# Motion planning: Beyond Navmeshes 

CMSC425.01 Spring 2019

## Administrivia

- Exam being graded
- Project 2 b concepts out, write up soon (add animations to 2 a)


## Today's questions

Big question: Making intelligent agents First question: Navigation

## Finding paths in polygonal configuration space

- Version 1: Navmesh
- Others?
- Version 7: Randomized placement (sampling)

(a)

(b)

(c)


## Finding paths in polygonal configuration space

- Version 1: Navmesh
- Others?
- Version 8: Rapidly-expanded Random Trees (RRTs)

(c)

(d)


## Computing shortest path

- Reduce navigation to path finding in graphs
- Directed?
- Weighted?

- $G=(V, E)$
- Vertices $V=\{u, v, \ldots\}$
- Edges $E=\{(u, v), \ldots\}$
- Weight function $\mathrm{w}(u, v) \rightarrow$ reals


## Computing shortest path

- Reduce navigation to path finding in graphs
- Directed?
- Weighted?

- Path sequence of nodes

$$
\text { - } P=\left\langle u_{0}, u_{1}, \ldots, u_{k}\right\rangle
$$

- Path cost

$$
\cdot \operatorname{cost}(P)=\sum_{i=0}^{k} w\left(u_{i}, u_{i+1}\right)
$$

- Lowest cost path $\partial(s, t)$


## First: what's the problem?

- Compute one shortest path?
- Compute all shortest paths to store?


## First: what's the problem?

- Compute path here to there?
- Find fastest way to home base?
- Reverse edges
- Find shortest path to all from home
- Find closest facility (health, etc)?
- Add Supernode connected to all facilities.
- Compute all shortest paths to store?
- Floyd-Warshall



## First: what's the problem?

- Find closest facility (health, etc)?
- Add Supernode connected to all facilities.



## Uninformed vs. informed search

- Uninformed - follow weights
- Pick next node on distance to $\mathrm{d}[\mathrm{u}]$
- Informed - add bias towards destination
- Heuristic
- Pick next node on distance to goal h(u)


## Uninformed search



Informed search


## Informed search

- Distance functions
- w(u,v) - distance node u to v
- d[u] - distance traversed from start to node u
- $\operatorname{dist}(\mathrm{u}, \mathrm{t})$ - distance from u to t
- $w(s, 1)=$ $\qquad$ $\operatorname{dist}(1, t)=$ $\qquad$
- $w(s, 2)=$ $\qquad$ $\operatorname{dist}(1, t)=$ $\qquad$



## Informed search

- Distance functions
- w(u,v) - distance node u to v
- d[u] - distance traversed from start to node u
- dist(u,t) - distance from $u$ to $t$
- $w(s, 1)=3 \quad \operatorname{dist}(1, t)=6$
- $w(s, 2)=3 \quad \operatorname{dist}(1, t)=4$
- $\operatorname{dist}(\mathrm{u}, \mathrm{t})$ is a heuristic



## Less perfect information?

- Can't see rest of graph until you expand it
- Need guess on what's to come
- dist( $u, t)$ as Euclidean distance
- Approximates actual cost



## Footnote

- Euclidean distance
- distE(p1,p2) = $\operatorname{sqrt}\left((x 1-x 2)^{\wedge} 2+(y 1-y 2)^{\wedge} 2\right)$
- Manhattan distance
- $\operatorname{dist} \mathrm{M}(\mathrm{p} 1, \mathrm{p} 2)=$

$$
a b s(x 1-x 2)+a b s(y 1-y 2)
$$



```
Dijkstra(G, s, t) {
    foreach (node u) { // initialize
            d[u] = +infinity; mark u undiscovered
    }
    d[s] = 0; mark s discovered // distance to source is 0
    repeat forever { // go until finding t
            let u be the discovered node that minimizes d[u]
            if (u == t) return d[t] // arrived at the destination
            else {
            for (each unfinished node v adjacent to u) {
                d[v] = min(d[v], d[u] + w(u,v)) // update d[v]
                mark v discovered
            }
            mark u finished // we're done with u
    }
}
```


## Example

- w(u,v) as given
- Start with d array as

- End with?



## Example

- w(u,v) as given
- Start with d array as

- End with?

| a | b | c | d | e | $z$ |
| :---: | :---: | :---: | :---: | :---: | :---: |
| 0 | 3 | 4 | 7 | 5 | 1 |



```
BestFirst(G, s, t) {
    foreach (node u) { // initialize
        d[u] = +infinity; mark u undiscovered
    }
    d[s] = 0; mark s discovered // distance to source is 0
    repeat forever { // go until finding t
            let u be the discovered node that minimizes dist(u,t)
            if (u == t) return d[t] // arrived at the destination
            else {
            for (each unfinished node v adjacent to u) {
                d[v] = min(d[v], d[u] + w(u,v)) // update d[v]
                mark v discovered
            }
            mark u finished // we're done with u
    }
}
```


## Best first bad case ...

- Trapped in local minimum


A*

- Pick next node to expand based on sum of distance so far and heuristic

$$
f(u)=d[u]+h(u)=d[u]+\operatorname{dist}(u, t)
$$

```
A-Star(G, s, t) {
    foreach (node u) { // initialize
        d[u] = +infinity; mark u undiscovered
    }
    d[s] = 0; mark s discovered // distance to source is 0
    repeat forever { // go until finding t
        let u be the discovered node that minimizes d[u] + dist(u,t)
        if (u == t) return d[t] // arrived at the destination
        else {
            for (each unfinished node v adjacent to u) {
                d[v] = min(d[v], d[u] + w(u,v)) // update d[v]
                mark v discovered
            }
            mark u finished // we're done with u
    }
}
```


## A* Example

- Manhattan distance


| A* Search - Each entry is $d[u]: f(u)^{\text {a }}$ |  |  |  |  |  |  |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| Stage | $d[s]$ | $d[a]$ | $d[b]$ | $d[c]$ | $d[d]$ | $d[e]$ | $d[f]$ | $d[g]$ | $d[h]$ | $d[t]$ |
| $h(u)$ | 15 | 13 | 15 | 17 | 12 | 10 | 9 | 8 | 5 | 0 |
| Init | 0:15 | $\infty: 13$ | $\infty: 15$ | $\infty: 17$ | $\infty: 12$ | $\infty: 10$ | $\infty$ :9 | $\infty$ :8 | $\infty$ : 5 | $\infty$ :0 |
| 1: $s$ | 0 | 8:13 | - | 2:17 | 3:12 | - | - | - | - | - |
| 2: d | $\downarrow$ | 8:13 | - | 2:17 | 3 | 5:10 | 6:9 | - | - | - |
| 3: e |  | 8:13 | - | 2:17 | $\downarrow$ | 5 | 6:9 | 7:8 | - | - |
| 4: $f$ |  | 8:13 | - | 2:17 |  | $\downarrow$ | 6 | 7:8 | - | 15:0 |
| 5: $t$ |  | 8:13 | - | 2:17 |  |  | $\downarrow$ | 7:8 | - | 15 |
| Final | 0 | 8 | $\infty$ | 2 | 3 | 5 | 6 | 7 | $\infty$ | 15 |

## Good heuristics

- For $\mathrm{A}^{*}$ to compute correctly the heuristic $\mathrm{h}(\mathrm{u})$ must be:
- Admissible:
$\mathrm{h}(\mathrm{u})$ never overestimates the graph distance from node $u$ to goal $t$
- Consistent:

$$
h\left(u^{\prime}\right)<=\operatorname{delta}\left(u^{\prime}, u^{\prime \prime}\right)+h\left(u^{\prime \prime}\right)
$$

