

CMSC 330: Organization of Programming Languages

Objects and Functional Programming

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OOP vs. FP

- ▶ Object-oriented programming (OOP)
 - Computation as interactions between objects
 - Objects encapsulate state, which is usually mutable
 - Accessed / modified via object's public methods
- ▶ Functional programming (FP)
 - Computation as evaluation of functions
 - Mutable data used to improve efficiency
 - Higher-order functions implemented as closures
 - Closure = function + environment

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Relating Objects and Closures

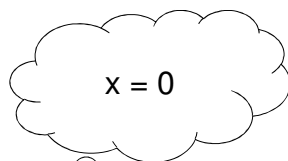
- ▶ An object...
 - Is a collection of fields (data)
 - ...and methods (code)
 - When a method is invoked
 - Method has implicit **this** parameter that can be used to access fields of object
- ▶ A closure...
 - Is a pointer to an environment (data)
 - ...and a function body (code)
 - When a closure is invoked
 - Function has implicit environment that can be used to access variables

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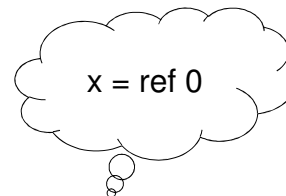
Relating Objects and Closures (cont.)

```
class C {  
  int x = 0;  
  void set_x(int y) { x = y; }  
  int get_x() { return x; }  
}
```



```
C c = new C();  
c.set_x(3);  
int y = c.get_x();
```

```
let make () =  
  let x = ref 0 in  
    ( (fun y -> x := y),  
      (fun () -> !x) )
```



```
fun y -> x := y  fun () -> !x
```

```
let (set, get) = make ();;  
set 3;;  
let y = get ();;
```

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Encoding Objects with Functions

- We can apply this transformation in general

```
class C { f1 ... fn; m1 ... mn; }
```

- becomes

```
let make () =
  let f1 = ...
  ...
  and fn = ... in
  ( fun ... , (* body of m1 *)
    ...
    fun ... , (* body of mn *)
  )
```

} Tuple
containing
closures

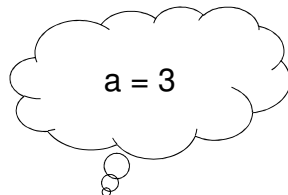
- make () is like the constructor
- The closure environment contains the fields

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Relating Closures and Objects

```
let app f x = f x
```



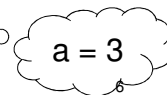
```
fun b -> a + b
```

```
let add a b = a + b;;
let f = add 3;;
app f 4;;
```

```
interface F {
  Integer eval(Integer y);
}
class C {
  static Integer app(F f, Integer x) {
    return f.eval(x);
  }
}
```

```
class G implements F {
  Integer a;
  G(Integer a) { this.a = a; }
  Integer eval(Integer y) {
    return new Integer(a + y);
  }
}
```

```
F adder = new G(3);
C.app(adder, 4);
```



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Encoding Functions with Objects

- ▶ We can apply this transformation in general

```
... (fun x -> (* body of fn *)) ...  
let h f ... = ...f y...
```

- becomes

```
interface F { Object eval(Object x); }  
class G implements F {  
    Object eval(Object x) { /* body of fn */ }  
}  
class C {  
    Typ h(F f, ...) {  
        ...f.eval(y)...  
    }  
}
```

- F is the interface to the callback
- G represents the particular function

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Code as Data

- ▶ Closures and objects are related
 - Both of them allow
 - Data to be associated with higher-order code
 - Pass code around the program
- ▶ The key insight in all of these examples
 - Treat code as if it were data
 - Allowing code to be passed around the program
 - And invoked where it is needed (as callback)
- ▶ Approach depends on programming language
 - Higher-order functions (OCaml, Ruby, Lisp)
 - Function pointers (C, C++)
 - Objects with known methods (Java)

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Code as Data (cont.)

