

CMSC 132: OBJECT-ORIENTED PROGRAMMING II



Algorithm Strategies

Department of Computer Science
University of Maryland, College Park

General Concepts

- **Algorithm strategy**
 - Approach to solving a problem
 - May combine several approaches
- **Algorithm structure**
 - Iterative → execute action in loop
 - Recursive → reapply action to subproblem(s)
- **Problem type**

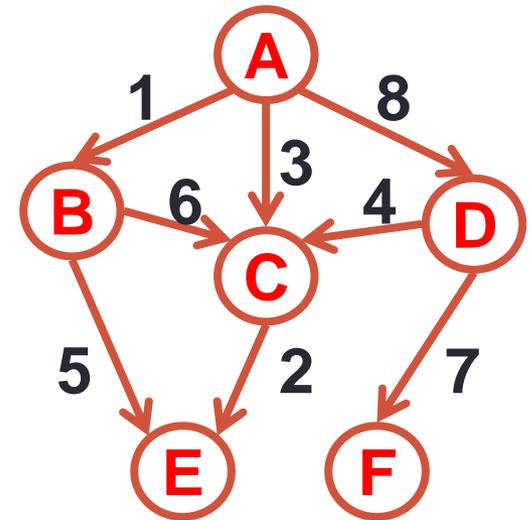
Problem Type

- **Satisfying**

- Find any satisfactory solution
- Example → Find path from **A** to **E**

- **Optimization**

- Find **best** solution (vs. cost metric)
- Example → Find **shortest** path from **A** to **E**



Some Algorithm Strategies

- Recursive algorithms
- Backtracking algorithms
- Divide and conquer algorithms
- Dynamic programming algorithms
- Greedy algorithms
- Brute force algorithms
- Branch and bound algorithms
- Heuristic algorithms

Recursive Algorithm

- Based on reapplying algorithm to subproblem
- Approach
 1. Solves **base case(s)** directly
 2. Recurs with a simpler subproblem
 3. May need to combine solution(s) to subproblems

Backtracking Algorithm

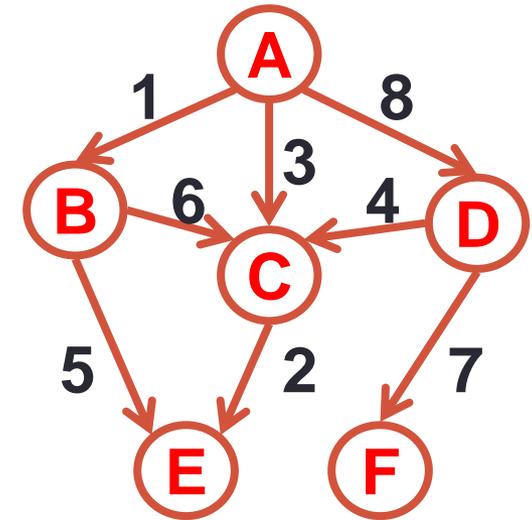
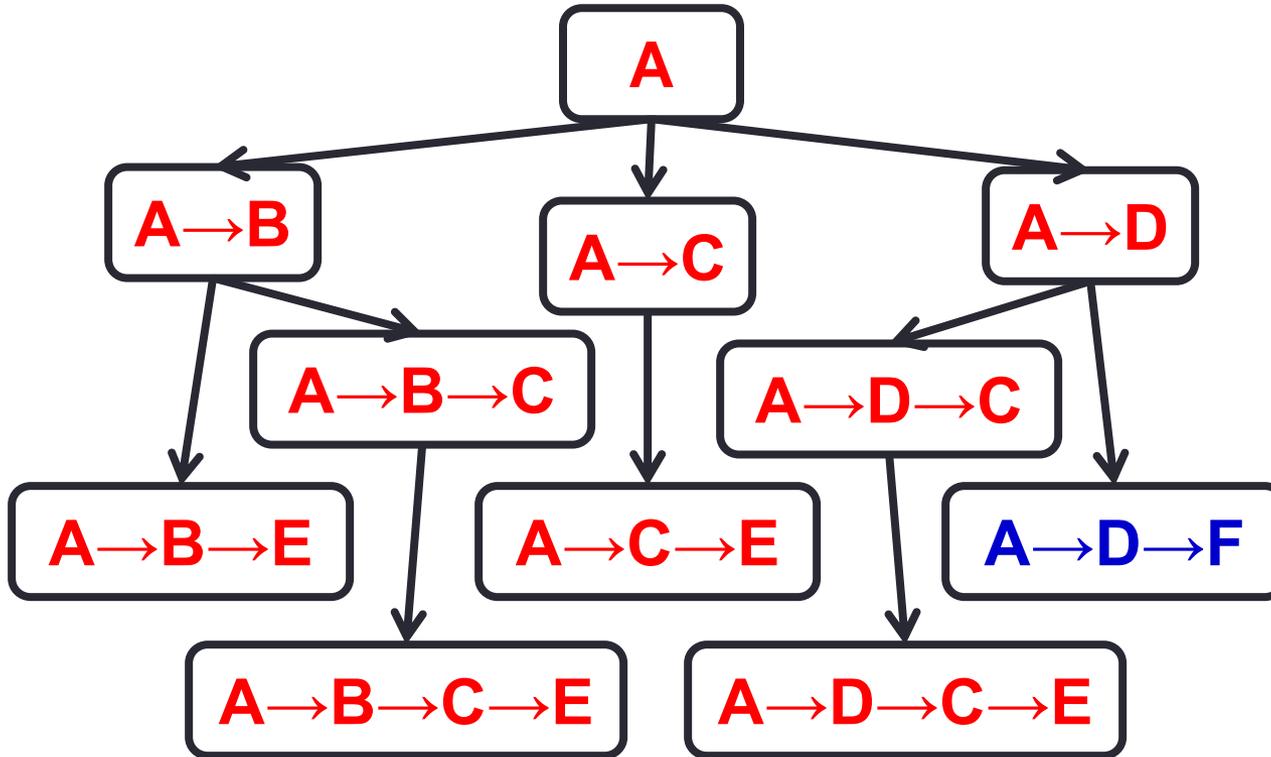
- Based on **depth-first** recursive search
- Approach
 1. Tests whether solution has been found
 2. If found solution, return it
 3. Else for each choice that can be made
 - a. Make that choice
 - b. Recur
 - c. If recursion returns a solution, return it
 4. If no choices remain, return failure
- Tree of alternatives → **search tree**

