

# CMSC 132: Object-Oriented Programming II

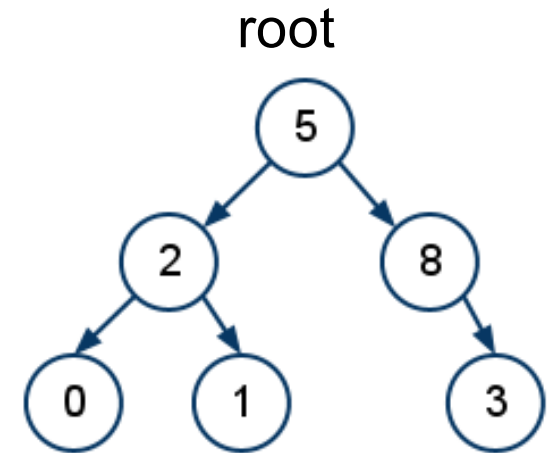
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## Binary Search Trees

# Quiz 1: What is the output?

---

```
void m(Node r) {  
    if(r==null) return;  
    print(r.key+",");  
    m(r.left);  
    m(r.right);  
}  
m(root);
```



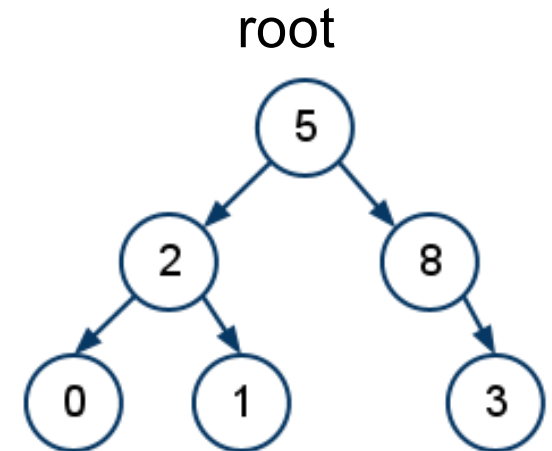
- A. 5,2,0,1,3,8
- B. 5,2,0,1,8,3
- C. 5,2,8
- D. 5,2,0,1

# Quiz 1: What is the output?

---

```
void m(Node r) {  
    if(r==null) return;  
    print(r.key+",");  
    m(r.left);  
    m(r.right);  
}  
m(root);
```

preOrder traversal

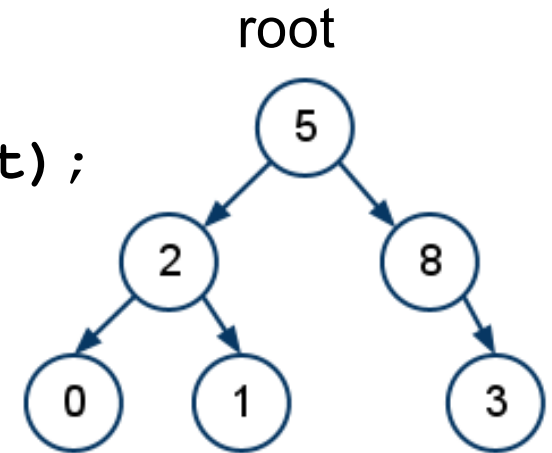


- A. 5,2,0,1,3,8
- B. 5,2,0,1,8,3
- C. 5,2,8
- D. 5,2,0,1

## Quiz 2: What is the output?

---

```
int m(Node r) {  
    if(r==null) return 0;  
    return r.key + m(r.left) + m(r.right);  
}  
m(root);
```

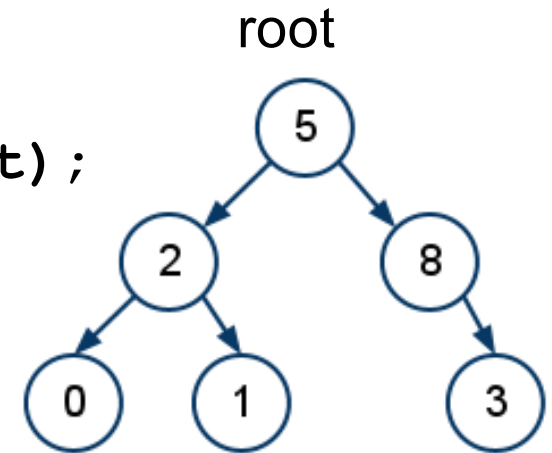


- A. 6
- B. 19
- C. 15
- D. 5

## Quiz 2: What is the output?

---

```
int m(Node r) {  
    if(r==null) return 0;  
    return r.key + m(r.left) + m(r.right);  
}  
m(root);
```



