

Announcements

- Program #5 Due Thursday, May 12 (8:00pm)
- Final Exam
 - Monday, May 16
 - 4:00-6:00 pm
 - in JMP 3201
- Tuesday, May 10 is the last new material for the semester – Thursday will be review
- Reading
 - Notes

Improving Code Performance

- This lecture is not about algorithms, but rather
 - how to convert algorithms into efficient code
 - how to refine code to make it run faster
- Helpful to know a bit about
 - what a compiler can and can't do
 - what takes time on the hardware
- Key Ideas
 - Be systematic and data driven in what you do
 - Easy to get into making random code changes
 - Program spend their time in
 - loops (or recursion)
 - a sequence of code executed once is fast
 - doing I/O

Know Your Compiler

- Compilers can be asked to "optimize" your code
 - -O flag (and -On) enables the optimizer
 - they are supposed to never break your code
 - only safe changes to the program are permitted
- Limits on compiler optimizations
 - should never alter correct programs
 - may discover latent bugs though
 - limited understanding of the program
 - you are much smarter than the best compiler
 - need to compile program quickly

Understanding Modern Processors

- **Pipelining**
 - parts of multiple instructions running at once
 - decode instruction, load from memory
- **Superscalar processors**
 - Can execute 2 or more instructions at once
- **Branch prediction**
 - Computer guess which way a branch will go
 - Allows pipeline to stay full
- **Cache memory**
 - keep copies of recently accessed memory in fast storage
 - recently accessed memory is much faster
- **Floating point takes much longer than integer**
 - typically 20-30x for the same operation

Sample Compiler Optimization

- Original Code

```
int i, j, k;  
....  
for (i=0; i < 200; i++) {  
    a[2*j+i] = j * k + i;  
}
```

- Optimized Code

```
int i, j, k;  
....  
int temp1 temp2;  
temp1= 2 * j;  
temp2 = j * k;  
for (i=0; i < 200; i++) {  
    a[temp1+i] = temp2 + i;  
}
```

Typical Compiler Optimizations

- Code motion
 - previous example
- Loop Unrolling
 - often faster if a loop body is a bit bigger
 - modern computers can do several instructions at once
 - convert:
for (i=0; i < n; i++)
 c[i] = a[i] + b[i];
– to:
for (i=0; i < n; i+=2) {
 c[i] = a[i] + b[i];
 c[i+1] = a[i+1] + b[i+1];
}
- Dead-Code Elimination
 - if can infer code is never run, remove it

Limits to Compiler Optimization

- Pointer aliases

```
void func1(int *xp, int *yp) {  
    *xp += *yp;  
    *xp += *yp;  
}
```

- simplify to ?

```
void func1(int *xp, int *yp) {  
    *xp += 2 * *yp;  
}
```

- Function Side effects

- given:

- return $f(x) + f(x) + f(x) + f(x)$;

- simplify to ?

- return $4 * f(x)$;

Common Code Changes - Loop Bound Checks

- Consider

```
for (i=0; i < strlen(data); i++) {  
    if (data[i] > 'A' && data[i] < 'Z') data[i] -= ('A' - 'a');  
}
```

- `strlen` is called each time, but doesn't change in loop
 - results in $O(N^2)$ performance

- Change to

```
int limit = strlen(data);  
for (i=0; i < limit; i++) {  
    if (data[i] > 'A' && data[i] < 'Z') data[i] -= ('A' - 'a');  
}
```

- Compiler could not figure this one out

Common Code Changes

Eliminate strlen completely

- Recall

```
int limit = strlen(data);  
for (i=0; i < limit; i++) {  
    if (data[i] > 'A' && data[i] < 'Z') data[i] -= ('A' - 'a');  
}
```

- Can eliminate strlen call

```
for (i=0, ptr = data; *ptr; ptr++) {  
    if (*ptr > 'A' && *ptr < 'Z') *ptr -= ('A' - 'a');  
}
```

Common Code Transformation - Don't write small functions

- Functions calls take time
 - Recall last week's code to setup calls stacks
- 1-2 line functions
 - take more time for the call than work
 - provide relatively little isolation of ideas
- Solution:
 - Define abstractions so work is larger than a few lines

Approach To Applying These Techniques

- For a big program
 - Hard to know where to start to optimize
- Start with gprof data
 - concentrate in parts of the program that take the most time
- Ahmdals Law
 - if tuning a function that takes α share of the time
 - and make it k times faster
 - $T_{\text{new}} = (1 - \alpha) T_{\text{old}} + (\alpha T_{\text{old}})/k$
 - Basically are limited by how big α is