

# CMSC 132: Object-Oriented Programming II

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## Hashing

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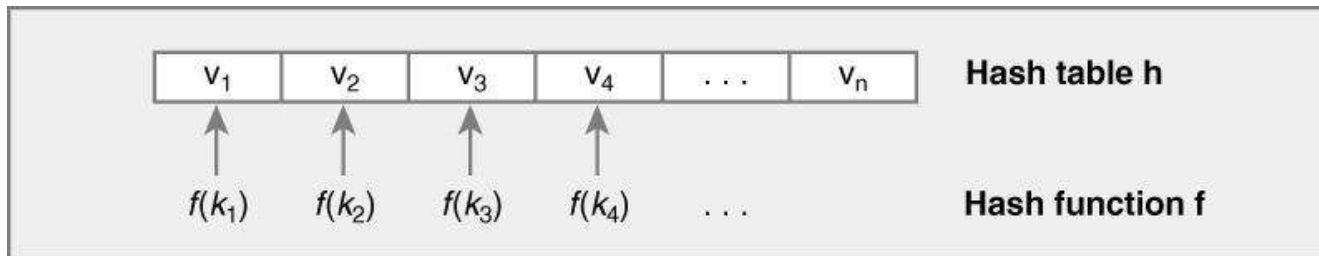
# Hashing

## ■ Hashing

- Hashing function  $\rightarrow$  function that maps data to a value (e.g., integer)
- Hash Code/Hash Value  $\rightarrow$  value returned by a hash function
- Hash functions can be used to speed up data access
- We can achieve  $O(1)$  data access using hashing

## ■ Approach

- Use **hash function** to convert key into number (**hash value**) used as index in **hash table**



# Hashing

## ■ Hash Table

- Array indexed using hash values
- Hash table  $A$  with size  $N$
- Indices of  $A$  range from 0 to  $N-1$
- Store in  $A[\text{hashValue} \% N]$

*Hash table  $h$*

|            |           |
|------------|-----------|
| $h[0]$     | $\Lambda$ |
| $h[1]$     | $\Lambda$ |
| $\dots$    | $\dots$   |
| $h[N - 1]$ | $\Lambda$ |

| <i>Location</i> | <i>Key</i> |
|-----------------|------------|
| 0               | $\Lambda$  |
| 1               | 10         |
| 2               | 15         |
| 3               | 20         |
| 4               | $\Lambda$  |
| $\dots$         |            |

# Hash Function

- Hash Function → Function for converting key into hash value
- For hash table of size N
  - Must reduce hash value to  $0..N - 1$
  - Can use modulo operator →  $\text{hash value} = \text{Math.abs}(\text{keyValue} \% N)$
- Example Problem
  - Assign 4 parking spaces to 4 people using
    - $h(\text{key}) = \text{keyValue} \% 4$
  - What happens if we have 4 spaces and 8 people?
    - Collision → Same hash value for multiple keys
- Bucket
  - Each table entry can be referred to as a bucket
  - In some implementations the bucket is represented by a list (those elements hashing to the same bucket are placed in the same list)
- Properties of a Good Hash Function
  - Distributes (scatters) values uniformly across range of possible values
  - It is not expensive to compute

# Scattering Hash Values

- Hash function should **scatter** hash values uniformly across range of possible values
  - Reduces likelihood of conflicts between keys
- $\text{Hash}(\text{<everything>}) = 0$ 
  - Satisfies definition of hash function
  - But not very useful (all keys at same location)
- Could use  $\text{Math.abs}(\text{keyValue} \% N)$ 
  - Might not distribute values well
  - Particularly if  $N$  is a power of 2

# Scattering Hash Values

- **Multiplicative congruency method**
  - Produces good hash values
  - Hash value =  $\text{Math.abs}((a * \text{keyValue}) \% N)$
  - Where
    - N is table size
    - a is large prime number

# Hash Function

## ■ Example

$\text{hash}(\text{"apple"}) = 5$

$\text{hash}(\text{"watermelon"}) = 3$

$\text{hash}(\text{"grapes"}) = 8$

$\text{hash}(\text{"kiwi"}) = 0$

$\text{hash}(\text{"strawberry"}) = 9$

$\text{hash}(\text{"mango"}) = 6$

$\text{hash}(\text{"banana"}) = 2$

## ■ Perfect hash function

- Unique values for each key

|   |            |
|---|------------|
| 0 | kiwi       |
| 1 |            |
| 2 | banana     |
| 3 | watermelon |
| 4 |            |
| 5 | apple      |
| 6 | mango      |
| 7 |            |
| 8 | grapes     |
| 9 | strawberry |