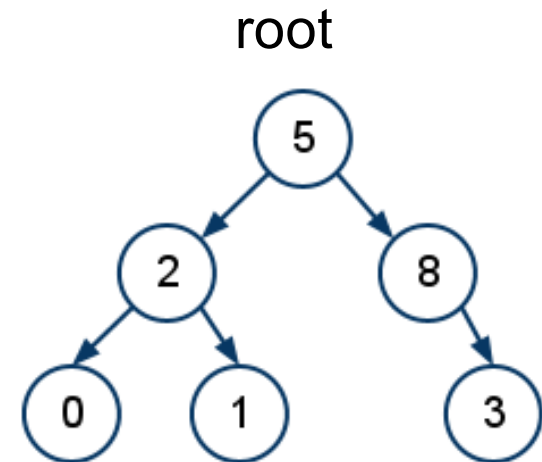
- A. 6
- B. 19
- C. 15
- D. 5

## Quiz 3: What is the output?

---

```
int m(Node r) {  
    if(r==null) return 0;  
    if(r.left==null && r.right==null) return r.key;  
    return m(r.left) + m(r.right);  
}  
m(root);
```

- A. 6
- B. 18
- C. 4
- D. 3

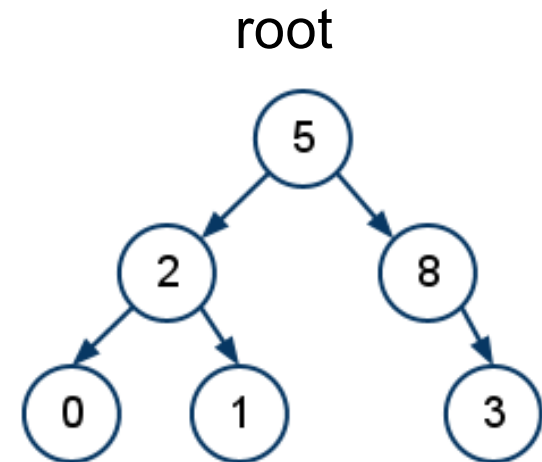


## Quiz 3: What is the output?

---

```
int m(Node r) {  
    if(r==null) return 0;  
    if(r.left==null && r.right==null) return r.key;  
    return m(r.left) + m(r.right);  
}  
m(root);
```

- A. 6
- B. 18
- C. 4
- D. 3



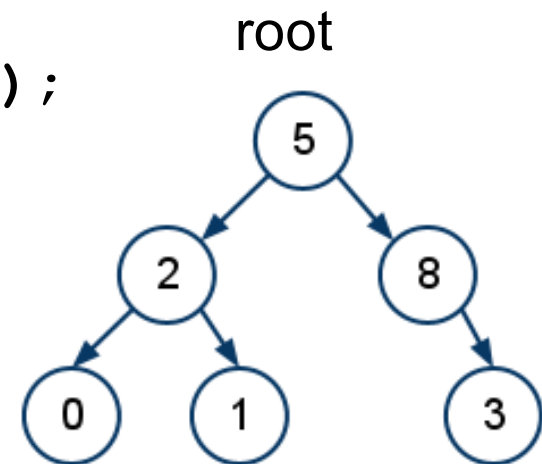
Sum of keys in leaf nodes

# Quiz 4: What is the output?

---

```
void m() {  
    Queue<Node> q = new LinkedList();  
    if(root==null) return;  
    q.offer(root);  
    while(!q.isEmpty()) {  
        Node t = q.poll();  
        if(t.right!=null) q.offer(t.right);  
        if(t.left!=null) q.offer(t.left);  
        print(t.key+",");  
    }  
}
```

m();



- A. 5,2,8,0,1,3
- B. 5,8,2,3,1,0
- C. 0,1,3,2,8,5
- D. 3,1,0,8,2,5



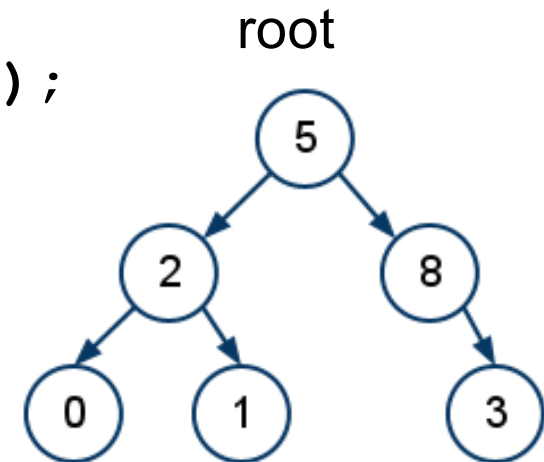
# Quiz 4: What is the output?