- ▶ This is a powerful programming technique
 - Solves a number of problems quite elegantly
 - Create new control structures (e.g., Ruby iterators)
 - Add operations to data structures (e.g., visitor pattern)
 - Event-driven programming (e.g., observer pattern)
 - Keeps code separate
 - Clean division between higher & lower-level code
 - Promotes code reuse
 - Lower-level code supports different callbacks

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An Integer List Abstraction in Java

```
public class MyList {  
    private class ConsNode {  
        int head; MyList tail;  
        ConsNode (int h, MyList l) { head = h; tail = l; }  
    }  
  
    private ConsNode contents;  
  
    public MyList () {  
        contents = null;  
    }  
  
    public MyList (int h, MyList l) {  
        contents = new ConsNode (h, l);  
    }  
  
    public MyList cons (int h) {  
        return (new MyList (h, this));  
    }  
  
    public int hd () {  
        return contents.head;  
    }  
  
    public MyList tl () {  
        return contents.tail;  
    }  
  
    public boolean isNull () {  
        return (contents == null);  
    }  
}
```

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Recall a Useful Higher-Order Function

```
let rec map f = function
  [] -> []
| (h::t) -> (f h)::(map f t)
```

- ▶ Map applies an arbitrary function **f**
 - To each element of a list
 - And returns the resulting modified list
- ▶ Can we encode this in Java?
 - Using object oriented programming

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A Map Method for Lists in Java

- ▶ Problem – Write a map method in Java
 - Must pass a function into another function
- ▶ Solution
 - Can be done using an object with a **known** method
 - Use **interface** to specify what method must be present

```
public interface IntFunction {
  int eval(int arg);
}
```

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A Map Method for Lists (cont.)

► Examples

- Two classes which both implement Function interface

```
class AddOne implements IntFunction {  
    int eval (int arg) {  
        return (arg + 1);  
    }  
}
```

```
class MultTwo implements IntFunction {  
    int eval(int arg) {  
        return (arg * 2);  
    }  
}
```

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The New List Class

```
class MyList {  
    ...  
    public MyList map (IntFunction f) {  
        if (this.isNull()) return this;  
        else return (this.tl().map(f).cons (f.eval (this.hd())));  
    }  
}
```

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Applying Map To Lists

- ▶ Then to apply the function, we just do

```
MyList l = ...;  
MyList l1 = l.map(new AddOne());  
MyList l2 = l.map(new MultTwo());
```

- We make a new object
 - That has a method that performs the function we want
- This is sometimes called a **callback**
 - Because **map** “calls back” to the object passed into it
- But it’s really just a higher-order function
 - Written more awkwardly

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We Can Do This for Fold Also!

- ▶ Recall fold

```
let rec fold f a = function  
  [] -> a  
| (h::t) -> fold f (f a h) t
```

- Fold accumulates a value (in **a**) as it traverses a list
- **f** is used to determine how to “fold” the head of a list into **a**
- ▶ This can be done in Java using an approach similar to **map**!

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A Fold Method for Lists in Java

- ▶ Problem – Write a fold method in Java
 - Must pass a function into another function
- ▶ Solution
 - Can be done using an object with a **known** method
 - Use **interface** to specify what method must be present

```
public interface IntBinFunction {  
    Integer eval(Integer arg1, Integer arg2);  
}
```

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A Fold Method for Lists (cont.)

- ▶ Examples
 - A classes which implements **IntBinFunction** interface

```
class Sum implements IntBinFunction {  
    Integer eval(Integer arg1, Integer arg2) {  
        return new Integer(arg1 + arg2);  
    }  
}
```

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The New New List Class

```
class MyList {
    ...
    public MyList map (IntFunction f) {
        if (this.isNull()) return this;
        else return (this.tl()).map(f).cons (f.eval (this.hd()));
    }

    public int fold (IntBinFunction f, int a) {
        if (this.isNull()) return a;
        else return (this.tl()).fold(f, f.eval(a, this.hd()));
    }
}
```

Applying Fold to Lists

- ▶ To apply the fold function, we just do this:

```
MyList l = ...;
int s = l.fold (new Sum(), 0);
```

- ▶ The result is that s contains the sum of the elements in l