Backtracking Algorithm - Reachability

- Find path in graph from **A** to **F**
 1. Start with currentNode = A
 2. If currentNode has edge to F, return path
 3. Else select neighbor node X for currentNode
 - Recursively find path from X to F
 - If path found, return path
 - Else repeat for different X
 - Return false if no path from any neighbor X

Backtracking Algorithm – Path Finding

- Search tree (path A to F)



Backtracking Algorithm – Map Coloring

- Color a map using **four** colors so adjacent regions do not share the same color
- Coloring map of countries
 - If all countries have been colored return success
 - Else for each color c of four colors and country n
 - If country n is not adjacent to a country that has been colored c
 - Color country n with color c
 - Recursively color country $n+1$
 - If successful, return success
 - Return failure
- Map from Wikipedia
 - https://en.wikipedia.org/wiki/Four_color_theorem#/media/File:Map_of_United_States_vivid_colors_shown.png

Divide and Conquer

- Based on dividing problem into subproblems
- Approach
 1. Divide problem into smaller subproblems
 - a. Subproblems **must be of same type**
 - b. Subproblems **do not need to overlap**
 2. Solve each subproblem recursively
 3. Combine solutions to solve original problem
- Usually contains two or more recursive calls

Divide and Conquer – Sorting

- **Quicksort**

- Partition array into **two parts** around pivot
- Recursively quicksort each part of array
- Concatenate solutions

- **Mergesort**

- Partition array into **two parts**
- Recursively mergesort each half
- Merge two sorted arrays into single sorted array

Dynamic Programming Algorithm

- **Based on remembering past results**
- Approach
 1. Divide problem into smaller subproblems
 - Subproblems **must be of same type**
 - Subproblems **must overlap**
 2. Solve each subproblem recursively
 - May simply look up solution (if previously solved)
 3. Combine solutions to solve original problem
 4. Store solution to problem
- Generally applied to optimization problems

Fibonacci Algorithm

- Fibonacci numbers
 - $\text{fibonacci}(0) = 1$
 - $\text{fibonacci}(1) = 1$
 - $\text{fibonacci}(n) = \text{fibonacci}(n-1) + \text{fibonacci}(n-2)$
- Recursive algorithm to calculate $\text{fibonacci}(n)$
 - If n is 0 or 1, return 1
 - Else compute $\text{fibonacci}(n-1)$ and $\text{fibonacci}(n-2)$
 - Return their sum
- Simple algorithm \rightarrow exponential time $O(2^n)$