---

```
void m() {  
    Queue<Node> q = new LinkedList();  
    if(root==null) return;  
    q.offer(root);  
    while(!q.isEmpty()) {  
        Node t = q.poll();  
        if(t.right!=null) q.offer(t.right);  
        if(t.left!=null) q.offer(t.left);  
        print(t.key+",");  
    }  
}
```

m();

- A. 5,2,8,0,1,3
- B. 5,8,2,3,1,0**
- C. 0,1,3,2,8,5
- D. 3,1,0,8,2,5

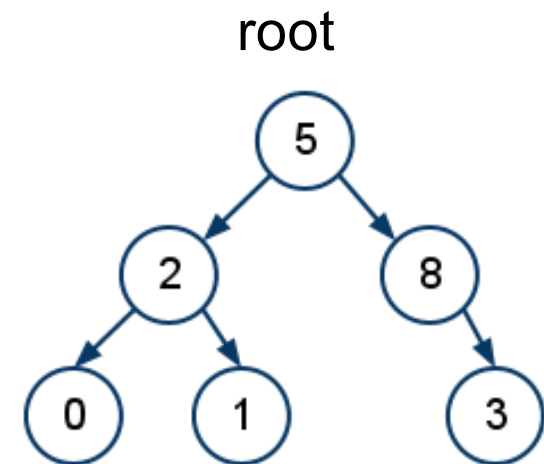


Level order right to left

# Quiz 5: What is the output?

---

```
void m(Node r, int n, int c){  
    if(r==null) return;  
    if(c == n){  
        System.out.print(r.key+",");  
    }  
    m(r.left, n, c+1);  
    m(r.right, n, c+1);  
}  
m(root,3,1);
```

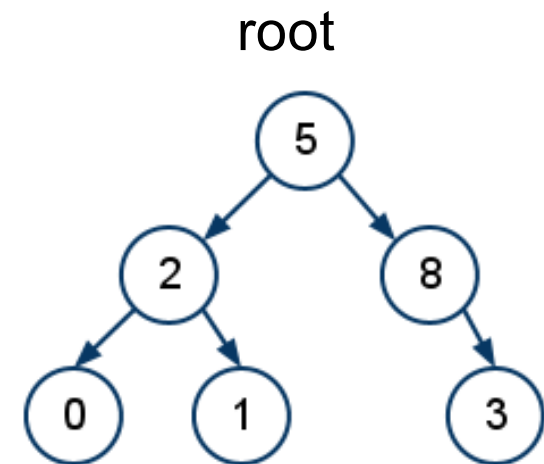


- A. 2,8
- B. 0,1,3
- C. 5,2,8,0,1,3
- D. 5,2,0,1,8,3

# Quiz 5: What is the output?

---

```
void m(Node r, int n, int c){  
    if(r==null) return;  
    if(c == n){  
        System.out.print(r.key+",");  
    }  
    m(r.left, n, c+1);  
    m(r.right, n, c+1);  
}  
m(root,3,1);
```



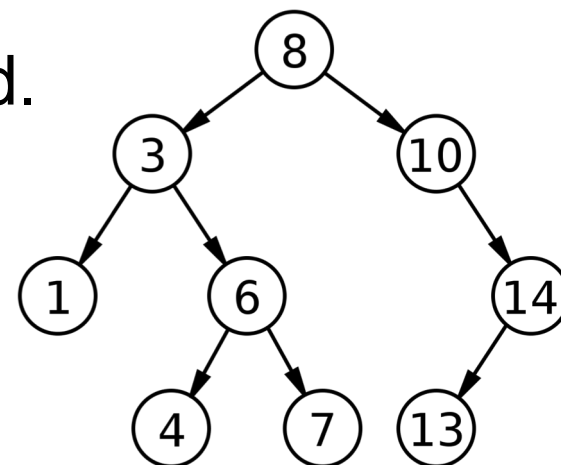
- A. 2,8
- B. 0,1,3
- C. 5,2,8,0,1,3
- D. 5,2,0,1,8,3

