CMSC330 Spring 2014 Midterm 2 Solutions

1. (16 pts) OCaml Types and Type Inference

Give the type of the following OCaml expressions:

a. (2 pts) fun b  -> b + 8  
   Type = int -> int  

b. (3 pts) fun b -> b 8  
   Type = (int -> 'a) -> 'a  

Write an OCaml expression with the following type:

c. (2 pts)  'a -> 'a -> 'a list list  
   Code = fun x y -> [ [x ; y] ]  

d. (3 pts)  (int -> int) -> (int -> int)  
   Code = fun x y -> (x 1) + y  

Give the value of the following OCaml expressions. If an error exists, describe it

e. (3 pts) (fun k t m -> [m; t]) 2 4 6  
   Value / Error = [6 ; 4]  

f. (3 pts) (fun (k::t::m) -> (m, t)) [2;4;6;8]  
   Value / Error = ( [6;8] , 4 )  

2. (10 pts) OCaml higher-order & anonymous functions

Using either map or fold and an anonymous function, write a curried function process
which when given a predicate function p, a compute function f, and a list lst, returns a list
of the results of applying f to the elements of lst, but only where p is true for that element
of the list. If p is false for an element of the list, it does not contribute to the final result.
The relative order of the elements in the new list must the same as in lst.

Your function must run in linear time. You may not use any library functions, with the
exception of the List.rev function, which reverses a list in linear time. You may not use
imperative OCaml (i.e., no ref variables). For example, the results of calling process will
be as follows:

Example:

let p x = x > 2;;  
let f x = x+1;;  
process p f [ ] = [ ]  
process p f [ 1; 2 ] = [ ]  
process p f [ 1; 3; 2; 4 ] = [ 4; 5 ]

let rec map f l = match l with  
  | [] -> []  
  | (h::t) -> (f h)::(map f t)

let rec fold f a l = match l with  
  | [] -> a  
  | (h::t) -> fold f (f h) t

let process p f lst = List.rev (fold (fun a h -> if (p h) then (f h)::a else a) [ ] lst)
3. (18 pts) OCaml programming

Write a function `allPairs` with type (`'a list -> 'a list list`) which given a list `lst` with `n` elements, returns a list of all `n^2` pairs of elements in `lst` (in any order), where each pair of elements is in a list of length 2.

You may not use any library functions, with the exception of the @ function (which concatenates two lists). You may use helper functions. You may use map or fold, but it is not required. Your function must execute in polynomial time. You may not use imperative OCaml constructs.

Examples:

```
let rec help1 x y = match y with
  [] -> []
| h::t -> [x;h] :: (help1 x t) ;;

let rec help2 all z = match z with
  [] -> []
| h::t -> (help1 h all) @ (help2 all t) ;;

let allPairs lst = help2 lst lst ;;
```

Sample Solutions

```
let rec help1 x y = match y with
  [] -> []
| h::t -> [x;h] :: (help1 x t) in
let rec help2 z = match z with
  [] -> []
| h::t -> (help1 h lst) @ (help2 t) in
help2 lst ;;

let help1 x y = fold (fun a h -> [x:h];:a) [] y ;;
let allPairs lst = fold (fun a h -> (help1 h lst) @ a) [] lst ;;

let help1 x = fold (fun a h -> [x:h];:a) [] lst in
  fold (fun a h -> (help1 h) @ a) [] lst ;;

let allPairs lst =
  let help1 x = map (fun h -> [x:h]) lst in
  fold (fun a h -> (help1 h) @ a) [] lst ;;

let allPairs lst =
  let help1 x = map (fun h -> [x:h]) lst in
  fold (fun a h -> (map (fun z -> [h;z]) lst) @ a) [] lst ;;
```
4. (8 pts) OCaml polymorphic types

Consider the OCaml type `bitwiseExp` implementing bitwise expressions (think of it as a binary tree where all leaves are 0’s or 1’s). For all answers, your code must work in linear time (i.e., avoid multiple passes over the tree). You are not allowed to use any OCaml library functions except +. You may use helper functions.

```ocaml
type bitwiseExp =
  Zero (* 0 *)
| One (* 1 *)
| Pair of bitwiseExp * bitwiseExp (* b₁ • b₂ *)
```

a. (2 pts) Write an OCaml expression with type bitwiseExp that is equivalent to the expression “0 • (1 • 0)”

```
Pair (Zero, Pair (One, Zero))
```

b. (6 pts) Write a function `onesExp` of type `(bitwiseExp -> int)` that takes a bitwise expression and returns the number of 1’s in the expression.

```ocaml
let rec onesExp b = match b with
  Zero -> 0
| One -> 1
| Pair (x,y) -> (onesExp x) + (onesExp y)
```
5. (10 pts) Context free grammars.
Consider the following grammar (S = start symbol and terminals = a b % &):
\[ S \rightarrow a | S\%S | \&T \]
\[ T \rightarrow a | b \]

a. (1 pt each) Indicate whether the following strings are generated by this grammar
i. %a  Yes  No  (circle one)
ii. a%&b  Yes  No  (circle one)
iii. &a&b  Yes  No  (circle one)

b. (3 pts) Draw a parse tree for the string “&b%a”

```
S /
/ \S % S /
/ \ & T a
/ |   b
```

c. (4 pts) Is the grammar is ambiguous? Provide proof if you believe it is ambiguous.

