

# Problem Set #2

Quantum Error Correction  
Instructor: Daniel Gottesman

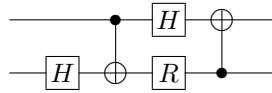
Due Tuesday, Mar. 5, 2024, 5 PM

**Note:** This problem set has problems on 2 pages.

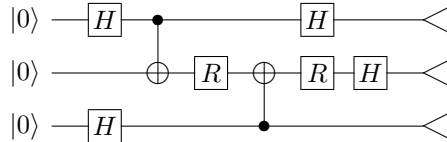
**Problem #1. Analyzing Clifford group circuits (25 points)**

In the following diagrams,  $R = R_{\pi/4}$  is the matrix  $\text{diag}(1, i)$  and  $H$  is the Hadamard transform.

- a) (10 points) For the following Clifford group circuit, compute the overall action on Paulis and use that to write down the  $4 \times 4$  unitary matrix performed by the circuit:



- b) (15 points) For the following Clifford group circuit, use Clifford simulation techniques to compute the full probability distribution of the 8 possible classical outputs after measuring all qubits in the computational basis:



**Problem #2. Twirling (20 points)**

Let  $S(\rho)$  be a quantum operation (a completely positive trace-preserving map) taking  $n$  qubits to  $n$  qubits. **Hint:** (For both parts) Any  $2^n \times 2^n$  matrix can be expanded in the basis of Pauli operators.

- a) (10 points) Consider the following quantum operation: Choose a uniformly random  $P \in \mathcal{P}_n / \{\pm I, \pm iI\}$  (i.e., a Pauli ignoring global phase). Apply  $P^\dagger$ , then  $S$ , then  $P$  (for the same  $P$ ). Show that, averaging over  $P$ , the resulting quantum operation is a Pauli channel. A Pauli channel is any channel  $\mathcal{S}$  such that  $\mathcal{S}(\rho) = (1 - \sum_P p_P) \rho + \sum_P p_P P \rho P^\dagger$ . (Where the sums are over non-trivial Paulis without phases,  $\mathcal{P}_n / \{\pm 1, \pm i\} \setminus \{I\}$ .)
- b) (10 points) Now instead of choosing a random Pauli, choose a random Clifford and do the same thing, i.e., uniformly random  $C \in \mathcal{C}_n / \{e^{i\phi} I\}$ , apply  $C^\dagger$ , then  $S$ , then  $C$ . Show that, averaging over  $C$ , the resulting quantum channel is a depolarizing channel. A depolarizing channel is a channel  $\mathcal{S}$  such that  $\mathcal{S}(\rho) = [1 - (2^n - 1)p] \rho + \sum_P p P \rho P^\dagger$ . (Where the sums are over non-trivial Paulis without phases,  $\mathcal{P}_n / \{\pm 1, \pm i\} \setminus \{I\}$ .) That is, an  $n$ -qubit depolarizing channel is a Pauli channel where all probabilities  $p_P$  are the same.

**Problem #3. Logical operations for qudit code (15 points)**

Consider the following stabilizer code for qudits (qudits with dimension  $p = 3$ ):

$$\begin{array}{cccc} X & X & Z & Z \\ Z & Z & X & X \end{array}$$

- a) (5 points) Show that these generators define a qudit stabilizer code. What are its parameters as a QECC?
- b) (5 points) Find a generating set for the logical Pauli group. (I.e., coset representatives for  $\overline{X}_i$  and  $\overline{Z}_i$ ).
- c) (5 points) For your choice of logical Pauli operators, write down the codeword for logical  $|\overline{0}\rangle$  states (for all logical qubits) expanded in the standard basis for the physical qubits.