Print nodes at level n

# Binary Search Tree

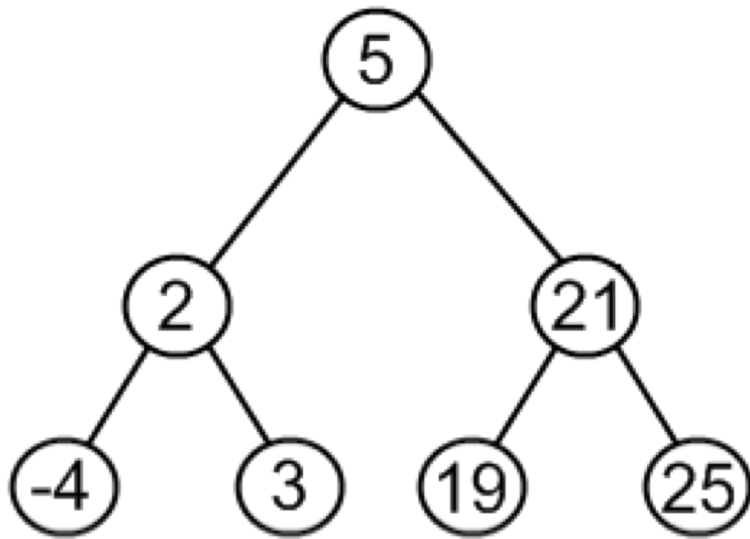
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- A BST is a binary tree
- Nodes are ordered in the following way:
  - Each node contains one key (also known as data)
  - Keys in the left subtree are less than the key in its parent node
  - Keys in the right subtree are greater than the key in its parent node
  - Duplicate keys are not allowed.

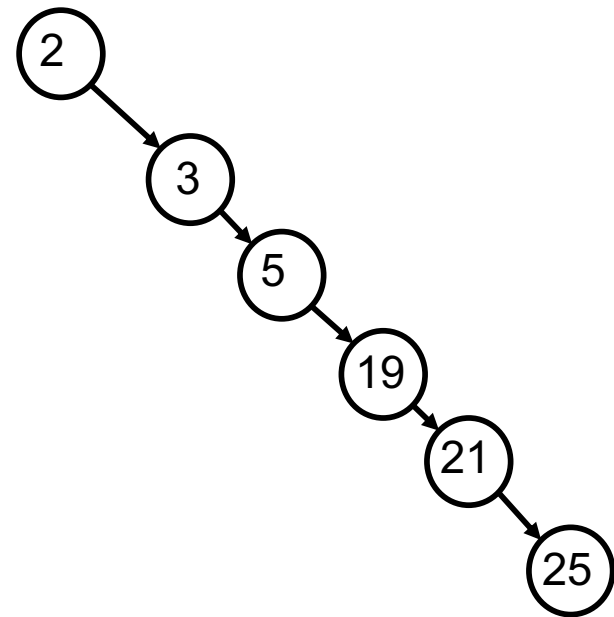


# Balanced vs Not Balanced BST

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**Balanced**  
**Height:  $O(\log n)$**