Grammar is ambiguous. Multiple parse trees for strings such as a%a%a.
6. (14 pts) Using context free grammars.
   a. (6 pts) Given the following production for B, create a grammar that generates
      OCaml tuples of strings (for the 3 strings Joe, Laura, and Mike). Your grammar
      should be able to generate strings such as (“Joe”), (“Joe”, “Mike”), (“Joe”,
      “Laura”, “Laura”), etc...

      \[ B \rightarrow \text{“Mike”} \mid \text{“Laura”} \mid \text{“Joe”} \]
      \[ S \rightarrow (A) \]
      \[ A \rightarrow B \mid B, A \]

   b. (8 pts) Consider the following grammar:

      \[ S \rightarrow c \mid d \mid e \mid S\%S \mid S\@S \mid \&S \]

      Modify the grammar above to make the \% operator left associative, the \@ operator
      right associative, and make \@ have the lowest precedence and \& have the highest
      precedence.

      \[ S \rightarrow A\@S \mid A \]
      \[ A \rightarrow A\%B \mid B \]
      \[ B \rightarrow \&B \mid C \quad (*B \rightarrow \&C \mid C \text{ is also acceptable } *) \]
      \[ C \rightarrow c \mid d \mid e \]
7. (15 pts) Parsing
Consider the following grammar, where S, A, B are nonterminals, and a, b, c, d are terminals.

\[
S \rightarrow a \mid bB \mid Ac \\
A \rightarrow SAd \mid \epsilon \\
B \rightarrow c \mid b
\]

a. (8 pts) Calculate FIRST sets for S, A, and B

FIRST(S) = \{ a, b, c \} \\
FIRST(A) = \{ a, b, c, \epsilon \} \\
FIRST(B) = \{ c, b \}

b. (7 pts) Using pseudocode, write only the parse_A function found in a recursive descent parser for the grammar. You may assume the functions parse_S, parse_B already exist.

Use the following utilities:

<table>
<thead>
<tr>
<th>lookahead</th>
<th>Variable holding next terminal</th>
</tr>
</thead>
<tbody>
<tr>
<td>match(x)</td>
<td>Function to match next terminal to x</td>
</tr>
<tr>
<td>error()</td>
<td>Reports parse error for input</td>
</tr>
</tbody>
</table>

parse_A() { // your code starts here

    if ((lookahead == 'a') || (lookahead == 'b') || (lookahead == 'c')) {
        parse_S();
        parse_A();
        match('d');
    } else {
        ;
    }
}
8. (9 pts) Multithreading

Consider the following multithreaded Java 1.4 code. Assume there are multiple producer and consumer threads being executed in the program, but only a single Buffer object. Questions about the “last statement executed” by two threads x & y refer to the most recently executed statement by those threads at some arbitrary time during the program execution. It does not mean the last statement executed by a thread before the thread exits. If a situation is possible, you need to give an example of how it is possible (e.g., thread x gets to statement a, then thread y gets to statement b). If a situation is not possible, you need to explain why.

```
class Buffer {
    void produce(o) {
        synchronized (this) {
            1. while (...) wait( );
            2. ...
            3. notifyAll( );
            4. ...
        }
    }

    Object consume( ) {
        synchronized (this) {
            5. while (...) wait( );
            6. ...
            7. notifyAll( );
            8. ...
        }
    }
}
```

a. (3 pts) Is it possible given two threads x and y for the last statement executed by thread x to be statement 5, and the last statement executed by thread y to be statement 7? Explain your answer.

Yes. Thread x reaches statement 5 and releases the lock by calling wait, thread y then acquires lock and reaches statement 7.

b. (3 pts) Is it possible given two threads x and y for the last statement executed by thread x to be statement 7, and the last statement executed by thread y to be statement 6? Explain your answer.

No. Whichever thread enters synchronized region first (and reaches statement 6 or 7) must still be holding the lock. The other thread cannot acquire the lock to enter the synchronized region.

c. (3 pts) Is it possible given two threads x and y for the last statement executed by thread x to be statement 6, and the last statement executed by thread y to be statement 1? Explain your answer.

Yes. Thread y reaches statement 1 and releases the lock by calling wait, thread x then acquires lock and reaches statement 6.