Dynamic Programming – Fibonacci

- Dynamic programming version of fibonacci(n)
 - If n is 0 or 1, return 1
 - Else solve fibonacci(n-1) and fibonacci(n-2)
 - Look up value if previously computed
 - Else recursively compute
 - Find their sum and store
 - Return result
- Dynamic programming algorithm $\rightarrow O(n)$ time
 - Since solving fibonacci(n-2) is just looking up value

Dynamic Programming – Shortest Path

Dijkstra's Shortest Path Algorithm

$S = \emptyset$

$C[X] = 0$

$C[Y] = \infty$ for all other nodes

while (not all nodes in S)

 find node K not in S with smallest $C[K]$

 add K to S

 for each node M not in S adjacent to K

$C[M] = \min (C[M], C[K] + \text{cost of } (K,M))$

**Stores results of
smaller subproblems**

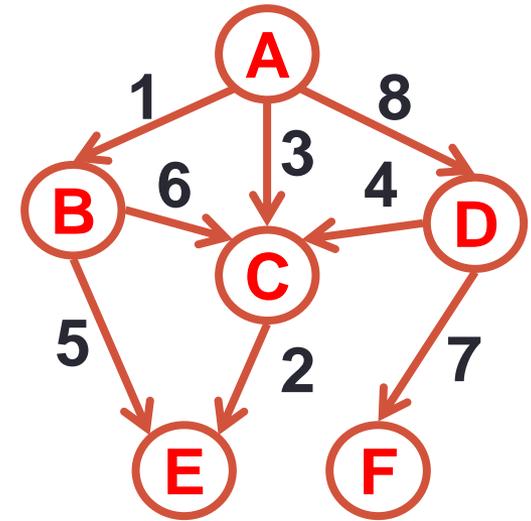
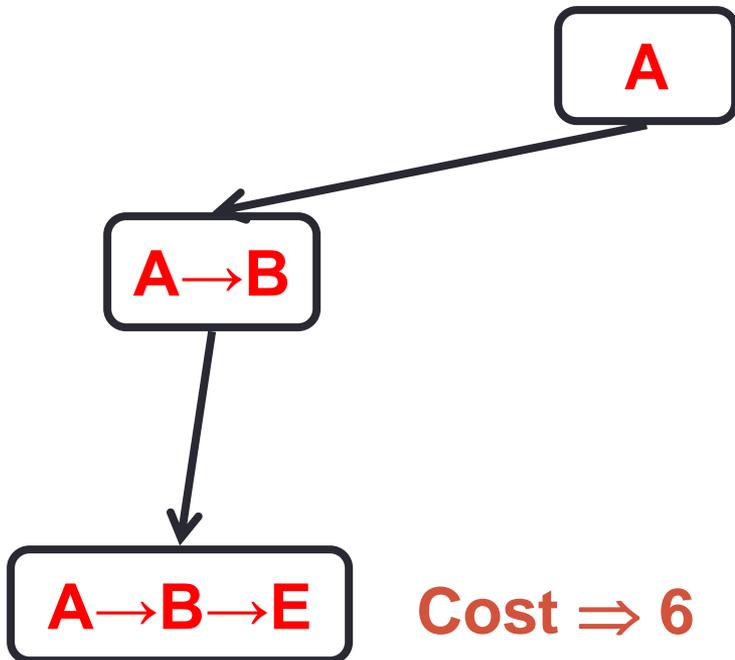


Greedy Algorithm

- Based on trying best current (local) choice
- Approach
 - At each step of algorithm
 - Choose best local solution
- Avoid backtracking, exponential time $O(2^n)$
- Hope local optimum lead to **global** optimum
- Example: Coin System
 - Coins – 30 20 15 1
 - Find minimum number of coins for 40
 - Greedy algorithm fails

