

CMSC433, Spring 2002 Programming Language Technology and Paradigms Distributed Computing & RMI

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Administrivia

- Project 3 due tomorrow
- Commentary for project 2 due next Wed.
 - submit as [project 12](#)
- Project 4 posted by Thursday
 - [distributed computing with RMI](#)
- Read Ch. 15, pages 973-979, in Eckel
 - [Readings web page also has additional distributed computing reading](#)
- Talk about exam Thursday
- Back to project 2 discussion

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Is distributing computing different?

- What kinds of distributed computing environments exist?
- Ways in which distributed computing is different
 - [Addressing objects](#)
 - [Latency](#)
 - [Partial failure](#)
 - [Concurrency](#)

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Distributed computing environments

- Usually refers to multiple CPU's
- Interprocessor communication via shared address space or message passing?
- On same chip, in same room, or across the Internet?
 - [latency, failure modes](#)

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Existing environments

- Seti @ Home, Entropia
- Server for search engine
- My laptop, PDA, cell phone, MP3 player, digital camera, ...

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Types of failure

- Machine sleeps
 - [wakes up, recovers state](#)
- Machine crash or failure
 - [machine may reboot and rejoin](#)
- Network partition
 - [network may heal](#)

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Uniform view of distributed objects

- Some objects are remote, some are local
 - Doesn't really matter to user of object
 - Objects might transparently migrate
- Design doesn't have to take object distribution into account
- Failure and performance issues don't belong in the design
- The interface doesn't change if an object is remote

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Uniform view

- Not appropriate for
 - wide area networks,
 - consumer electronics,
 - portable devices
- Appropriate for some local area networks
 - but robust distributed applications plan for failure
 - even if all objects local
 - means in interfaces, not just implementation

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Memory access

- Can we make the fact that an object is remote transparent?
- Perhaps for objects
 - What about int's ?
 - What about char *'s?
- If you can't directly access fields and create pointers to them
 - then it's not transparent

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Partial failure

- Computers fail
- OS's crash
- Networks fail
- PDA's get turned off or taken out of the room
- Often no warning, and each hardware/software component can fail independently of all others
 - and no other component may be able to tell exactly what happened

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Queue example

- Want to add x to remote queue q
 - `q.enqueue(x)`
- Operation could fail
- Want to reliably enqueue x

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Queue example

```
while (true) {  
    try {  
        q.enqueue(x);  
        break;  
    }  
    catch (RemoteException e) {}  
}
```

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Partial failure

- Object was enqueued, but failure occurred during return message
- Could enqueue x multiple times
- Solution?
 - Need a request tag so that duplicate enqueue requests can be detected
- Real question here is how much does it cost to make the remote call look the same as a local call

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Concurrency

- Distributed computations mandate concurrency
 - objects at different locations can't be stopped from executing concurrently
 - even less control than with multi-threading
 - no single point of control (hardware, OS, thread scheduler, etc.)

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