CMSC 330 Organization of Programming Languages

OCaml Higher Order Functions

CMSC330 -Spring 2022

Anonymous Functions

Recall code blocks in Ruby

(1..10).each { |x| print x }

- Here, we can think of { |x| print x } as a function
- We can do this (and more) in OCaml

Anonymous Functions

- As with Ruby, passing around functions is common
 - So often we don't want to bother to give them names
- Use fun to make a function with no name



Anonymous Functions

- Syntax
 - fun <u>x1</u> ... <u>xn</u> -> e
- Evaluation
 - An anonymous function is an expression
 - In fact, it is a value no further evaluation is possible
 - As such, it can be passed to other functions, returned from them, stored in a variable, etc.
- Type checking
 - (fun x1 ... xn -> e):(t1 -> ... -> tn -> u)

when e: u under assumptions x1: t1, ..., xn: tn.

 $\square (Same rule as let <math>f x1 \dots xn = e)$

Quiz 1: What does this evaluate to?

let y = (fun x
$$->$$
 x+1) 2 in
(fun z $->$ z-1) y

A. Error
B. 2
C. 1
D. 0

D. 0

Quiz 1: What does this evaluate to?

A. Error
B. 2
C. 1
D. 0

Quiz 2: What is this expression's type ?

$$(fun x y -> x) 2 3$$

- A. Type error
- B. int
- C. int -> int -> int
- D. 'a -> 'b -> 'a

Quiz 2: What is this expression's type ?

$$(fun x y -> x) 2 3$$

A. Type error
B. int
C. int -> int -> int
D. 'a -> 'b -> 'a

Functions and Binding

 Functions are first-class, so you can bind them to other names as you like

let f x = x + 3;;

let g = f;;

- g 5 = 8
- In fact, let for functions is a syntactic shorthand
 let f x = body
 - is semantically equivalent to

let $f = fun x \rightarrow body$

Example Shorthands

- let next x = x + 1
 - Short for let next = fun $x \rightarrow x + 1$

- let plus x y = x + y
 - Short for let plus = fun x y -> x + y

Quiz 3: What does this evaluate to?

- **A**. 0
- B. 1
- C. 2
- D. Error

Quiz 3: What does this evaluate to?

A. 0

- B. 1
- C. 2
- D. Error

Defining Functions Everywhere

```
let move 1 x =
  let left x = x - 1 in (* locally defined fun *)
  let right x = x + 1 in (* locally defined fun *)
  if 1 then left x
 else right x
;;
let move' l x = (* equivalent to the above *)
  if 1 then (fun y \rightarrow y - 1) x
 else (fun y \rightarrow y + 1) x
```

Pattern Matching With Fun

• match can be used within fun

```
(fun 1 \rightarrow match 1 with (h::_) \rightarrow h) [1; 2] = 1
```

But use named functions for complicated matches

May use standard pattern matching abbreviations, too
 (fun (x, y) -> x+y) (1,2) = 3

Passing Functions as Arguments

In OCaml you can pass functions as arguments





The Map Function

OCaml's map is a higher order function; like Ruby's collect

 map f 1 takes a function f and a list 1, applies function f to each element of 1, and returns a list of the results (preserving order)

let add_one x = x + 1
let negate x = -x
map add_one [1; 2; 3] = [2; 3; 4]
map negate [9; -5; 0] = [-9; 5; 0]

How can we implement Map?

```
let rec add1all 1 =
  match 1 with
  [] -> []
  | h::t ->
    (add one h):: add1al1 t
```

```
let rec negall l =
  match l with
   [] -> []
   | h::t ->
    (neg h):: negall t
```

```
let rec map f l =
  match l with
  [] -> []
  | h::t -> (f h)::(map f t)
```

Implementing map

let rec map f l =
 match l with
 [] -> []
 | h::t -> (f h)::(map f t)

• What is the type of map?



Implementing map

let rec map f l =
 match l with
 [] -> []
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• What is the type of map?