Greedy Algorithm – Shortest Path

- Example (Shortest Path from A to E)
 - Choose lowest-cost neighbor



- Does not yield overall (global) shortest path

Greedy Algorithm – MST

Kruskal's Minimal Spanning Tree Algorithm

sort edges by weight (from least to most)

tree = \emptyset

for each edge (X,Y) **in order**

if it does not create a cycle

add (X,Y) to tree

stop when tree has $N-1$ edges



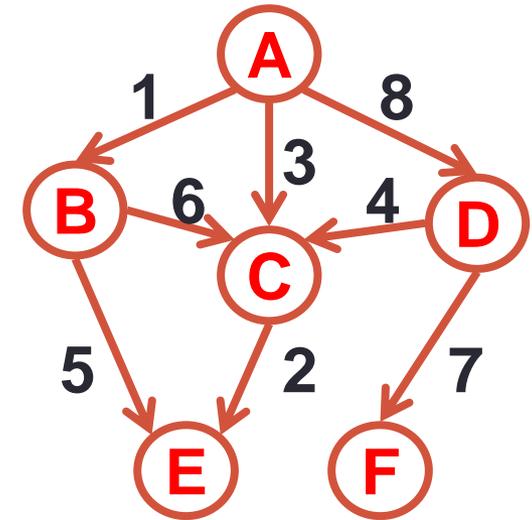
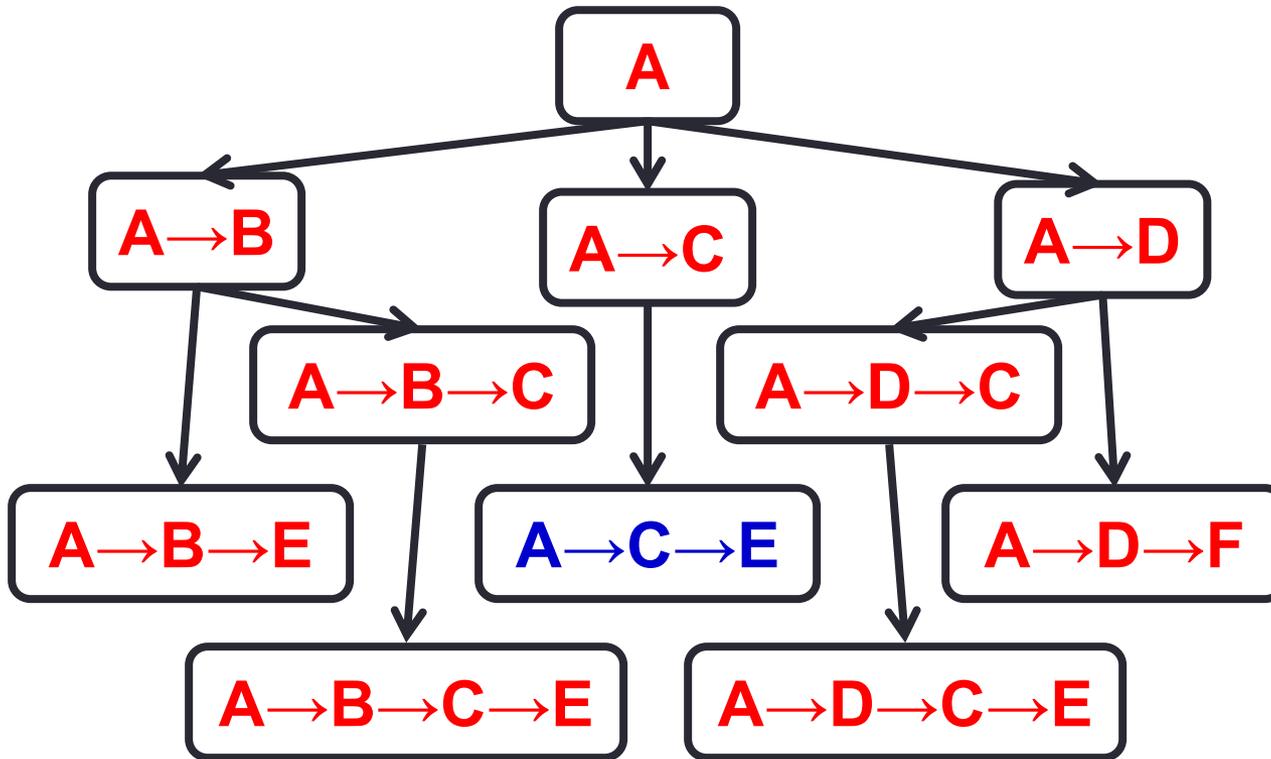
**Picks best
local solution
at each step**

Brute Force Algorithm

- **Based on trying all possible solutions**
- Approach
 - Generate and evaluate possible solutions until
 - Satisfactory solution is found
 - Best solution is found (if can be determined)
 - All possible solutions found
 - Return best solution
 - Return failure if no satisfactory solution
- Generally, most expensive approach

Brute Force Algorithm – Shortest Path

- Example (From A to E)



