

Final Exam

*Open book and notes; In class**Tuesday, Dec. 19th, 1030a*

- ⊕ *Do not forget to write your name on the first page. Initial each subsequent page.*
- ⊕ *Be **neat** and **precise**. I will not grade answers I cannot read.*
- ⊕ *You should draw simple figures if you think it will make your answers clearer.*
- ⊕ *Good luck and remember, brevity is the soul of wit*

- All problems are mandatory
- I cannot stress this point enough: **Be precise**. If you have written something incorrect along with the correct answer, you should **not** expect to get all the points. I will grade based upon what you **wrote**, not what you **meant**.
- Maximum possible points: 50.

Name: _____

Problem	Points
1	
2	
3	
4	
5	
Total	

1. Nomenclature

(a) Describe the following: (2 points each)

- Burst error

- Distribution System (DS)

- DNS Zone

- Go-Back-N

- Exposed Node Problem

2. Network and Transport

(a) List two uses of the AS-path attribute in BGP. (3 points)

(b) Why does TCP require the `TIME-WAIT` state? (3 points)

(c) Suppose you enable the TCP window scale (WSCALE) option (and use maximum scaling), but do *not* enable the Timestamps (TS). What is the maximum throughput you can achieve on a 40Gbps, 100ms RTT end-to-end link? Explain why. (4 points)

3. MAC Protocols

- [illegible]

4. DNS/Application Layer

- (a) Why do typographic errors form a major fraction of queries received at DNS root servers? (2 points)

- (b) What is the per-peer state kept by BitTorrent peers? (Describe each state, as opposed to just naming them.) (2 points)

- (c) Suppose *all* the peers in a BitTorrent swarm are behind NATs. (How) Would the protocol function? (3 points)

- (d) Consider a Chord alternate that replicates items r times by “naming” an item using r different hash functions, and inserting each item at these r names. Compare this design to Chord’s original technique of replicating at $r - 1$ successors. (3 points)

5. General

- (a) A CRC with divisor degree n can detect all n bit burst errors. Can a n -bit CRC detect all errors in a message with n bits of data? Why/why not? (2 points)
- (b) The Ethernet CRC polynomial is $x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11} + x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x^1 + 1$. Can it detect an odd number of errors? If yes, explain why. If not, give an example of a transmitted message and received where an odd number of errors is not detected. (4 +1 bonus points)
- (c) Suppose you administer the domain `ispyon.everyone` and have delegated a subdomain `secret.ispyon.everyone`. (How) can you ensure that you can log every IP address that looks up a name in the delegated domain? (4 +1 bonus points)