

# CMSC 330: Organization of Programming Languages

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Tail Recursion

# Reverse

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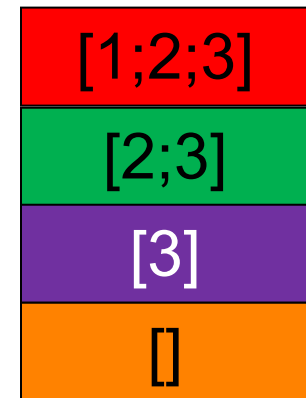
```
let rec rev l = match l with
  [] -> []
  | (x::xs) -> (rev xs) @ [x]
```

- Pushes a stack frame on each recursive call

**rev [1;2;3]**

```
→ (rev [2;3]) @ [1]
→ ((rev [3]) @ [2]) @ [1]
→ (((rev []) @ [3]) @ [2]) @ [1]
→ ((([] @ [3]) @ [2]) @ [1]
→ ([3] @ [2]) @ [1]
→ [3;2] @ [1]
→ [3;2;1]
```

Stack: values of **l**



# A Clever Version of Reverse

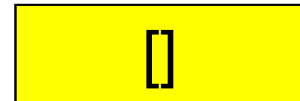
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```
let rec rev_helper l a = match l with
  [] -> a
  | (x::xs) -> rev_helper xs (x::a)
let rev l = rev_helper l []
```

- No need to push a frame for each call!

```
rev [1;2;3] →
rev_helper [1;2;3] [] →
rev_helper [2;3] [1] →
rev_helper [3] [2;1] →
rev_helper [] [3;2;1] →
[3;2;1]
```

Stack: values of `l`



# Tail Recursion

---

- Whenever a function ends with a recursive call, it is called **tail recursive**
  - Its “tail” is recursive
- Tail recursive functions can be implemented **without requiring a stack frame for each call**
  - **No intermediate variables need to be saved**, so the compiler overwrites them
- Typical pattern is to use an **accumulator** to build up the result, and return it in the base case

# Compare rev and rev\_helper

---

```
let rec rev l =  
  match l with  
  [] -> []  
  | (x::xs) -> (rev xs) @ [x]
```

*Waits for recursive call's result to compute final result*

```
let rec rev_helper l a =  
  match l with  
  [] -> a  
  | (x::xs) -> rev_helper xs (x::a)
```

*final result is the result of the recursive call*

# Exercise: Finish Tail-recursive Version

---

```
let rec sumlist l =  
  match l with  
  [] -> 0  
  | (x::xs) -> (sumlist xs) + x
```

*Tail-recursive version:*

```
let sumlist l =  
  let rec helper l a =  
    match l with  
    [] -> a  
    | (x::xs) -> helper xs (x+a) in  
  helper l 0
```

# Quiz #1

---

True/false: `map` is tail-recursive.

```
let rec map f = function
  [] -> []
| (h::t) -> (f h) :: (map f t)
```

- A. True
- B. False

# Quiz #1

---

True/false: `map` is tail-recursive.

```
let rec map f = function
  [] -> []
| (h::t) -> (f h) :: (map f t)
```

A. True

**B. False**



## Quiz #2

---

True/false: `fold_left` is tail-recursive

```
let rec fold_left f a = function
  [] -> a
| (h::t) -> fold f (f a h) t
```

- A. True
- B. False

## Quiz #2

---

True/false: `fold_left` is tail-recursive

```
let rec fold_left f a = function
  [] -> a
| (h::t) -> fold f (f a h) t
```

**A. True**

B. False

## Quiz #3

---

True/false: `fold_right` is tail-recursive

```
let rec fold_right f l a =  
  match l with  
  | [] -> a  
  | (h::t) -> f h (fold_right f t a)
```

- A. True
- B. False

## Quiz #3

---

True/false: `fold_right` is tail-recursive

```
let rec fold_right f l a =  
  match l with  
  | [] -> a  
  | (h::t) -> f h (fold_right f t a)
```

A. True

**B. False**

# Tail Recursion is Important

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- Pushing a call frame for each recursive call when operating on a list is dangerous
  - One stack frame for each list element
  - Big list = **stack overflow!**
- So: **favor tail recursion when inputs could be large** (i.e., recursion could be deep). E.g.,
  - Prefer `List.fold_left` to `List.fold_right`
    - Library documentation should indicate tail recursion, or not
  - Convert recursive functions to be tail recursive

# Tail Recursion Pattern (1 argument)

---

let *func* x =

let rec helper arg acc =

if (*base case*) then acc

else

let arg' = (*argument to recursive call*)

let acc' = (*updated accumulator*)

helper arg' acc' in (\* end of helper fun \*)

helper x (*initial val of accumulator*)

::  
;

# Tail Recursion Pattern with `fact`

---

let `fact` x =

let rec helper arg acc =

if `arg = 0` then acc

else

let arg' = `arg - 1` in

let acc' = `acc * arg` in

helper arg' acc' in (\* end of helper fun \*)

helper x `1`

;;

# Tail Recursion Pattern with `rev`

---

```
let rev x =
```

```
  let rec rev_helper arg acc =
```

```
    match arg with [] -> acc
```

```
    | h::t ->
```

```
      let arg' = t in
```

```
      let acc' = h::acc in
```

```
      rev_helper arg' acc' in (* end of helper fun *)
```

```
  rev_helper x []
```

```
;;
```

*Can generalize to more than one argument, and multiple cases for each recursive call*



# Quiz #4

---

True/false: this is a tail-recursive `map`

```
let map f l =  
  let rec helper l a =  
    match l with  
    [] -> a  
    | h::t -> helper t ((f h)::a)  
  in helper l []
```

- A. True
- B. False

# Quiz #4

---

True/false: this is a tail-recursive `map`

```
let map f l =  
  let rec helper l a =  
    match l with  
    [] -> a  
    | h::t -> helper t ((f h)::a)  
  in helper l []
```

A. True

**B. False** (elements are reversed)

# A Tail Recursive map

---

```
let map f l =  
  let rec helper l a =  
    match l with  
    [] -> a  
    | h::t -> helper t ((f h)::a)  
  in rev (helper l [])
```

Could instead change  $(f\ h)::a$  to be  $a@(f\ h)$

**Q:** Why is the above implementation a better choice?

**A:**  $O(n)$  running time, not  $O(n^2)$  (where  $n$  is length of list)

# Outlook: Is Tail Recursion General?

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- A function that is tail-recursive returns **at most once** (to its caller) when completely finished
  - The final result is exactly the result of a recursive call; no stack frame needed to remember the current call
- Is it possible to convert an *arbitrary program* into an equivalent one, except where **no call ever returns?**
  - Yes. This is called **continuation-passing style**
  - We will look at this later, if we have time