

# CMSC 330: Organization of Programming Languages

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Lets, Tuples, Records

# Let Expressions

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- Enable binding variables in other expressions
  - These are different from the `let` *definitions* we've been using at the top-level
- They are expressions, so they have a value
- Syntax
  - `let x = e1 in e2`
  - **x** is a *bound variable*
  - **e1** is the *binding expression*
  - **e2** is the *body expression*

# Let Expressions

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- Syntax

- `let  $x$  =  $e1$  in  $e2$`

- Evaluation

- Evaluate  $e1$  to  $v1$

- Substitute  $v1$  for  $x$  in  $e2$  yielding new expression  $e2'$

- Evaluate  $e2'$  to  $v2$

- Result of evaluation is  $v2$

## Example

```
let x = 3+4 in 3*x
```

```
➤ let x = 7 in 3*x
```

```
➤ 3*7
```

```
➤ 21
```

# Let Expressions

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- Syntax
  - `let  $x = e1$  in  $e2$`
- Type checking
  - If  `$e1 : t1$`  and  `$e2 : t$`  (assuming  `$x : t1$` )
  - Then `let  $x = e1$  in  $e2 : t$`
- Example: `let  $x = 3+27$  in  $x*3$` 
  - `$3+27 : int$`
  - `$x*3 : int$`  (assuming  `$x:int$` )
  - so `let  $x = 3+27$  in  $x*3 : int$`

# Let Definitions vs. Let Expressions

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- At the top-level, we write
  - `let x = e;;` (\* no in e2 part \*)
  - This is called a let *definition*, not a let *expression*
    - Because it doesn't, itself, evaluate to anything
- Omitting `in` means “from now on”:
  - # `let pi = 3.14;;`
  - (\* pi is now *bound* in the rest of the top-level scope \*)

# Top-level expressions

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- We can write any expression at top-level, too
  - `e;;`
  - This says to evaluate `e` and then ignore the result
    - Equivalent to `let _ = e;;`
    - Useful when `e` has an effect, such as reading/writing a file, printing to the screen, etc.

```
let x = 37;;  
let y = x + 5;;  
print_int y;;  
print_string "\n";;
```

- When run, outputs 42 to the screen

# Let Expressions: Scope

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- In `let x = e1 in e2`, variable `x` is *not* visible outside of `e2`

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;  
print_float pi;;
```

error: `pi` not bound

bind `pi` (only) in body of `let`  
(which is `pi *. 3.0 *. 3.0`)

# Binding in other languages

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- Compare to similar usage in Java/C



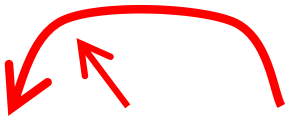
```
let pi = 3.14 in  
  pi *. 3.0 *. 3.0;;  
pi;; (* pi unbound! *)
```

```
{  
  float pi = 3.14;  
  
  pi * 3.0 * 3.0;  
}  
pi; /* pi unbound! */
```





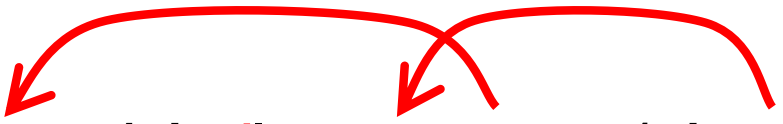
# Examples – Scope of Let bindings

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- $x;;$   
– (\* Unbound value x \*)  

- $\text{let } x = 1 \text{ in } x + 1;;$   
– (\* 2 \*)  

- $\text{let } x = x \text{ in } x + 1;;$   
– (\* Unbound value x \*)  


# Examples – Scope of Let bindings

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- $\text{let } x = 1 \text{ in } (x + 1 + x) \ ;;$   
– (\* 3 \*)  

- $(\text{let } x = 1 \text{ in } x + 1) \ ;;\ x \ ;;$   
– (\* Unbound value x \*)  

- $\text{let } x = 4 \text{ in } (\text{let } x = x + 1 \text{ in } x) \ ;;$   
– (\* 5 \*)  


# Shadowing Names

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- **Shadowing** is rebinding a name in an inner scope to have a different meaning
  - May or may not be allowed by the language

## C

```
int i;  
  
void f(float i) {  
    {  
        char *i = NULL;  
        ...  
    }  
}
```

## OCaml

```
let g = 3;;  
let g x = x + 3;;
```

## Java

```
void h(int i) {  
    {  
        float i; // not allowed  
        ...  
    }  
}
```

# Let Expressions in Functions

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- You can use **let** inside of functions for local vars

```
let area r =  
  let pi = 3.14 in  
  pi *. r *. r
```

- And you can use many **lets** in sequence

```
let area d =  
  let pi = 3.14 in  
  let r = d /. 2.0 in  
  pi *. r *. r
```

# Nested Let Expressions

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- Uses of `let` can be nested in OCaml
  - Nested bound variables (`pi` and `r`) invisible outside
- Similar scoping possibilities C and Java

```
let res =  
  (let area =  
    (let pi = 3.14 in  
     let r = 3.0 in  
      pi *. r *. R) in  
   area /. 2.0) ;;
```

```
float res;  
{ float area;  
  { float pi = 3.14  
    float r = 3.0;  
    area = pi * r * r;  
  }  
  res = area / 2.0;  
}
```

# Quiz 1

---

Which of these is **not** an expression that evaluates to 3?

A. `let x=3`

B. `let x=2 in x+1`

C. `let x=3 in x`

D. `3`

# Quiz 1

---

Which of these is **not** an expression that evaluates to 3?

A. **let x=3 ---> not an expression**

B. **let x=2 in x+1**

C. **let x=3 in x**

D. **3**

## Quiz 2: What does this evaluate to?

---

