

CMSC 250

Discrete Structures

Induction

An Example - with an inequality

Example

Claim:

$$(\forall n \in \mathbb{Z}^{\geq 3}) [2n + 1 < 2^n]$$

Another Example with an inequality

Example

$$(\forall n \in \mathbb{Z}^{\geq 2})(\forall x \in \mathbb{R}^+) \left[1 + nx \leq (1 + x)^n \right]$$

An example with a recurrence relation

- Assume the following definition of a sequence:

$$a_1 = 1$$

$$(\forall k \geq 2) \left[a_k = a_{k-1} + (2k - 1) \right]$$

- Prove: $(\forall n \geq 1) \left[a_n = n^2 \right]$

Recurrence relation and summation together

- Assume the following definition of a sequence:

$$a_0 = 1$$

$$\text{For } n \geq 1 : a_n = \left[\sum_{i=0}^{n-1} a_i \right] + 1$$

- Prove: $(\forall n \geq 0) [a_n = 2^n]$

Another recurrence to consider

- Assume the following definition of a function:

$$\begin{aligned} a_0 &= 1 & a_1 &= 1 & a_2 &= 3 \\ (\forall k \in \mathbb{Z}^{\geq 3}) [a_k &= a_{k-1} + a_{k-2} - a_{k-3}] \end{aligned}$$

Claim: $(\forall n \in \mathbb{Z}^{\geq 0}) [a_n \in \mathbb{Z}^{odd}]$

Can we prove this with induction?

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What happens if we knock the first three dominos off the table manually?

Strong induction

For (ordinary) induction:

- We assume one particular domino falls and then show that the next one must also fall.

For strong induction:

- We assume all of the dominos before domino $k + 1$ fall and then show domino $k + 1$ must also fall.

This is frequently easier because our inductive hypothesis is stronger (we are assuming more stuff).

Strong induction

Definition

Claim: $(\forall n \in \mathbb{N})[P(n)]$

Proof:

I will apply strong induction on n

Base case: Show $P(0)$, $P(1)$, $P(2)$, \dots , $P(7)$ directly.

Inductive hypothesis: For some $k \geq 7$: Assume $P(i)$ holds for all $i \leq k$.

Inductive Step: Prove $P(k+1)$ must also be true, based on your assumption that P holds for all previous values.

Why Does Strong Induction Work?

$$P(1)$$

$$(\forall n \geq 2)[P(1) \wedge P(2) \wedge \cdots \wedge P(n-1) \rightarrow P(n)]$$

- 1) Have $P(1)$
- 2) Have $P(1) \rightarrow P(2)$
- 3) From $P(1)$ and $P(1) \rightarrow P(2)$ have $P(2)$.
- 4) Have $P(1) \wedge P(2) \rightarrow P(3)$.
- 5) From $P(1), P(2)$ and $P(1) \wedge P(2) \rightarrow P(3)$ have $P(3)$.
etc.

Now prove the recurrence relation property, using strong induction

- Here is the function definition again:

$$\begin{aligned} a_0 &= 1 & a_1 &= 1 & a_2 &= 3 \\ (\forall k \in \mathbb{Z}^{\geq 3}) [a_k &= a_{k-1} + a_{k-2} - a_{k-3}] \end{aligned}$$

Claim: $(\forall n \in \mathbb{Z}^{\geq 0}) [a_n \in \mathbb{Z}^{odd}]$

Another recurrence relation example

- Assume the following definition of a function:

$$a_0 = 1 \qquad a_1 = 2$$

$$(\forall k \in \mathbb{Z}^{\geq 2})[a_k = a_{k-1} + a_{k-2}]$$

- Prove the following definition property, using strong induction:

$$(\forall n \in \mathbb{Z}^{\geq 0})[a_n \leq 2^n]$$

Another example- a divisibility property

- Assume the following definition of a recurrence relation:

$$a_0 = 0$$

$$a_1 = 7$$

$$(\forall i \geq 2)[a_i = 2a_{i-1} + 3a_{i-2}]$$

- Prove using strong induction that all elements in this relation have this property:

$$(\forall n \in N)[a_n \equiv 0 \pmod{7}]$$