

CMSC 330: Organization of Programming Languages

Lets, Tuples, Records

Let Expressions

- Enable binding variables in other expressions
 - These are different from the `let` *definitions* we've been using at the top-level
- They are expressions, so they have a value
- **Syntax**
 - `let x = e1 in e2`
 - **x** is a *bound variable*
 - **e1** is the *binding expression*
 - **e2** is the *body expression*

Let Expressions

- Syntax

- `let x = $e1$ in $e2$`

- Evaluation

- Evaluate $e1$ to $v1$

- Substitute $v1$ for x in $e2$ yielding new expression $e2'$

- Evaluate $e2'$ to $v2$

- Result of evaluation is $v2$

Example

```
let x = 3+4 in 3*x
```

```
➤ let x = 7 in 3*x
```

```
➤ 3*7
```

```
➤ 21
```

Let Expressions

- Syntax
 - `let x = $e1$ in $e2$`
- Type checking
 - If `$e1 : t1$` and `$e2 : t$` (assuming `$x : t1$`)
 - Then `let x = $e1$ in $e2 : t$`
- Example: `let x = $3+27$ in $x*3$`
 - `$3+27 : int$`
 - `$x*3 : int$` (assuming `$x:int$`)
 - **so** `let x = $3+27$ in $x*3 : int$`

Let Definitions vs. Let Expressions

- At the top-level, we write
 - `let x = e;;` (* no `in e2` part *)
 - This is called a let *definition*, not a let *expression*
 - Because it doesn't, itself, evaluate to anything
- Omitting `in` means “from now on”:
 - # `let pi = 3.14;;`
 - (* `pi` is now *bound* in the rest of the top-level scope *)

Top-level expressions

- We can write any expression at top-level, too
 - `e;;`
 - This says to evaluate `e` and then ignore the result
 - Equivalent to `let _ = e;;`
 - Useful when `e` has an effect, such as reading/writing a file, printing to the screen, etc.

```
let x = 37;;  
let y = x + 5;;  
print_int y;;  
print_string "\n";;
```

- When run, outputs 42 to the screen

Let Expressions: Scope

- In `let x = e1 in e2`, variable `x` is *not* visible outside of `e2`

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;  
print_float pi;;
```

error: `pi` not bound

bind `pi` (only) in body of `let`
(which is `pi *. 3.0 *. 3.0`)

Binding in other languages




- Compare to similar usage in Java/C

```
let pi = 3.14 in
  pi *. 3.0 *. 3.0;;
pi;; (* pi unbound! *)
```




```
{
  float pi = 3.14;

  pi * 3.0 * 3.0;
}
pi; /* pi unbound! */
```


Examples – Let

- $x;;$
– (* Unbound value x *)

- $\text{let } x = 1 \text{ in } x + 1;;$
– (* 2 *)

- $\text{let } x = x \text{ in } x + 1;;$
– (* Unbound value x *)


Examples – Let

- $\text{let } x = 1 \text{ in } (x + 1 + x) \ ;;$
– (* 3 *)

- $(\text{let } x = 1 \text{ in } x + 1) \ ;;\ x \ ;;$
– (* Unbound value x *)

- $\text{let } x = 4 \text{ in } (\text{let } x = x + 1 \text{ in } x) \ ;;$
– (* 5 *)


Shadowing Names

- **Shadowing** is rebinding a name in an inner scope to have a different meaning
 - May or may not be allowed by the language

C

```
int i;

void f(float i) {
    {
        char *i = NULL;
        ...
    }
}
```

OCaml

```
let g = 3;;
let g x = x + 3;;
```

Java

```
void h(int i) {
    {
        float i; // not allowed
        ...
    }
}
```

Let Expressions in Functions

- You can use **let** inside of functions for local vars

```
let area r =  
  let pi = 3.14 in  
  pi *. r *. r
```

- And you can use many **lets** in sequence

```
let area d =  
  let pi = 3.14 in  
  let r = d /. 2.0 in  
  pi *. r *. r
```

Nested Let Expressions

- Uses of `let` can be nested in OCaml
 - Nested bound variables (`pi` and `r`) invisible outside
- Similar scoping possibilities C and Java

```
let res =  
  (let area =  
    (let pi = 3.14 in  
     let r = 3.0 in  
      pi *. r *. R) in  
   area /. 2.0) ;;
```

```
float res;  
{ float area;  
  { float pi = 3.14  
    float r = 3.0;  
    area = pi * r * r;  
  }  
  res = area / 2.0;  
}
```

Quiz 1

Which of these expressions does **not** evaluate to 3?

A. `let x=3`

B. `let x=2 in x+1`

C. `let x=3 in x`

D. `3`

E. `let f x = x+1 in f 2`

Quiz 1

Which of these expressions does **not** evaluate to 3?

A. let x=3 ---> not an expression

B. let x=2 in x+1

C. let x=3 in x

D. 3

E. let f x = x+1 in f 2

Quiz 2: What does this evaluate to?

```
let x = 2 in  
let y = 2 in  
x + y
```

- A. 2
- B. 3
- C. 4
- D. 5

Quiz 2: What does this evaluate to?

```
let x = 2 in  
let y = 2 in  
x + y
```

- A. 2
- B. 3
- C. 4
- D. 5

Quiz 3: What does this evaluate to?

