

Synchronization issues

- Locks
- **synchronized** statements and methods
- **wait** and **notify**
- Deadlock

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Locks

- *All* objects can be locked
- Only one thread can hold a lock on an object
 - other threads block until they can acquire it
- If your thread already holds a lock on an object
 - can lock it a second time
 - object not unlocked until both locks released
- No way to only attempt to acquire a lock

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Synchronized methods

- A method can be synchronized
 - add **synchronized** modifier before return type
- Obtains a lock on object referenced by **this**, before executing method
 - releases lock when method completes
- For a **static synchronized** method
 - locks the class object

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Synchronized statement

- **synchronized (obj) { statements }**
- Obtains a lock on **obj** before executing statements in block
- Releases lock once block completes
- Provides finer grained control than synchronized method
- Allows locking arguments to a method

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Synchronization example

```
class SyncTest extends Thread {
    String msg;
    public SyncTest(String s) {
        msg = s;
        start();
    }
    public void run() {
        synchronized (SyncTest.class) {
            System.out.print("{ " + msg);
            try { Thread.sleep(1000); }
            catch (InterruptedException e) {}
            System.out.println("}");
        }
    }
    public static void main(String[] args) {
        new SyncTest("Hello"); new SyncTest("Synchronized");
        new SyncTest("World");
    }
}
```

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Wait and Notify

- Must be called inside **synchronized** method or block of statements (on the synchronized object)
- **a.wait()**
 - gives up lock(s) on **a**
 - adds thread to wait set for **a**
 - suspends thread
- **a.wait(int m)**
 - limits suspension to **m** milliseconds

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Wait and Notify (cont.)

- **a.notify()** resumes one thread from **a**'s wait list
 - and removes it from wait set
 - no control over which thread
- **a.notifyAll()** resumes *all* threads on **a**'s wait list
- resumed threads must reacquire lock before continuing
- **wait** doesn't give up locks on any other objects
 - e.g., acquired by methods that called this one

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Producer/Consumer Example

```
// Won't work. Needs to handle
// InterruptedException

public class ProducerConsumer {
    private boolean ready = false;
    private Object obj;

    public ProducerConsumer() { }
    public ProducerConsumer(Object o) {
        obj = o; ready = true; }

    synchronized void produce(Object o) {
        while (ready) wait();
        obj = o; ready = true;
        notifyAll(); }

    synchronized Object consume() {
        while (!ready) wait();
        ready = false;
        notifyAll();
        return obj; }
}
```

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Changed example

```
synchronized void produce(Object o) {
    while (full) {
        wait();
        if (full) notify(); }
    obj = o; full = true;
    notify();
}

synchronized Object consume() {
    while (!full) {
        wait();
        if (!full) notify(); }
    full = false;
    notify();
    return obj;
}
```

Bad design – no
guarantee about who
will get woken up

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A Better Solution

```
void produce(Object o) {
    while (buffFull) { synchronized (empty) {
        try { empty.wait(); }
        catch (InterruptedException e) { }
    }}
    obj = o; buffFull = true;
    synchronized (full) {
        full.notify(); }
}

Object consume() {
    while (!buffFull) { synchronized (full) {
        try { full.wait(); }
        catch (InterruptedException e) { }
    }}
    Object o = obj; buffFull = false;
    synchronized (empty) {
        empty.notify(); }
}
```

Use two objects,
empty and **full**, to
allow two different
wait sets

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notify() vs. notifyAll()

- Very tricky to use notify() correctly
 - notifyAll() much safer
- Need:
 - all waiters are equal
 - each notify only needs to wake up 1 thread
 - handle **InterruptedException** correctly

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InterruptedException Example

- Threads t1 and t2 are waiting
- Thread t3 performs a notify
 - thread t1 is selected
- Before t1 can acquire lock, t1 is interrupted
- t1's call to wait throws InterruptedException
 - t1 doesn't process notification
 - t2 doesn't wake up

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Handling InterruptedException

```
synchronized (this) {  
    while (!ready) {  
        try { wait(); }  
        catch (InterruptedException e) {  
            notify();  
            throw e; }  
        // do whatever  
    }  
}
```

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Deadlock

- Quite possible to create code that deadlocks
 - Thread 1 holds lock on **A**
 - Thread 2 holds lock on **B**
 - Thread 1 is trying to acquire a lock on **B**
 - Thread 2 is trying to acquire a lock on **A**
 - Deadlock!
- Not easy to detect when deadlock has occurred
 - other than by the fact that nothing is happening

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A common multi-threading bug

- Threads might cache values
- Obtaining a lock forces the thread to get fresh values
- Releasing a lock forces the thread to flush out all pending writes
- **volatile** variables are never cached
- **sleep(...)** doesn't force fresh values
- Many compilers don't perform these optimizations
 - but some do (Hotspot?)
- Problem might also occur with multiple CPUs

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Guidelines to simple/safe multi-threaded programming

- Synchronize access to shared data
- Don't hold a lock on more than one object at a time
 - could cause deadlock
- Hold a lock for as little time as possible
 - reduces blocking waiting for locks
- While holding a lock, don't call a method you don't understand
 - e.g., a method provided by someone else, especially if you can't be sure what it locks

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Guidelines (cont.)

- Have to go beyond these guidelines for more complex situations
 - but need to understand threading and synchronization well
- Recommended book for more details:
 - *Concurrent Programming in Java*, by Doug Lea

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