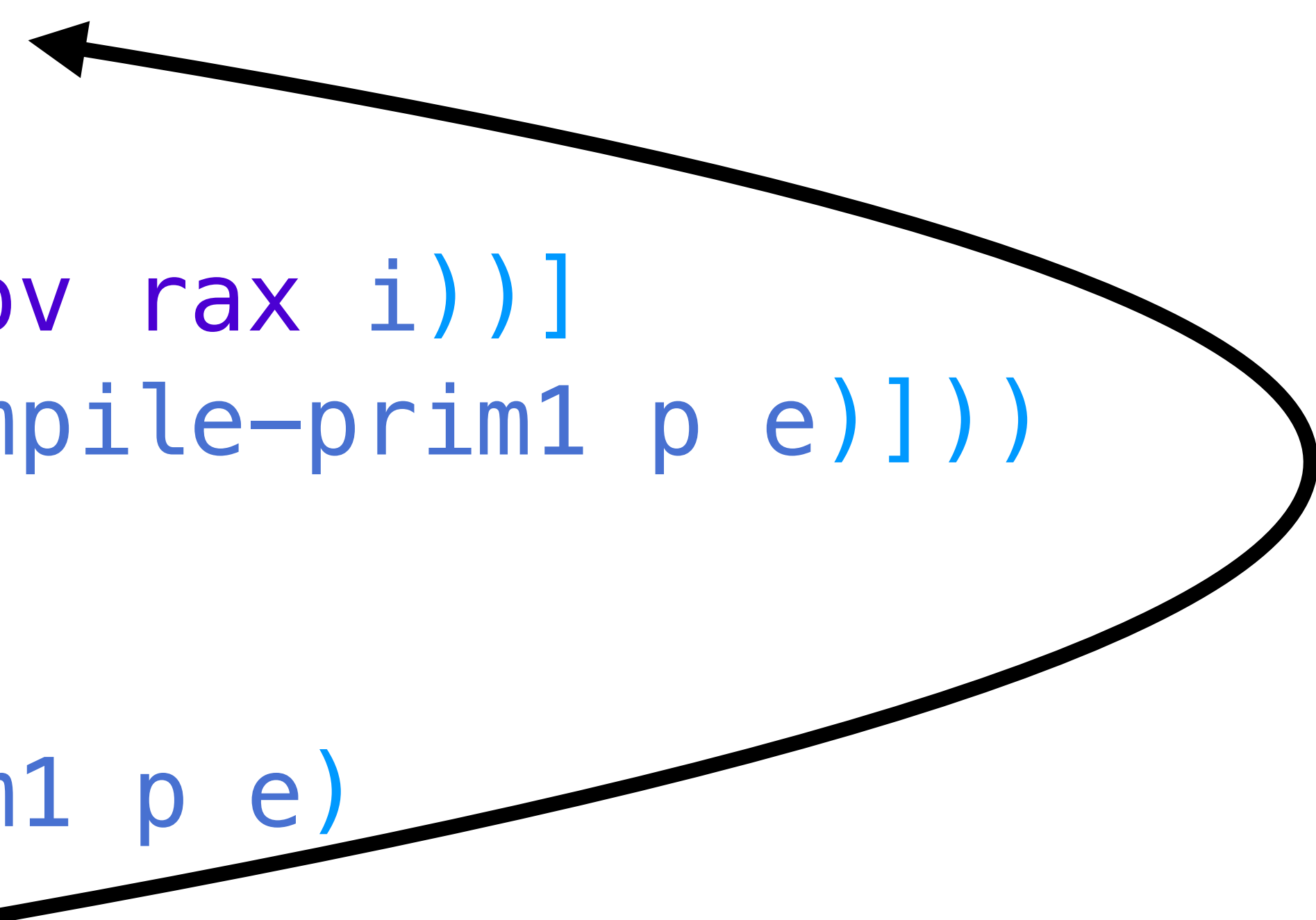
```
;; Invariant:  
;; Produces instructions that when run,  
;; leave e's value in rax  
(compile-e e)
```

Blackmail compiler

Compiling expressions

:: Expr -> Asm

```
(define (compile-e e)  
  (match e  
    [(Lit i) (seq (Mov rax i))]  
    [(Prim1 p e) (compile-prim1 p e)]))
```



:: Op1 Expr -> Asm

```
(define (compile-prim1 p e)  
  (seq (compile-e e)  
        (compile-op1 p)))
```


Blackmail compiler

Compiling primitives

```
:: Op1 -> Asm
(define (compile-op1 p)
  (match p
    ['add1 (Add rax 1)]
    ['sub1 (Sub rax 1)]))
```

Blackmail compiler

Example

```
> (asm-interp (compile (parse '(add1 (add1 40)))))  
42
```

For next time

- Read Con, Dupe notes
- Study Abscond, Blackmail notes notes
- Take ELMS quiz

CMSC 430 - 10 Sept 2026

Our first compilers

- Quick Blackmail review
- Con - conditional expressions
- Dupe - booleans, revised conditional