

CMSC433, Fall 2002 Programming Language Technology and Paradigms Distributed Computing

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Is distributing computing different?

- What kinds of distributed computing environments exist?
Ways in which distributed computing is
 - Addressing objects
 - Latency
 - Partial failure
 - Concurrency

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Distributed computing environments

- Usually refers to multiple CPU's
communication via shared
address space or message passing?
On same chip, in same room, or across the
 - latency, failure modes

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Existing environments

- Seti @ Home, Entropia
- Server for search engine
- My laptop, PDA, cell phone, MP3 player, digital camera, ...

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Types of failure

- Machine sleeps
 - wakes up, recovers state
- Machine crash or failure
 - machine may reboot and rejoin
- Network partition
 - network may heal

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Uniform view of distributed objects

- Some objects are remote, some are local
 - Doesn't really matter to user of object
Objects might transparently migrate
- Design doesn't have to take object distribution

Failure and performance issues don't belong in the

The interface doesn't change if an object is remote

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Uniform view

- Not appropriate for
 - wide area networks,
 - consumer electronics,
 - portable devices
- Appropriate for some local area networks
 - but robust distributed applications plan for

even if all objects local
means in interfaces, not just implementation

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Memory access

- Can we make the fact that an object is remote transparent?
Perhaps for objects
 - What about int's ?
 - What about char *'s?
- If you can't directly access fields and create
 - then it's not transparent

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Partial failure

- Computers fail
- OS's crash
- Networks fail
- PDA's get turned off or taken out of the

Often no warning, and each
hardware/software component can fail
independently of all others

- and no other component may be able to tell
exactly what happened

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Queue example

- Want to add x to remote queue q
 - q.enqueue(x)
- Operation could fail
- Want to reliably enqueue x

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Queue example

```
while (true) {  
    try {  
        q.enqueue(x);  
        break;  
    }  
    catch (RemoteException e) {}  
}
```

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- Object was enqueued, but failure occurred during return message
- Could enqueue x multiple times
- Solution?
 - Need a request tag so that duplicate requests can be detected
- Real question here is how much does it cost to make the remote call look the same as a

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Concurrency

- Distributed computations mandates
 - objects at different locations can't be stopped from executing concurrently
 - even less control than with multi-threading
 - no single point of control (hardware, OS, thread)

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Latency

- Making a call to an object on a remote machine is much more expensive than a
 - several orders of magnitude
- Leading to overall system performance problems if there are “too many” remote

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