Another Example

Apply a list of functions to list of ints

```
let neg x = -x;;
let add_one x = x+1;;
let double x = x + x;;
let fs = [neg; add_one; double];;
let lst = [1;2;3];;
```

```
map (fun f -> map f lst) fs
[[-1; -2; -3]; [2; 3; 4]; [2; 4; 6]]
(neg 1) (neg 2) (neg 3) (add_one 1) ... (double 1) ...
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```

map, as a cartoon



map is included in the standard List module, i.e., as List.map

Quiz 4: What does this evaluate to?

map (fun
$$x \rightarrow x * 4$$
) [1;2;3]

- A. [1.0; 2.0; 3.0]
- B. [4.0; 8.0; 12.0]
- C. Error
- D. [4; 8; 12]

Quiz 4: What does this evaluate to?

map (fun
$$x \rightarrow x * 4$$
) [1;2;3]

- A. [1.0; 2.0; 3.0]
- B. [4.0; 8.0; 12.0]
- C. Error
- D. [4; 8; 12]

map ??? [1; 0; 3] = [true; false; true]

- A. fun $x \rightarrow$ true
- B. fun $x \rightarrow x = 0$
- C. fun $x \rightarrow x != 0$
- D. fun $x \rightarrow x = (x != 0)$

map ??? [1; 0; 3] = [true; false; true]

- A. fun $x \rightarrow$ true
- B. fun $x \rightarrow x = 0$
- C. fun $x \rightarrow x != 0$
- D. fun $x \rightarrow x = (x != 0)$ int bool Note type error!

fold (and foldr)

Two Recursive Functions

Sum a list of ints

Concatenate a list of strings

let rec concat l =
 match l with
 [] -> ""
 | h::t -> h ^ (concat t)

sum [1;2;3;4];;

-: int = 10

concat ["a";"b";"c"];;
- : string = "abc"

Notice Anything Similar?

Sum a list of ints

Concatenate a list of strings

- let rec sum l = let rec concat l =
 match l with
 [] -> 0
 [] -> (+) h (sum t) | h::t -> (^) h (concat t)
- The structure of the two functions is the same!
- Only the parts in red differ
 - What to return for an empty list (0 or "")
 - What function to apply to h and the result of the recursive call (+ or ^)
- foldr abstracts these similarities using higher order functions

CMSC330 Spring 2022 Note: (+) treats + as a prefix function, so 1+2 = (+) 1 2 31

The foldr Function

Sum a list of ints

Concatenate a list of strings:

let rec sum l = let rec concat l =
 match l with
 [] -> 0
 [] -> (+) h (sum t) | h::t -> (^) h (concat t)

```
let rec foldr f a l =
  match l with
  [] -> a
  | h::t -> f h (foldr f a t)
```

let sum l = foldr (+) 0 l let concat l = foldr (^) "" l

So, What is foldr?

- foldr is a function that
 - takes a function of two arguments, a final value, and a list
 - processes the list by applying the function to the head and the recursive application of the function to the rest of the list, returning the final value for the empty list

foldr f v [v1; v2; ...; vn] =
 f v1 (f v2 (...(f vn v)...))

S0 foldr add 0 [1;2;3;4] =
 add 1 (add 2 (add 3 (add 4 0))) = 10

• List.fold_right in the standard library is foldr, but with the order of its last two parameters reversed, i.e.,

fold_right f [v1; v2; ...; vn] v =
 f v1 (f v2 (...(f vn v)...))

Fold (aka fold_left)

- The List module also defines fold_left
 - which we will just call fold



Computes **f** on the *accumulator* **a** and the head **h**, then passes the result as the accumulator to the recursive call

• Similar to foldr, but changes the order of operations

```
let rec foldr f a l =
  match l with
  [] -> a
  | h::t -> f h (foldr f a t)
```

What does **fold** do?

```
let rec fold f a l =
  match l with
  [] -> a
  | h::t -> fold f (f a h) t
```

We just built the sum function!

Fold (aka fold_left)

• What does fold do?

let rec fold f a l =
 match l with
 [] -> a
 | h::t -> fold f (f a h) t

- fold f v [v1; v2; ...; vn]
- = fold f (f v v1) [v2; ...; vn]
- = fold f (f (f v v1) v2) [...; vn]
- = f (f (f (f v v1) v2) ...) vn

• e.g., fold add 0 [1;2;3;4] =
add (add (add (add 0 1) 2) 3) 4 = 10

Another Example

```
let rec fold f a l =
  match l with
  [] -> a
  | h::t -> fold f (f a h) t
```

We just built the length function!

Using Fold to Build Reverse

```
let rec fold f a l =
  match l with
  [] -> a
  | h::t -> fold f (f a h) t
```

Let's build the reverse function with fold!

```
let prepend a x = x::a
fold prepend [] [1; 2; 3; 4] \rightarrow
fold prepend [1] [2; 3; 4] \rightarrow
fold prepend [2; 1] [3; 4] \rightarrow
fold prepend [3; 2; 1] [4] \rightarrow
fold prepend [4; 3; 2; 1] [] \rightarrow
[4; 3; 2; 1]
```

let f x y = if x > y then x else y in fold f 0 [3;4;2]

- **A**. 0
- B. true
- C. 2
- D. 4

let f x y = if x > y then x else y in fold f 0 [3;4;2]

- **A**. 0
- B. true
- C. 2
- D. 4

Quiz 7: What does this evaluate to?

fold (fun a y
$$->$$
 a-y) 0 [3;4;2]

A. -9
B. -1
C. [2;4;3]
D. 9

Quiz 7: What does this evaluate to?

fold (fun a y -> a-y) 0 [3;4;2]

A. -9
B. -1
C. [2;4;3]
D. 9

Type of fold_left, fold_right



Summary: Left-to-right vs. right-to-left

- fold_left f v [v1; v2; ...; vn] =
 f (f (f (f v v1) v2) ...) vn
- fold_right f [v1; v2; ...; vn] v =
 f v1 (f v2(... (f vn v) ...))
- fold_left (fun x y -> x y) 0 [1;2;3] = -6
 since ((0-1)-2)-3) = -6

fold_right [1;2;3] (fun x y -> x - y) 0 = 2
since 1-(2-(3-0)) = 2

When to use one or the other?

- Many problems lend themselves to fold_right
- But it does present a performance disadvantage
 - The recursion builds of a deep stack: One stack frame for each recursive call of fold_right
- An optimization called tail recursion permits optimizing fold_left so that it uses no stack at all
 - We will see how this works in a later lecture!

Combining map and fold

- Idea: map a list to another list, and then fold over it to compute the final result
 - Basis of the famous "map/reduce" framework from Google, since these operations can be parallelized

