Abstract Data Types

• Expose signature
  • operators to create, combine, and observe values

• Hide representation & implementations

• Advantages
  • Good engineering (libraries, code reuse, local reasoning)
  • Optimizations (choose implementation w.r.t. your code)

• Limitations
  • Require static typing
  • Different representations cannot mix
Objects

- Object-oriented programming (OOP)
  - Computation as interactions between objects

- An object...
  - Is a collection of fields (data)
  - ...and methods (code)
  - When a method is invoked
    - Method has implicit this parameter that can be used to access fields of object
Relating ADTs to Objects

Abstract Data Types:
define behavior for all sets

Objects:
- encoded as records
- fields define how this object interacts with other objects

```
type iSet = {
          isEmpty  : bool;
          insert   : int -> iSet;
          contains : int -> bool;
        } 
```

```
type set
val empty    : set

val isEmpty  : set -> bool
val insert   : set * int -> set
val contains : set * int -> bool
```

Observations on set argument
module ISetF : ISET =
struct
  type set = bool * (int -> bool)
  let empty = (true, fun x -> false)

  let isEmpty (e,_) = e

  let insert((_,f), i) =
    (false, if f i then f else
     fun j -> i = j || f j)

  let contains((_,f), i) = f i
end;;
Objects by Encoding

• We will consider an encoding of objects similar to the Object $\leftrightarrow$ FP encoding we saw before
  • Objects contain methods
    ➢ But we will define “top level” methods that take the object as an argument. These will call the object’s methods
  • This will let us see how to encode the self/this parameter
    ➢ Not something we looked at before

• In our examples, objects are *not* imperative
  • “Updates” to an object just return a new object with the update made
ISet as an object

type iSet = {  
    isEmpty : bool;  
    insert : int -> iSet;  
    contains : int -> bool;  
}

# let (x:iSet) = ... ;;  
- ...  
# let y = x.insert 1;;  
- ...  
# let z = insert_obj(y,2);;;  
- ... (* calls y’s insert method *)  
# y.contains 1;;  
- : bool = true  
# z.contains 2;;  
- : bool = false
Example: Insert Set

let insert_obj(s, n) =
    if s.contains(n) then s else {
        isEmpty = false;
        contains = fun i -> (i = n || s.contains(i));
        insert   = fun i -> insert_obj (this, i);
    }

Question: How can we encode this?
Answer: Use the fixpoint combinator Y
fix : (('a -> 'b) -> 'a -> 'b) -> 'a -> 'b
Example: Insert Set

let rec insert_obj(s, n) =
  if s.contains(n) then s else (fix (fun this _ -> { 
    isEmpty = false;
    contains = fun i -> (i = n || s.contains(i));
    insert   = fun i -> insert_obj (this, i);
  })) ()

Question: How can we encode this?
Answer: Use the fixpoint combinator Y

fix : (((a -> b) -> a -> b) -> a -> b
Quiz 1: Which is the Empty Set?

A. let empty_obj _ = (fix (fun this _ -> { isEmpty = false;
contains = fun i -> false;
insert = fun i -> insert_obj(this, i) })) ()

B. let empty_obj _ = (fix (fun this _ -> { isEmpty = true;
contains = fun i -> false;
insert = fun i -> insert_obj(this, i) })) ()

C. let empty_obj _ = (fix (fun this _ -> { isEmpty = true;
contains = fun i -> false;
insert = fun i -> this })) ()
Quiz 1: Which is the Empty Set?

B. let empty_obj _ = (fix (fun this _ -> {
  isEmpty = true;
  contains = fun i -> false;
  insert = fun i -> insert_obj(this, i)})) ()
Quiz 2: Which produces the union?

A. let union_obj(s1, s2) = (fix (fun this _ -> {  
isEmpty = s1.isEmpty || s2.isEmpty;  
contains = fun i -> s1.contains(i) || s2.contains(i);  
insert = fun i -> insert_obj(this, i)})) ()

B. let union_obj(s1, s2) = (fix (fun this _ -> {  
isEmpty = s1.isEmpty || s2.isEmpty;  
contains = fun i -> s1.contains(i) && s2.contains(i);  
insert = fun i -> insert_obj(this, i)})) ()

C. let union_obj(s1, s2) = (fix (fun this _ -> {  
isEmpty = s1.isEmpty && s2.isEmpty;  
contains = fun i -> s1.contains(i) || s2.contains(i);  
insert = fun i -> insert_obj(this, i)})) ()
Autognosis: An object can only access other objects through their public interface.

