CMSC 330: Organization of Programming Languages

Lets, Tuples, Records
Let Expressions

• Enable binding variables in other expressions
  – These are different from the let definitions we’ve been using at the top-level

• They are expressions, so they have a value

• Syntax
  – `let x = e1 in e2`
  – `x` is a bound variable
  – `e1` is the binding expression
  – `e2` is the body expression
Let Expressions

• Syntax
  - let \( x = e_1 \) in \( e_2 \)

• Evaluation
  - Evaluate \( e_1 \) to \( v_1 \)
  - Substitute \( v_1 \) for \( x \) in \( e_2 \) yielding new expression \( e_2' \)
  - Evaluate \( e_2' \) to \( v_2 \)
  - Result of evaluation is \( v_2 \)

Example

```
let x = 3+4 in 3*x
let x = 7 in 3*x
3*7
21
```
Let Expression Example

let x = 3+27 in x*3

- 3+27 : int
- x*3 : int (assuming x:int)
- so let x = 3+27 in x*3 : int
Let Definitions vs. Let Expressions

• At the top-level, we write
  – `let x = e ;; (* no in e2 part *)`
  – This is called a let definition, not a let expression
    • Because it doesn’t, itself, evaluate to anything

• Omitting `in` means “from now on”:
  # `let pi = 3.14 ;;`
  (* pi is now bound in the rest of the top-level scope *)
Top-level expressions

• We can write any expression at top-level, too
  – `e;;`
  – This says to evaluate `e` and then ignore the result
    • Equivalent to `let _ = e;;`
    • Useful when `e` has a side effect, such as reading/writing a file, printing to the screen, etc.

```ml
let x = 37;;
let y = x + 5;;
print_int y;;
print_string "\n";;
```

• When run, outputs 42 to the screen
Let Expressions: Scope

- In `let x = e1 in e2`, variable `x` is *not* visible outside of `e2`.

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;
print_float pi;;
```

- `bind pi (only) in body of let` (which is `pi *. 3.0 *. 3.0`)
- error: `pi` not bound
Binding in other languages

• Compare to similar usage in Java/C

```ocaml
let pi = 3.14 in
  pi *. 3.0 *. 3.0;;
pi;; (* pi unbound! *)
```

```c
{float pi = 3.14;
  pi * 3.0 * 3.0;
}
pi; /* pi unbound! */
```
Examples – Scope of Let bindings

• \texttt{x;;}
  – (* Unbound value x *)

• \texttt{let x = 1 in x + 1;;}
  – (* 2 *)

• \texttt{let x = x in x + 1;;}
  – (* Unbound value x *)
Examples – Scope of Let bindings

• let x = 1 in (x + 1 + x) ;;
  – (* 3 *)

• (let x = 1 in x + 1) ;; x;;
  – (* Unbound value x *)

• let x = 4 in (let x = x + 1 in x) ;;
  – (* 5 *)
Shadowing Names

- **Shadowing** is rebinding a name in an inner scope to have a different meaning
  - May or may not be allowed by the language

```c
int i;
void f(float i) {
    { char *i = NULL;
      ...
    }
}
```

```java
void h(int i) {
    {
        float i; // not allowed
        ...
    }
}
```

```ocaml
let x = 3;;
let g x = x + 3;;
```
Shadowing, by the Semantics

- What if $e_2$ is also a `let` for $x$?
  - Substitution will **stop** at the $e_2$ of a shadowing $x$

Example

```plaintext
let x = 3+4 in let x = 3*x in x+1
  ➢ let x = 7 in let x = 3*x in x+1
  ➢ let x = 3*7 in x+1
  ➢ let x = 21 in x+1
  ➢ 21+1
  ➢ 22
```

Not substituted, since it is shadowed by the inner `let`
Let Expressions in Functions

• You can use `let` inside of functions for local vars

```ml
let area r =
    let pi = 3.14 in
    pi *. r *. r
```

– And you can use many `lets` in sequence

```ml
let area d =
    let pi = 3.14 in
    let r = d /. 2.0 in
    pi *. r *. r
```
Shadowing (of Locals) Discouraged

• You can use shadowing to simulate mutation (variable update)

```ocaml
let rec f x n =  
  if x = 0 then 1  
  else  
    let x = x - 1 in (* shadowed *)  
    n * (f x n)
```

• But avoiding shadowing can be clearer, so we recommend not using it
  – With no shadowing, if you see a variable x, you know it hasn’t been ”changed,” no matter where it appears
  – if you want to “update” n, use a new name n1, n’, etc.
Nested Let Expressions

- Uses of `let` can be nested in OCaml
  - Nested bound variables (\(\pi\) and \(r\)) invisible outside

- Similar scoping possibilities C and Java

```ocaml
let res =
    (let area =
        (let pi = 3.14 in
         let r = 3.0 in
         pi *. r *. r)
        in
         area /. 2.0));;
```

```c
float res;
{ float area;
    { float pi = 3.14
      float r = 3.0;
      area = pi * r *
      r;
    }
    res = area / 2.0;
}
```
Nested Let Style: Generally Avoid

- Oftentimes a nested binding can be rewritten in a more linear style
  - Easier to understand
- Can go too far: namespace pollution
  - Avoiding adding unnecessary variable bindings to top-level

```ml
let res =
  (let area =
    (let pi = 3.14 in
     let r = 3.0 in
     pi *. r *. r)
   in
    area /. 2.0);

let pi = 3.14;;
let r = 3.0;;
let area = pi *. r *. r;;
let res = area /. 2.0;;
```
Quiz 1

Which of these is **not** an expression that evaluates to 3?

A. let x=3  
B. let x=2 in x+1  
C. let x=3 in x  
D. 3
Quiz 1

Which of these is not an expression that evaluates to 3?

A. `let x=3` --> not an expression
B. `let x=2 in x+1`
C. `let x=3 in x`
D. 3
Quiz 2: What does this evaluate to?