- Examines all paths in the graph

Brute Force Algorithm – TSP

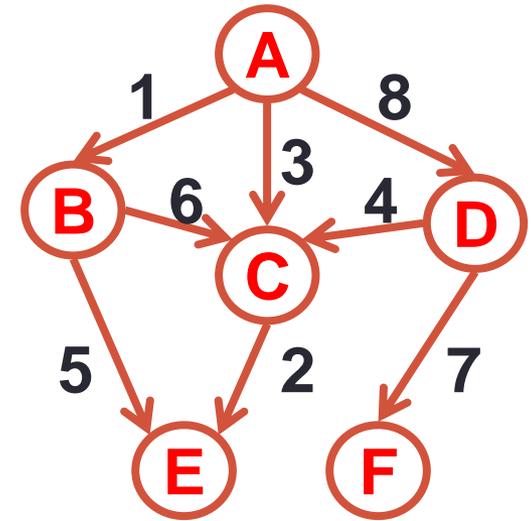
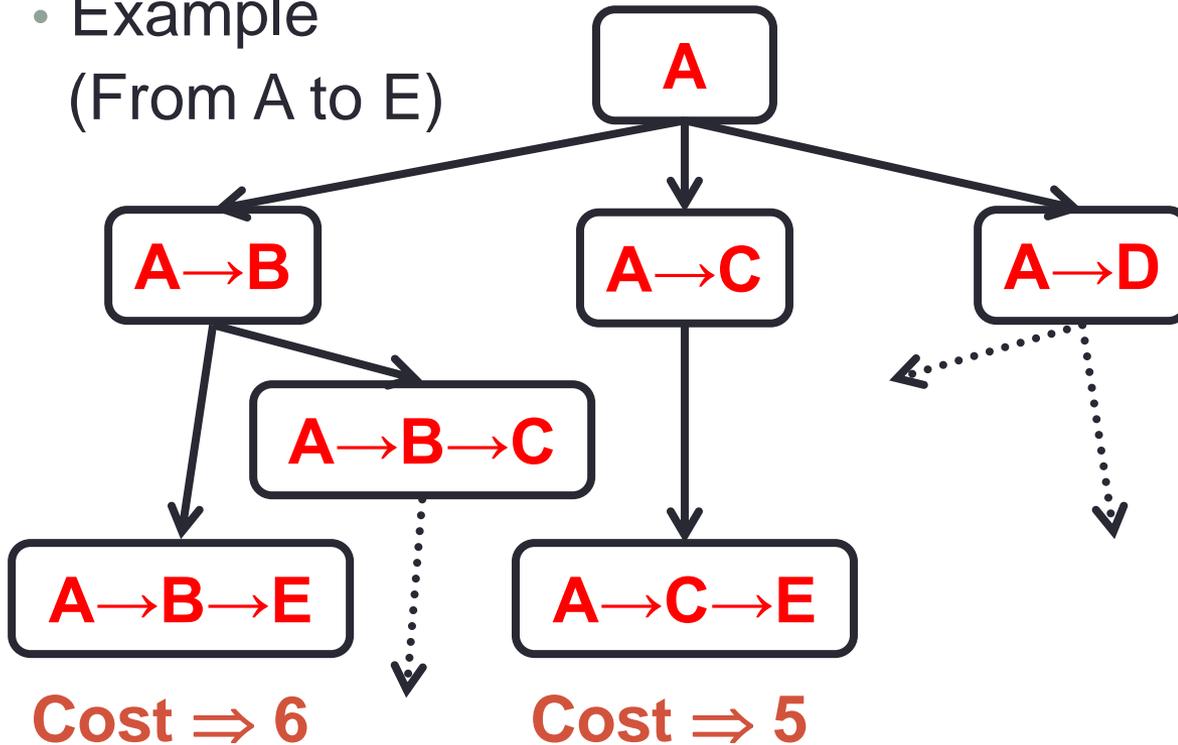
- Traveling Salesman Problem (TSP)
 - Given weighted undirected graph (map of cities)
 - Find lowest cost path visiting all nodes (cities) once
 - No known polynomial-time general solution
- Brute force approach
 - Find all possible paths using recursive backtracking
 - Calculate cost of each path
 - Return lowest cost path
 - Complexity $O(n!)$

Branch and Bound Algorithm

- Based on limiting search using current solution
- Approach
 - Track best current solution found
 - Eliminate (**prune**) partial solutions that can not improve upon best current solution
- Reduces amount of backtracking
 - Not guaranteed to avoid exponential time $O(2^n)$