Problem: The representation of s1 s2 is unknown, so … … cannot be used for optimizations.

Solution: Expose more implementation details.

```
C. let union_obj(s1, s2) = (fix (fun this _ -> { 
       isEmpty   = s1.isEmpty && s2.isEmpty;
       contains  = fun i -> s1.contains(i) || s2.contains(i);
       insert    = fun i -> insert_obj(this, i)})) ()
```
Lack of Autognosis in ADTs

**Autognosis**: An object can only access other objects through their public interface.

- Each implementation can inspect representation
  - Used for optimizations (eg., sorted lists)

- Different implementations cannot interact
  - eg., cannot union `SeIt` and `ISetBST`
Quiz 3: The Even Set

A. `even_obj _ = (fix (fun this _ ->  {`
   `  isEmpty   = false;`
   `  contains  = fun i -> i mod 2 = 0;`
   `  insert    = fun i -> insert_obj(this, i)})) ()`

B. `even_obj _ = (fix (fun this _ ->  {`
   `  isEmpty   = true;`
   `  contains  = fun i -> true;`
   `  insert    = fun i -> insert_obj(this, i)})) ()`

C. `even_obj _ = (fix (fun this _ ->  {`
   `  isEmpty   = i mod 2 = 0;`
   `  contains  = fun i -> i mod 2 = 0;`
   `  insert    = fun i -> insert_obj(this, i)})) ()`
The Even Set: Flexibility

Flexibility:
- Objects accept any value that implements required methods.
- They can be easily extended with new (here infinite) representations.

A. even_obj _ = (fix (fun this _ -> {
  isEmpty = false;
  contains = fun i -> i mod 2 = 0;
  insert   = fun i -> insert_obj(this, i)})()) ()
Computation via *dynamic binding*: The function to be called is selected from the object record. Every object has a different `contains` field. Difficult to reason.
Computations are Object Interactions

\[
\text{union_obj(}
\text{even_obj(),}
\text{empty_obj().insert(3)})\text{.contains(3)}
\]

\[
\text{contains} = \text{fun } i \rightarrow i \mod 2 = 0
\]

\[
\text{contains} = \text{fun } i \rightarrow false
\]

\[
\text{contains} = \text{fun } i \rightarrow i == 3 || false
\]

\[
\text{contains} = \text{fun } i \rightarrow i == 3 || false || i \mod 2 = 0
\]
Quiz 4: ADTs vs. Objects

Used for data abstraction
(i.e., separate behavior from implementation)

A. ADTs  B. Objects
C. None    D. Both
Quiz 4: ADTs vs. Objects

Used for data abstraction (i.e., separate behavior from implementation)

**ADTs** use Type Abstraction:
expose a type whose representation is hidden.

**Objects** use Procedural Abstraction:
expose procedures available on each object.

A. ADTs  B. Objects  C. None  D. Both
Quiz 5: ADTs vs. Objects

Require Static Type System

A. ADTs
B. Objects
C. None
D. Both
ADTs require Static Types

A. ADTs  B. Objects
C. None  D. Both

```ocaml
type set
val empty : set

val isEmpty : set -> bool
val insert : set * int -> set
val contains : set * int -> bool
```
Quiz 5: ADTs vs. Objects

Require Static Type System

ADTs rely on Static Type System to define the Type Abstraction

Objects define abstractions via records: can be used in Static & Dependently typed lang.

A. ADTs  B. Objects
C. None    D. Both
Quiz 6: ADTs vs. Objects

Allow mixing different representations.

A. ADTs  
B. Objects  
C. None  
D. Both
Quiz 6: ADTs vs. Objects

Allow mixing different representations.

ADTs doesn’t permit interaction of different representations. Representation can be inspected, which is used for optimizations.

Objects cannot inspect representations (example union). Different representations can be mixed.

A. ADTs  B. Objects  C. None  D. Both
ADTs vs. Objects

ADTs are easier to reason.
- The implementation is statically known (unlike dynamic binding)
- Obey theories in math (abstract algebra) and PL (existential typing)

Objects are flexible and easy to extend.
- Only requirement to have appropriate fields