

CMSC 351: Practice Questions for Final Exam

These are practice problems for the upcoming final exam. You will be given a sheet of notes for the exam. Also, go over your homework assignments. **Warning:** This does not necessarily reflect the length, difficulty, or coverage of the actual exam.

Problem 1. Let $A[1, \dots, n]$ be an array of n numbers (some positive and some negative).

- (a) Give an algorithm to find which *three* numbers have sum closest to zero. Make your algorithm as efficient as possible. Write it in pseudo-code.
- (b) Analyze its running time.

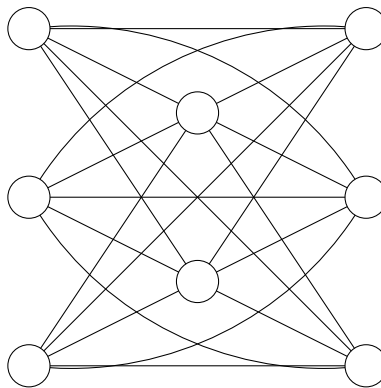
Problem 2. Assume that you developed an algorithm to find the (index of the) $n/3$ smallest element of a list of n elements in $2n$ comparisons.

- (a) Using the algorithm (as a black box), give an algorithm, efficient in the worst case, to find the k th smallest element of a list.
- (b) Write down a recurrence for (a bound on) the number of comparisons it executes in the worst case.
- (c) Solve the recurrence (using constructive induction). Find the high order term exactly (but you do not need any low order terms).
- (d) Using the (black box) algorithm for finding the $n/3$ smallest element and using the ideas and results of Parts (a), (b), and (c), give an efficient algorithm to find (the index of) two elements, the k_1 th smallest and the k_2 smallest (for inputs k_1 and k_2). The algorithm description can be very high level and brief.
- (e) How many comparisons does it use? Find the high order term exactly (but you do not need any low order terms). Give a brief justification.

Problem 3. Assume you want to sort a list of n numbers that are in k groups, so that the smallest n/k are first, then the next smallest n/k are second, etc.

- (a) Give a decision tree based lower bound on the time to produce a single sorted list. (You could base a lower bound on the fact that there really are k independent sorting problems. Do NOT use this lower bound argument.)
- (b) Give an algorithm for sorting this list. How fast is the algorithm? (This is an upper bound.)

Problem 4. A graph is tripartite if the vertices can be partitioned into three sets so that there are no edges internal to any set. The *complete* tripartite graph, $K(a, b, c)$, has three sets of vertices with sizes a , b , and c and all possible edges between each pair of sets of vertices. $K(3, 2, 3)$ is pictured below. A *Hamiltonian* cycle in a graph is a cycle that traverses every vertex exactly once.



- (a) For which values of n does $K(1, 1, n)$ have a Hamiltonian cycle. Justify your answer.
- (b) For which values of n does $K(1, n, n)$ have a Hamiltonian cycle. Justify your answer.
- (c) For which values of n does $K(n, n, n)$ have a Hamiltonian cycle. Justify your answer.

Problem 5. Let $G = (V, E)$ be an undirected graph. A *triangle* is a set of three vertices such that each pair has an edge.

- (a) Give an efficient algorithm to find all of the triangles in a graph.
- (b) How fast is your algorithm?

Problem 6. This problem is more open-ended than you would see on an exam: If you do not know how to play Sudoku, look it up. Normally, Sudoku is played on a 9×9 grid.

- (a) Generalize Sudoku to larger grids.
- (b) State the (generalized) Sudoku game as a decision problem.
- (c) Show that the decision version of (generalized) Sudoku is in NP.
- (d) Show that if you can solve the decision version of (generalized) Sudoku in polynomial time, you can solve a (generalized) Sudoku puzzle in polynomial time.

Problem 7. In a graph $G = (V, E)$ edge (x, y) *touches* vertices x and y .

A *newtonian cluster* in a graph $G = (V, E)$ is a subset of the edges such that every vertex is touched by at least one edge. The *size of a newtonian cluster* is the number of edges in the subset.

- (a) Give an example of a graph that has a newtonian cluster of size four but not of size three.
- (b) Let C be a newtonian cluster of G . What can you say about the minimum size of C as a function of the number of vertices n ? Justify.
- (c) A *minimal newtonian cluster* is a newtonian cluster such that if any edge is removed it is no longer a newtonian cluster. Let C be a minimal newtonian cluster of G . What can you say about the maximum size of C as a function of the number of vertices n ? Justify.
- (d) The *newtonian cluster problem* is given a graph $G = (V, E)$ and an integer k does G have a newtonian cluster of size (at most) k . Show that the newtonian cluster problem is in **NP**. What is the certificate?

Problem 8. The *Weighted Independent Set Problem (WISP)* is, given a graph $G = (V, E)$ with integer weights on the nodes, find a subset of the nodes whose sum of weights is as large as possible such that there is no edge between any pair of nodes in the subset.

- (a) WISP is an optimization problem. Define a decision version of WISP.
- (b) Show that the decision version is in **NP**. What is the certificate?
- (c) Show that if you could solve the optimization version in polynomial time that you could also solve the decision version in polynomial time.
- (d) Show that if you could solve the decision version in polynomial time that you could also solve the optimization version in polynomial time. HINT: First find the weight of an optimal independent set.