# CMSC 330: Organization of Programming Languages

#### Closures and Iterators In Rust

CMSC 330 - Spring 2021

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## **Closures**

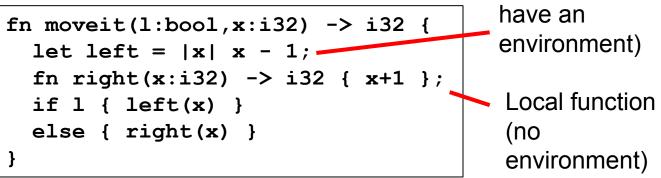
- Syntax
  - |x1[:t1]<sup>?</sup>, ..., xn[:tn]<sup>?</sup>| [-> u]<sup>?</sup> e
    - Type annotations are optional will be inferred if absent
- Evaluation
  - A closure is a value
- Type checking
  - has type (t1, ..., tn) -> u

when **e** : **u** under assumptions **x1** : **t1**, ..., **xn** : **tn** 

Not curried

## **Using Closures/Functions Locally**

Rust has local functions, and closures



Closure (may

OCaml local functions/closures

```
let moveit l x =
  let left = fun x -> x - 1 in
  let right = fun x -> x + 1 in
  if l then left x
  else right x
```

## Limits of Type Inference

Rust infers non-polymorphic types

OCaml infers polymorphic types

```
let f = fun x -> x in (* 'a -> 'a *)
let x = id 1 in
let y = id "hi" in (* OK *) ...
```

- More details on closures at the end, including polymorphism
  - Now for something (not so completely) different

#### Iteration using the Iterator Trait

• Recall an earlier example:

```
let a = vec![10, 20, 30, 40, 50];
for e in a.iter() {
    println!("the value is: {}", e);
}
```

 The iter() method returns an *iterator*, i.e., a value with the Iterator trait

```
trait Iterator {
  type Item; //this is an associated type
  fn next(&mut self) -> Option<Self::Item>;
  ... //default method impls
}
```

## Unpacking the **for** syntax

- Each call to next advances the iterator
  - So it has to be **mut**

```
let a = vec![10, 20];
let mut iter = a.iter();
assert_eq!(iter.next(), Some(&10));
assert_eq!(iter.next(), Some(&20));
assert_eq!(iter.next(), None);
```

- calls to next produce immutable references to the values in a
  - else may call into\_iter or iter\_mut on a to get different sorts of references

#### **Iterator** Adaptors

- We can make one iterator from another
  - An iterator is consumed as it used; it is *lazy*
- This is a pattern for higher order programming
  - i.map(f) produces an iterator returning f(e) for each of i's elements e
  - i.filter(f) produces iterator for i's elements e such that
    f(e) == true
  - i.collect() converts an iterator into a vector
  - i.fold(a,f) is like OCaml's fold\_right
    - fold\_right f a v where v is the list corresponding to i

- zip, sum, ...

#### Examples

```
let a = vec![10,20];
let i = a.iter();
let j = i.map(|x| x+1).collect(); //[11,21]
let k = a.iter().fold(0,|a,x| x-a); //10
for e in a.iter().filter(|&&x| x == 10) {
    println!("{}",e);
} //prints 10
```

## Quiz 1: Output of the following code

```
fn main() {
    let a = [0, 1, 2, 3, 4, 5];
    let mut iter2 = a.iter().map(|x| 2 * x);
    iter2.next();
    let t2 = iter2.next();
    println!("{:?}", t2)
}
```

- A. Some(0)
- B. Some(1)
- C. Some(2)
- D. Some(4)

## Quiz 1: Output of the following code

```
fn main() {
    let a = [0, 1, 2, 3, 4, 5];
    let mut iter2 = a.iter().map(|x| 2 * x);
    iter2.next();
    let t2 = iter2.next();
    println!("{:?}", t2)
}
```

A. Some(0)B. Some(1)C. Some(2)D. Some(4)

#### **Iterator Notes**

- You can make your own iterators too
  - Implement the Iterator trait
  - Several examples in the Rust Book
- Iterators perform extremely well
  - Better that for loops with explicit indexes!
  - This is because Rust aggressively optimizes the code it generates, e.g., by unrolling the iteration loop
  - So feel free to program using map, fold, zip, etc.

## Back to Closures: Passing as Arguments

- Each closure has a distinct type
  - Even if two closures have the same signature, their types are considered different
    - Such types are called *generative* types
- To specify the type of a closure (for a function parameter, say), use generics with trait bounds
  - Fn t (will describe later)
  - FnMut t
  - FnOnce t
- Functions (defined with fn f...) implement the above trait bounds too

#### Using the Fn Trait Trait bound on **T** to specify type of **f** fn app int $\langle T \rangle (f:T,x:i32) \rightarrow i3$ where $T: Fn(i32) \rightarrow i32$ ł f(x) fn main() { println!("{}",app\_int((|x| x-1),1));

But cannot write

fn app\_int(f:(i32) -> i32,x:i32) -> i32
{ f(x) }

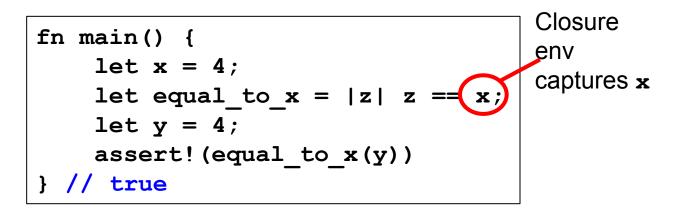
 Can also use function trait bounds in struct, enum, etc. definitions

CMSC 330 - Spring 2021

## Using the Fn Trait Polymorphically

```
fn app<T,U,W>(f:T,x:U) -> W
    where T:Fn(U) -> W
{
    f(x)
}
fn main() {
    println!("{}",app((|x| x-1),1));//i32
    let s = String::from("hi ");
    println!("{}",app(|x| x+"there",s));//String
}
```

# **Capturing Free Variables**



- Note: fails if equal\_to\_x defined as a local function
  - Local functions do not have an environment
- Complication: What if **x** is owned?
  - Capturing it could move it or borrow (mut or immut)
  - Use various **Fn**X traits to specify what to do

# **Distinguishing Fn Trait Bounds**

- FnOnce t (where t is a func type)
  - Consumes the variables it captures from its enclosing scope (i.e., moves or copies them)
  - Thus can only be called once
    - The call consumes ownership
- FnMut t
  - Borrows captured variables mutably
- Fn t
  - Borrows captured variables immutably, or copies
    - equal\_to\_x copied x due to its Copy trait
  - Try this bound first; follow the compiler's advice if it doesn't work

#### Example use of FnOnce

let x = String::from("hi"); let add\_x = |z| x+z; //captures x; is FnOnce println!("x = {}",x); //fails let s = add\_x(" there");//consumes closure let t = add\_x(" joe");//fails, add\_x consumed