```plaintext
let x = 2 in
x = 3
```

A. 3
B. 2
C. true
D. false
Quiz 2: What does this evaluate to?

```
let x = 2 in
x = 3
```

A. 3  
B. 2  
C. true  
D. false
Quiz 3: What does this evaluate to?

```
let x = 3 in
let y = x+2 in
let x = 8 in
x+y
```

A. 13  
B. 8  
C. 11  
D. 18
Quiz 3: What does this evaluate to?

```
let x = 3 in
let y = x+2 in
let x = 8 in
x+y
```

A. 13
B. 8
C. 11
D. 18
**let Specializes match**

More general form of let allows patterns:

- **let** \( p = e_1 \) in \( e_2 \)
  
  where \( p \) is a pattern. If \( e_1 \) fails to match that pattern then an exception is thrown

This pattern form of **let** is equivalent to

- **match** \( e_1 \) with \( p \rightarrow e_2 \)

Examples

- **let** \([x] = [1]\) in \(1::x\)  (* evaluates to \([1;1]\) *)
- **let** \(h::\_ = [1;2;3]\) in \(h\)  (* evaluates to 1 *)
- **let** () = **print_int** 5 in 3  (* evaluates to 3 *)
Tuples

• Constructed using \((e_1, \ldots, e_n)\)
• Deconstructed using pattern matching
  – Patterns involve parens and commas, e.g., \((p_1, p_2, \ldots)\)
• Tuples are similar to C structs
  – But without field labels
  – Allocated on the heap
• Tuples can be heterogeneous
  – Unlike lists, which must be homogenous
  – \((1, ["string1";"string2"]))\) is a valid tuple
Tuple Types

• Tuple types use * to separate components
  – Type joins types of its components

• Examples
  – \((1, 2)\) :
  – \((1, "string", 3.5)\) :
  – \((1, ["a"; "b"], 'c')\) :
  – \([(1,2)]\) :
  – \([(1, 2); (3, 4)]\) :
  – \([(1,2); (1,2,3)]\) :
Tuple Types

• Tuple types use * to separate components
  – Type joins types of its components

• Examples
  – (1, 2) : int * int
  – (1, "string", 3.5) : int * string * float
  – (1, ["a"; "b"], 'c') : int * string list * char
  – [(1,2)] : (int * int) list
  – [(1, 2); (3, 4)] : (int * int) list
  – [(1,2); (1,2,3)] : error

Because the first list element has type int * int, but the second has type int * int * int – list elements must all be of the same type
Pattern Matching Tuples

# let plusThree t =
  match t with
    (x, y, z) -> x + y + z;;
plusThree : int*int*int -> int = <fun>

# let plusThree' (x, y, z) = x + y + z;;
plusThree' : int*int*int -> int = <fun>

# let addOne (x, y, z) = (x+1, y+1, z+1);;
addOne : int*int*int -> int*int*int = <fun>

# plusThree (addOne (3, 4, 5));;
  : int = 15

Remember, **semicolon** for lists, **comma** for tuples

- [1, 2] = [(1, 2)] which is a list of size one
- (1; 2) Warning: This expression should have type unit
Tuples Are A Fixed Size

• This OCaml definition
  
  ```ocaml
  # let foo x = match x with
  (a, b) -> a + b
  | (a, b, c) -> a + b + c;;
  ```

• Tuples of different size have different types

  - (a, b) has type: 'a * 'b * 'c
  - (a, b, c) has type: 'a * 'b
Records

• Records: identify elements by name
  – Elements of a tuple are identified by position

• Define a record type before defining record values

  ```haskell
  type date = { month: string; day: int; year: int };
  ```

• Define a record value

  ```haskell
  # let today = { day=16; year=2017; month="f"^"eb" };;
  today : date = { day=16; year=2017; month="feb" };;
  ```
Destructing Records

```ocaml
type date = { month: string; day: int; year: int } let today = { day=16; year=2017; month="feb" };;
```

- **Access by field name or pattern matching**

```ocaml
print_string today.month;; (* prints feb *)
(* patterns *)
let { month=_; day=d } = today in
let { year } = today in
let _ = print_int d in  (* prints 16 *)
print_int year;;  (* prints 2017 *)
```

- **Notes:**
  - In record patterns, you can skip or reorder fields
  - You can use the field name as the bound variable
Quiz 4: What does this evaluate to?

let get (a,b) = a+b in
get 1 2

A. 3
B. 2
C. 1
D. type error
Quiz 4: What does this evaluate to?

```haskell
let get (a,b) = a+b in
get 1 2
```

A. 3
B. 2
C. 1
D. type error – `get` takes one argument (a pair)
Quiz 5: What does this evaluate to?

let get x y =
  match x with
    (a,b) -> a+y
in
get (1,2) 1

A. 3
B. type error
C. 2
D. 1
Quiz 5: What does this evaluate to?

```ocaml
let get x y =
    match x with
    (a,b) -> a+y
in
get (1,2) 1
```

A. 3
B. type error
C. 2
D. 1
Quiz 6: What is the type of \texttt{shift}?

\begin{verbatim}
type point = {x:int; y:int}
let shift p =
    match p with
    { x=px; y=py } -> [px;py]
\end{verbatim}

A. \texttt{point -> int list}
B. \texttt{int list -> int list}
C. \texttt{point -> point}
D. \texttt{point -> bool list}
Quiz 6: What is the type of `shift`?

```plaintext
type point = {x:int; y:int}

let shift p =
    match p with
    { x=px; y=py } -> [px;py]
```

A. point -> int list
B. int list -> int list
C. point -> point
D. point -> bool list