CMSC 430, Feb 18th 2020

Extort

• Assignment #2

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- Two issues:
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 - Without ELMS/Canvas, how can I best communicate grades?

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• Why?

• Recap from last time:

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(add1 #f)

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- Previously, these were undefined
 - In our interpreter we would get a failure because of the errors from the underlying Racket execution
 - In our compiler we'd get junk

Extort

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• Our language **extort** the same as **dupe** *except* we address errors explicitly

• No changes:

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• e ::= ... | if e e e | zero? e

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• Why don't we need to change the AST?

• Type mismatches in **dupe** were undefined behavior

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 - o Do we have to make them defined?
 - \circ What are the pros/cons?

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Errors Rule

- Let's add some, knowing that it's not strictly necessary
- Our semantics now relate programs to *answers* instead of values
 - *answers* are either *values* (as before), or *errors*
- We'll just show the new rules, none of the others have changed.

• Where can errors occur (currently)?

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E[(add 1 b), err]]

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E[(add 1 b), err]]

 $E[(sub1 \ b), err]]$

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E[(zero? b), err]]

• Is that it?

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(if (zero? #f) 1 2)

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 $E[\![e, err]\!]$

E[(zero? e), err]]

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 $\frac{E[e, err]}{E[(zero? e), err]]}$ $\frac{E[e, err]]}{E[(add1 e), err]]}$

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E[[e, err]] E[(zero? e), err]] E[[e, err]] E[(add1 e), err]] E[[e, err]] E[(sub1 e), err]]

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E[[*e*, err]] E[(zero? e), err]]*E*[[*e*, err]] E[(add 1 e), err]]*E*[[*e*, err]] E[(sub1 e), err]]E[e, err]E[[(if $e e_0 e_1$), err]]

Rules are the easy part

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• How can our implementations match these rules?

Let's look at the interpreter

We'll do that in the terminal, as it's starting to get a bit too cumbersome

Let's experiment

extort> (require "extort_interp.rkt")

Now the compiler.

Now the compiler.

• What needs to change, if anything?

Now the compiler.

- What needs to change, if anything?
- What should the error message be?

• Things need to happen in the RTS and compiler.

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 Runtime system?

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 Runtime system?
 - Compiler?

Let's take a look at the RTS and compiler

Assignment 3

• Is live

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- Due next Tuesday.

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 - Please tell your fellow students to check the webpage periodically
 - If there are any issues that might make you unable to do the assignment on time, *talk to me*