```
let x = 3 in  
let y = 4 in  
let x = 8 in  
x = 10-y
```

- A. 6
- B. true
- C. 12
- D. false

Quiz 3: What does this evaluate to?

```
let x = 3 in  
let y = 4 in  
let x = 8 in  
x = 10-y
```

- A. 6
- B. true
- C. 12
- D. false

Quiz 4: What does this evaluate to?

```
let x = 10 in  
let y = x+2 in  
let x = 8 in  
y
```

- A. 5
- B. 12
- C. 10
- D. false

Quiz 4: What does this evaluate to?

```
let x = 10 in  
let y = x+2 in  
let x = 8 in  
y
```

- A. 5
- B. 12
- C. 10
- D. false

Tuples

- **Constructed** using (e_1, \dots, e_n)
- **Deconstructed** using pattern matching
 - Patterns involve parens and commas, e.g., (p_1, p_2, \dots)
- Tuples are similar to C structs
 - But without field labels
 - Allocated on the heap
- Tuples can be heterogenous
 - Unlike lists, which must be homogenous
 - $(1, ["string1"; "string2"])$ is a valid tuple

Tuple Types

- Tuple types use `*` to separate components
 - Type joins types of its components
- Examples
 - `(1, 2)` :
 - `(1, "string", 3.5)` :
 - `(1, ["a"; "b"], 'c')` :
 - `[(1,2)]` :
 - `[(1, 2); (3, 4)]` :
 - `[(1,2); (1,2,3)]` :

Tuple Types

- Tuple types use `*` to separate components
 - Type joins types of its components
- Examples
 - `(1, 2) : int * int`
 - `(1, "string", 3.5) : int * string * float`
 - `(1, ["a"; "b"], 'c') : int * string list * char`
 - `[(1,2)] : (int * int) list`
 - `[(1, 2); (3, 4)] : (int * int) list`
 - `[(1,2); (1,2,3)] : error`
 - Because the first list element has type `int * int`, but the second has type `int * int * int` – list elements must all be of the same type

Pattern Matching Tuples

```
# let plusThree t =  
  match t with  
    (x, y, z) -> x + y + z;;  
plusThree : int*int*int -> int = <fun>  
  
# let plusThree' (x, y, z) = x + y + z;;  
plusThree' : int*int*int -> int = <fun>  
  
# let addOne (x, y, z) = (x+1, y+1, z+1);;  
addOne : int*int*int -> int*int*int = <fun>  
  
# plusThree (addOne (3, 4, 5));;  
- : int = 15
```

Remember, **semicolon** for lists, **comma** for tuples

- `[1, 2]` = `[(1, 2)]` *which is a list of size one*
- `(1; 2)` *Warning: This expression should have type unit*

More Examples With Tuples

- `let sum ((a, b), c) = (a+c, b+c)`
 - `sum ((1, 2), 3) = (4, 5)`
- `let plusFirstTwo (x::y::_ , a) = (x + a, y + a)`
 - `plusFirstTwo ([1; 2; 3], 4) = (5, 6)`
- `let tls (_::xs, _::ys) = (xs, ys)`
 - `tls ([1; 2; 3], [4; 5; 6; 7]) = ([2; 3], [5; 6; 7])`

Tuples Are A Fixed Size

- This OCaml definition

- `# let foo x = match x with`

- `(a, b) -> a + b`

- `| (a, b, c) -> a + b + c;;`

- Would yield this error message

- This pattern matches values of type `'a * 'b * 'c`

- but is here used to match values of type `'d * 'e`

- Tuples of different size have different types
 - Thus never more than one match case with tuples

Records

- Records: identify elements by **name**
 - Elements of a tuple are identified by **position**
- Define a **record type** before defining record values

```
type date = { month: string; day: int; year: int }
```

- **Construct** a record
 - { **f1=e1**; ...; **fn=en** } : evaluates **e1** to **en**, assigns results to the given fields
 - Fields do not have to be written in order

```
# let today = { day=16; year=2017; month="f"^"eb" };;  
today : date = { day=16; year=2017; month="feb" };;
```

Destructing Records

```
type date = { month: string; day: int; year: int }  
let today = { day=16; year=2017; month="feb" };;
```

- **Access** by **field name** or **pattern matching**

```
print_string today.month;; (* prints feb *)  
(* patterns *)  
let { month=_; day=d } = today in  
let { year } = today in  
let _ = print_int d in      (* prints 16 *)  
print_int year;;          (* prints 2017 *)
```

- **Notes:**
 - In patterns, you can skip or reorder fields
 - You can use the field name as the bound variable

Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in  
get 1 2
```

- A. 3
- B. type error
- C. 2
- D. 1

Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in  
get 1 2
```

A. 3

B. type error – get's first argument must be a pair

C. 2

D. 1

Quiz 6: What does this evaluate to?

```
let get (x,y) =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

Quiz 6: What does this evaluate to?

```
let get (x,y) =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

A. 3

B. type error – get takes only one argument

C. 2

D. 1

Quiz 7: What is the type of `shift`?

```
type point = {x:int; y:int}
let shift { x=px; y=py } =
  {x=px+1; y=py+1};;
```

- A. `point -> bool list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> int list`

Quiz 7: What is the type of `shift`?

```
type point = {x:int; y:int}
let shift { x=px; y=py } =
  {x=px+1; y=py+1};;
```

- A. `point -> bool list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> int list`