

CMSC 216
Introduction to Computer Systems
Lecture 12

Administrivia

- Read Reek, Chapter 12
- Start reading Bryant and O'Hallaron Chapter 3 (IA-32 instruction set architecture) and 4.1 (Y86 subset)

Project 3b

- Let's talk about the project

Chapters 3 and 4.1, Bryant and O'Hallaron

ASSEMBLY LANGUAGE (CONT.)

Y86 integer instructions

Instruction	Effect	Description
addl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] + \text{Reg}[S]$	Addition
subl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] - \text{Reg}[S]$	Subtract
andl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] \& \text{Reg}[S]$	Bitwise AND
xorl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] \wedge \text{Reg}[S]$	Bitwise XOR
multl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] * \text{Reg}[S]$	Multiplication*
divl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] / \text{Reg}[S]$	Integer division*
modl S,D	$\text{Reg}[D] \leftarrow \text{Reg}[D] \% \text{Reg}[S]$	Remainder*

- All these instructions operate on two integers, and set the condition code flags appropriately
- Notice that only registers are used
- Instructions marked with an asterisk (*) are extensions to Y86 we've added to the ones in the book

Integer instruction example

- Assembler output:

```
0x000: 308003000000 | irmovl $3,%eax    # a = 3
0x006: 308305000000 | irmovl $5,%ebx    # b = 5
0x00c: 6003          | addl    %eax,%ebx  # b = a + b
0x00e: f308          | wrint   %eax
0x010: 308620000000 | irmovl $32,%esi   # 32 == ' '
0x016: f168          | wrch    %esi
0x018: f338          | wrint   %ebx
0x01a: 30860a000000 | irmovl $10,%esi   # 10 == '\n'
0x020: f168          | wrch    %esi
0x022: 10           | halt
```

- Simulator run:

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...

- Notice these instructions are destructive; they overwrite the second operand
 - Need to make copies if you need old values

Condition codes

- Performing integer operations causes various flags to be set, describing the attributes of the result of the operation
- These are used by other, subsequent instructions to perform conditional branching
- The three we are concerned with are:
 - OF: overflow flag; did the operation overflow?
 - SF: sign flag; is the result negative?
 - ZF: zero flag; is the result zero?

Branch instructions

- These are used to perform the effect of if statements, loops, and switches
- When encountered, if a certain condition is true, control flow will then go to the address specified, rather than advancing to the next instruction
 - The address of the next instruction to be executed is held in the program counter; in many architectures, this is held in an accessible register (not so with Y86).
- Labels
 - We use labels (name followed by colon) to represent target addresses
 - When the assembler encounters a label (e.g., **Loop:**), the address of the labeled instruction is used by the assembler to replace all references to that label in the program
 - Labels do not need to be declared before use

Y86 branch instructions

Instruction	Branch if...	Description
<code>jmp Label</code>	1	Unconditional jump
<code>jle Label</code>	$(SF \wedge OF) \vee ZF$	Jump if less than or equal to zero
<code>jlt Label</code>	$SF \wedge OF$	Jump if less than zero
<code>je Label</code>	ZF	Jump if equal to zero
<code>jne Label</code>	$\sim ZF$	Jump if not equal to zero
<code>jge Label</code>	$\sim (SF \wedge OF)$	Jump if greater than or equal to zero
<code>jgt Label</code>	$\sim (SF \wedge OF) \ \& \ \sim ZF$	Jump if greater than zero

- Each instruction relies on the condition codes set by the most recent integer instruction

Translating branches

- Since assembly has no else statements, we need to use branching + code reorganization to get if/else
- Consider the following C code:

```
if (i == j)
    printf("=") ;
else
    printf("X") ;
printf("\n") ;
```

```
if (i == j)
    goto Equal ;
    printf("X") ;
    goto EndIf ;
Equal :
    printf("=") ;
EndIf :
    printf("\n") ;
```

Translating branches, cont.

- We can then take the labeled C code and translate it in a fairly straightforward fashion:

```
    subl %eax,%ebx    # i:%eax,j:%ebx
```

```
    je Equal
```

```
    irmovl $88,%ecx  # else block
```

```
    wrch %ecx
```

```
    jmp EndIf
```

```
Equal: irmovl $61,%ecx # true block
```

```
    wrch %ecx
```

```
EndIf: irmovl $10,%ecx # after if/else
```

```
    wrch %ecx
```

Branch Example

- Assembler output:

```
0x000: 308003000000 |      irmovl $3,%eax      # a = 3
0x006: 308304000000 |      irmovl $4,%ebx      # b = 4
0x00c: 2006          |      rrmovl %eax,%esi     # s = a
0x00e: 6136          |      subl   %ebx,%esi     # s = s - b
0x010: 751c000000   |      jge    Else         # if s >= 0 jump
                                |
0x015: f308          |      wrint  %eax          # printf("%d", a)
0x017: 701e000000   |      jmp    Endif        # jump
                                |
0x01c: f338          | Else: wrint  %ebx          # printf("%d", b)
                                |
0x01e: 30860a000000 | Endif: irmovl $10,%esi    #
0x024: f168          |      wrch   %esi          # printf("\n")
0x026: 10           |      halt
```

- Simulator output:

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Stopped in 10 steps at PC = 0x27. Exception 'HLT', CC Z=0 S=1 O=0

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