# Hash Function

## ■ Suppose now

hash("apple") = 5

hash("watermelon") = 3

hash("grapes") = 8

hash("kiwi") = 0

hash("strawberry") = 9

hash("mango") = 6

hash("banana") = 2

hash("orange") = 3

## ■ Collision

- Same hash value for multiple keys

0

kiwi

1

2

banana

3

watermelon

4

5

apple

6

mango

7

8

grapes

9

strawberry



# Beware of % (Modulo Operator)

- The % operator is integer remainder

$$x \% y == x - y * (x / y)$$

- Result may be negative

$$-|y| < x \% y < +|y|$$

- $x \% y$  has same sign as  $x$

- $-3 \% 2 = -1$

- $-3 \% -2 = -1$

- Use `Math.abs( x % N )` and not `Math.abs( x ) % N`

- About absolute value in Java

- `Math.abs(Integer.MIN_VALUE) == Integer.MIN_VALUE !`

- Will happen 1 in  $2^{32}$  times (on average) for random int values

# Hashing in Java

- Object class has built-in support for hashing
  - Method `int hashCode( )` provides
    - Numerical hash value for any object
    - 32-bit signed int
- Default `hashCode( )` implementation
  - Usually just address of object in memory
- Can override with new user definition
  - Must work with `equals( )`
  - Must satisfy the “hash code contract”

# Java Hash Code Contract

- Java Hash Code Contract

if `a.equals(b) == true`, then we must **guarantee**  
`a.hashCode() == b.hashCode()`

- Inverse is not true

`!a.equals(b)` does not imply  
`a.hashCode() != b.hashCode()`

(Though Java libraries may be more efficient)

- Converse is also not true

`a.hashCode() == b.hashCode()`  
does not imply `a.equals(b) == true`

- `hashCode()`

- Must return same value for object in each execution, provided information used in `equals()` comparisons on the object is not modified

# When to Override hashCode

- You must write classes that satisfy the Java Hash Code Contract
- You will run into problems if you don't satisfy the Java Hash Code Contract and use classes that rely on hashing (e.g., HashMap, HashSet)
  - Possible problem – You add an element to a set but cannot find it during a lookup operation
  - See code distribution example
- Does the default equals and hashCode satisfy the contract? **Yes!**
- If you implement the Comparable interface you should provide the appropriate equals method which leads to the appropriate hashCode method

# Java hashCode()

- **Implementing hashCode( )**
  - **Include only information used by equals( )**
    - Else 2 “equal” objects → different hash values
  - **Using all / more of information used by equals( )**
    - Help avoid same hash value for unequal objects
- **Example hashCode( ) functions**
  - **For pair of Strings**
    - 1<sup>st</sup> letter of 1<sup>st</sup> str
    - 1<sup>st</sup> letter of 1<sup>st</sup> str + 1<sup>st</sup> letter of 2<sup>nd</sup> str
    - Length of 1<sup>st</sup> str + length of 2<sup>nd</sup> str
    - $\sum \text{letter(s) of 1}^{\text{st}} \text{ str} + \sum \text{letter(s) of 2}^{\text{nd}} \text{ str}$

# Art and Magic of hashCode()

- There is no “right” hashCode function
  - Art involved in finding good hashCode function
  - Also for finding hashCode to hashBucket function
- From java.util.HashMap

```
static int hashBucket(Object x, int N) {  
    int h = x.hashCode();  
    h += ~(h << 9);  
    h ^= (h >>> 14);  
    h += (h << 4);  
    h ^= (h >>> 10);  
    return Math.abs(h % N);  
}
```

# Pre-filtering with hashCode

- hashCode( ) provides **pre-filter** for equals( )
  - Check equals( ) only if hashCode( ) is identical
  - Example
    - if ( a.hashCode( ) == b.hashCode( ) )  
          result = a.equals( b );  
          else result = false;
  - Efficient if hashCode( ) is faster than equals( )