```
let countone 1 =
   fold (fun a h -> if h=1 then a+1 else a) 0 1
let countones ss =
   let counts = map countone ss in
   fold (fun a c -> a+c) 0 counts
countones [[1;0;1]; [0;0]; [1;1]] = 4
countones [[1;0]; []; [0;0]; [1]] = 2
```

fold & map More examples, practice

Map Example 1: Permute a list

```
let permute lst =
  let rec rm x l = List.filter ((<>) x) l
  and insertToPermute lst x =
    let t = rm x lst in
    List.map ((fun a b \rightarrow a::b) \times)(permuteall t)
  and permuteall lst =
    match 1st with
    |[]->[]
    [[x]->[[x]]
    ->List.flatten(List.map (insertToPermute lst) lst)
  in permuteall lst
;;
 # permute [1;2;3];;
  -: int list list =
  [[1; 2; 3]; [1; 3; 2]; [2; 1; 3]; [2; 3; 1]; [3; 1; 2]; [3; 2; 1]]
```

Map Example 2: Power Set

```
let populate a b =
    if b=[] then [[a]]
    else let t = List.map (fun x->a::x) b in
        [a]::t@b
;;
```

- # populate 1 [[2];[3]];;
- : int list list =
- [[1]; [1; 2]; [1; 3]; [2]; [3]]

```
let powerset lst = List.fold_right populate lst []
;;
```

```
# powerset [1;2;3];;
- : int list list = [[1]; [1; 2]; [1; 2; 3]; [1; 3]; [2]; [2; 3]; [3]]
```

Fold Example 1: Product of an int list

let mul x y = x * y;; let lst = [1; 2; 3; 4; 5];; fold mul 1 lst - : int = 120



Fold Example 2: Count elements of a list satisfying a condition

```
countif (fun x -> x > 0) [30;-1;45;100;0];;
```

```
-: int = 3
```

Fold Example 3: Collect even numbers in the list

let f acc y = if (y mod 2) = 0 then y::acc
 else acc;;

fold f [] [1;2;3;4;5;6];;

- : int list = [6; 4; 2] Reversed

Fold Example 4: Inner Product

first compute list of pair-wise products, then sum up

```
[x1;x2;x3]*[y1;y2;y3] = x1*y1 + x2*y2 + x3*y3
let rec map2 f a b =
   match (a,b) with
   |([1,[1)->([1))|
   (h1::t1,h2::t2)->(f h1 h2):: (map2 f t1 t2)
   ->invalid arg "map2";;
let product v1 v2 =
      fold (+) 0 (map2 ( * ) v1 v2);;
# val product : int list -> int list -> int = <fun>
product [2;4;6] [1;3;5];;
#-: int = 44
```

Fold Example 5: Find the maximum from a list

```
maxList [3;10;5];;
- : int = 10
```

```
(*
maxList [3;10;5]
fold max 3 [10:5]
fold max (max 3 10) [5]
fold max (max 10 5) []
fold max 10 []
10 *)
```

Quiz: Sum of sublists

Given a list of int lists, compute the sum of each int list, and return them as list.

For example:

- sumList [[1;2;3];[4];[5;6;7]]
- : int list = [6; 4; 18]

let sumList = map (fold (+) 0);;

sumList [[1;2;3];[4;5;6];[10]];;

- : int list = [6; 15; 10]

Quiz: Maximum contiguous sublist

Given an int list, find the contiguous sublist, which has the largest sum and return its sum.

Example:

Input: [-2,1,-3,**4,-1,2,1**,-5,4] Output: 6 Explanation: [4,-1,2,1] has the largest sum = 6

Quiz: Maximum contiguous sublist

```
let f (m, acc) h = (
   let m = max m (acc + h) in
   let x = if acc < 0 then 0 else acc in
   (m, x+h)
;;
let submax lst = let (max_so_far, max current) =
      fold f(0,0) lst in
      max_so_far
;;
submax [-2; 1; -3; 4; -1; 2; 1; -5; 4];;
-: int = 6
```

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