```
let x = 2 in  
let y = 3 in  
x + y
```

- A. 2
- B. 3
- C. 4
- D. 5



## Quiz 2: What does this evaluate to?

---

```
let x = 2 in  
let y = 3 in  
x + y
```

- A. 2
- B. 3
- C. 4
- D. 5

## Quiz 3: What does this evaluate to?

---

```
let x = 6 in  
let y = 4 in  
let x = 8 in  
x = 10-y
```

- A. 6
- B. true
- C. 12
- D. false

## Quiz 3: What does this evaluate to?

---

```
let x = 6 in  
let y = 4 in  
let x = 8 in  
x = 10-y
```

- A. 6
- B. true
- C. 12
- D. false

## Quiz 4: What does this evaluate to?

---

```
let x = 3 in  
let y = x+2 in  
let x = 8 in  
y
```

- A. 5
- B. 12
- C. 10
- D. false

## Quiz 4: What does this evaluate to?

---

```
let x = 3 in
let y = x+2 in
let x = 8 in
y
```

- A. 5
- B. 12
- C. 10
- D. false

# Tuples

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- **Constructed** using `(e1, ..., en)`
- **Deconstructed** using pattern matching
  - Patterns involve parens and commas, e.g., `(p1,p2, ...)`
- Tuples are similar to C structs
  - But without field labels
  - Allocated on the heap
- Tuples can be heterogenous
  - Unlike lists, which must be homogenous
  - `(1, ["string1";"string2"])` is a valid tuple

# Tuple Types

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- Tuple types use `*` to separate components
  - Type joins types of its components
- Examples
  - `(1, 2)` :
  - `(1, "string", 3.5)` :
  - `(1, ["a"; "b"], 'c')` :
  - `[(1,2)]` :
  - `[(1, 2); (3, 4)]` :
  - `[(1,2); (1,2,3)]` :

# Tuple Types

---

- Tuple types use **\*** to separate components
  - Type joins types of its components
- Examples
  - `(1, 2) : int * int`
  - `(1, "string", 3.5) : int * string * float`
  - `(1, ["a"; "b"], 'c') : int * string list * char`
  - `[(1,2)] : (int * int) list`
  - `[(1, 2); (3, 4)] : (int * int) list`
  - `[(1,2); (1,2,3)] : error`

Because the first list element has type `int * int`, but the second has type `int * int * int` – list elements must all be of the same type



# Pattern Matching Tuples

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```
# let plusThree t =  
  match t with  
    (x, y, z) -> x + y + z;;  
plusThree : int*int*int -> int = <fun>  
  
# let plusThree' (x, y, z) = x + y + z;;  
plusThree' : int*int*int -> int = <fun>  
  
# let addOne (x, y, z) = (x+1, y+1, z+1);;  
addOne : int*int*int -> int*int*int = <fun>  
  
# plusThree (addOne (3, 4, 5));;  
- : int = 15
```

Remember, **semicolon** for lists, **comma** for tuples

- `[1, 2] = [(1, 2)]` which is a list of size one
- `(1; 2)` Warning: This expression should have type unit

# Tuples Are A Fixed Size

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- This OCaml definition

```
- # let foo x = match x with  
  (a, b) -> a + b  
  | (a, b, c) -> a + b + c;;
```

- Would yield this error message

– This pattern matches values of type 'a \* 'b \* 'c  
but is here used to match values of type 'd \* 'e

- Tuples of different size have different types  
– Thus never more than one match case with tuples

# Records

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- Records: identify elements by **name**
  - Elements of a tuple are identified by **position**
- Define a **record type** before defining record values

```
type date = { month: string; day: int; year: int }
```

- **Construct** a record
  - { ***f1=e1***; ...; ***fn=en*** } : evaluates ***e1*** to ***en***, assigns results to the given fields
    - Fields do not have to be written in order

```
# let today = { day=16; year=2017; month="f"^"eb" };;  
today : date = { day=16; year=2017; month="feb" };;
```

# Destructing Records

---

```
type date = { month: string; day: int; year: int }  
let today = { day=16; year=2017; month="feb" };;
```

- **Access** by **field name** or **pattern matching**

```
print_string today.month;; (* prints feb *)  
(* patterns *)  
let { month=_; day=d } = today in  
let { year } = today in  
let _ = print_int d in      (* prints 16 *)  
print_int year;;           (* prints 2017 *)
```

- Notes:
  - In record patterns, you can skip or reorder fields
  - You can use the field name as the bound variable

## Quiz 5: What does this evaluate to?

---

```
let get (a,b) y = a+y in  
get 1 2
```

- A. 3
- B. type error
- C. 2
- D. 1

## Quiz 5: What does this evaluate to?

---

```
let get (a,b) y = a+y in  
get 1 2
```

- A. 3
- B. type error – get's first argument must be a pair
- C. 2
- D. 1

## Quiz 6: What does this evaluate to?

---

```
let get (x,y) =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

## Quiz 6: What does this evaluate to?

---

```
let get (x,y) =  
  match x with  
    (a,b) -> a+y  
in  
get (1,2) 1
```

- A. 3
- B. type error – get takes only one argument
- C. 2
- D. 1



## Quiz 7: What is the type of `shift`?

---

```
type point = {x:int; y:int}
let shift { x=px; y=py } =
  {x=px+1; y=py+1};;
```

- A. `point -> bool list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> int list`

## Quiz 7: What is the type of `shift`?

---

```
type point = {x:int; y:int}
let shift { x=px; y=py } =
  {x=px+1; y=py+1};;
```

- A. `point -> bool list`
- B. `int list -> int list`
- C. `point -> point`
- D. `point -> int list`