Branch & Bound Alg. – Shortest Path

- Example
(From A to E)



- Starting with $A \rightarrow B \rightarrow E$
- Pruned paths beginning with $A \rightarrow B \rightarrow C$ & $A \rightarrow D$

Branch and Bound – TSP

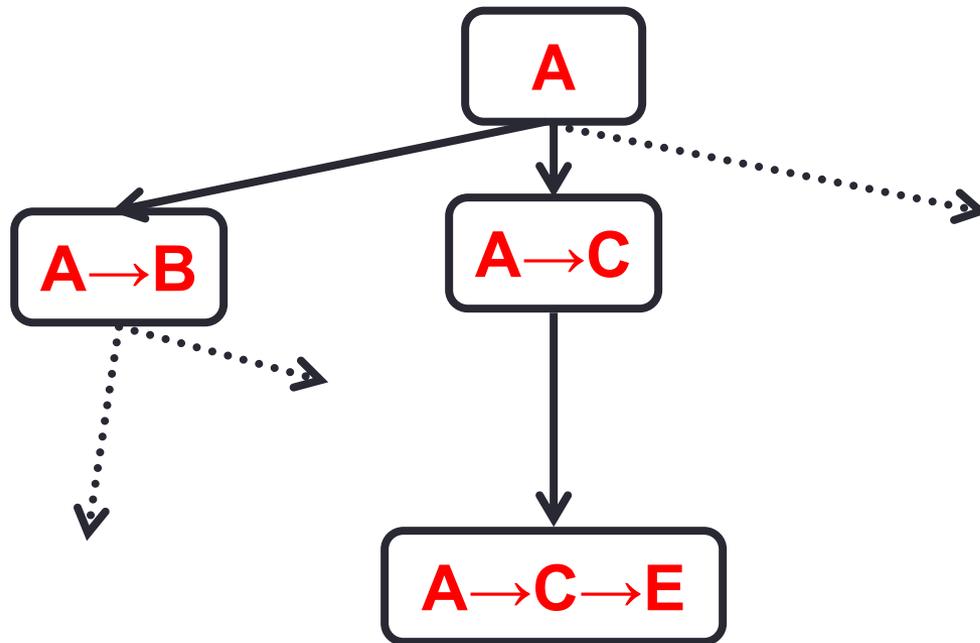
- Branch and bound algorithm for TSP
 - Find possible paths using recursive backtracking
 - Track cost of best current solution found
 - Stop searching path **if cost > best current solution**
 - Return lowest cost path
- If good solution found early, can reduce search
- May still require exponential time $O(2^n)$

Heuristic Algorithm

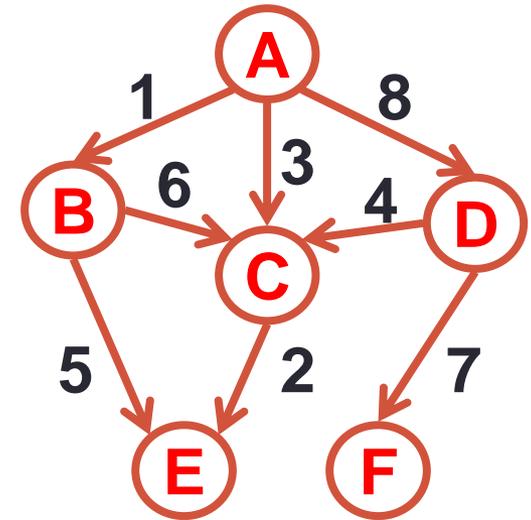
- **Based on trying to guide search for solution**
- Heuristic \Rightarrow “rule of thumb”
- Approach
 - Generate and evaluate possible solutions
 - Using “rule of thumb”
 - Stop if satisfactory solution is found
- Can reduce complexity
- Not guaranteed to yield best solution

Heuristic – Shortest Path

- Example (From A to E)
 - Try only edges with cost < 5



- Worked...in this case



Heuristic Algorithm – TSP

- Heuristic algorithm for TSP
 - Find possible paths using recursive backtracking
 - Search 2 lowest cost edges at each node first
 - Calculate cost of each path
 - Return lowest cost path from first 100 solutions
- Not guaranteed to find best solution
- Heuristics used frequently in real applications

Summary

- Wide range of strategies
- Choice depends on
 - Properties of problem
 - Expected problem size
 - Available resources