**Not Balanced**  
**Height:  $O(n)$**

# BST Class

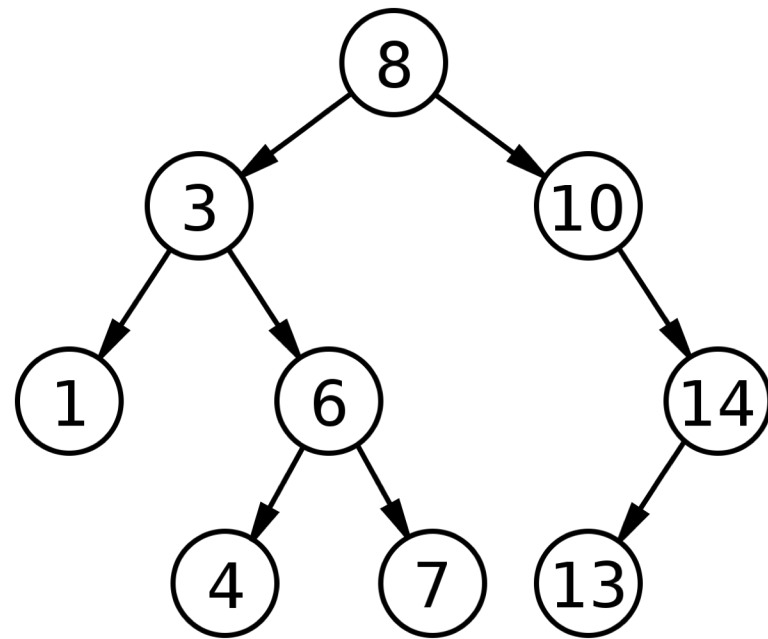
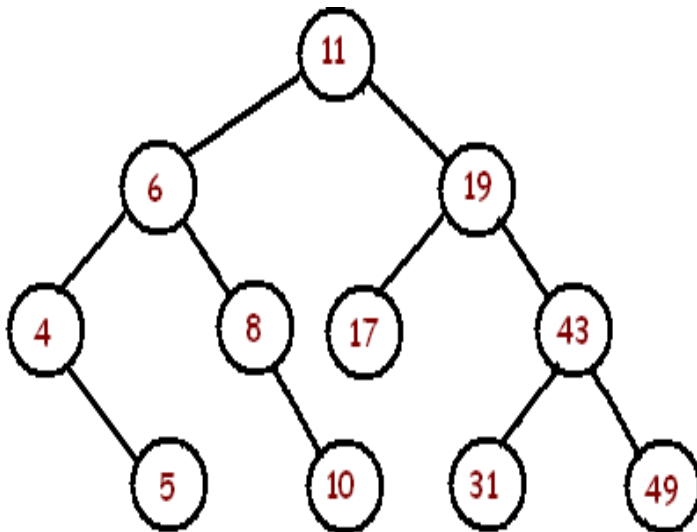
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```
public class BST<Key extends Comparable<Key>, Value>{
    private Node root;
    private class Node{
        private Key key;
        private Value value;
        private Node left, right;
        public Node(Key k, Value v){
            key = k;
            value = v;
        }
    }
}
```

# Search

---

- Since a binary search tree with  $n$  nodes has a minimum of  $O(\log n)$  levels, it takes at least  $O(\log n)$  comparisons to find a particular node.
- Unfortunately, a binary search tree can degenerate to a linked list, reducing the search time to  $O(n)$ .



# Search

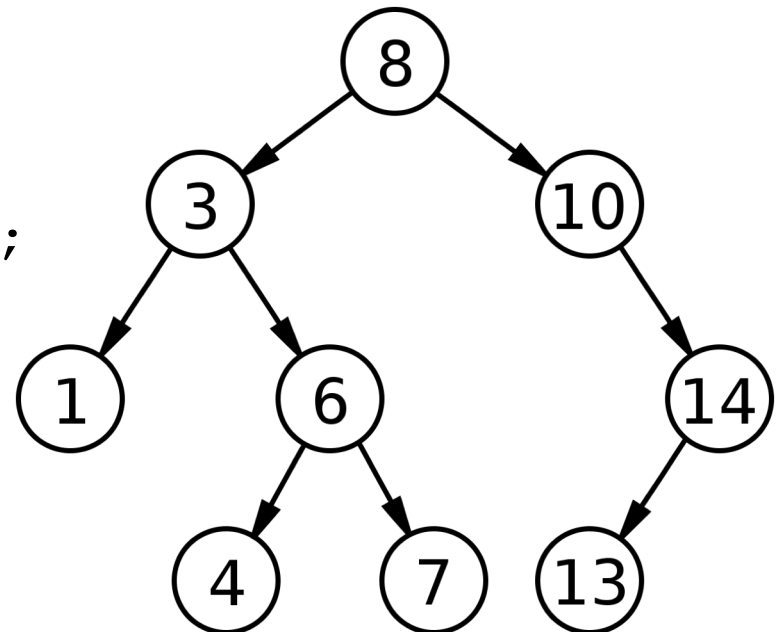
---

```
public Value get(Key key){  
    return get(root, key);  
}
```

```
private Value get(Node x, Key key){  
    if(x == null) return null;  
    int cmp = key.compareTo(x.key);  
    if(cmp < 0)  
        return get(x.left, key);  
    else if(cmp > 0)  
        return get(x.right, key);  
    else return x.value;  
}
```

Average:  $O(\log n)$

Worst:  $O(n)$



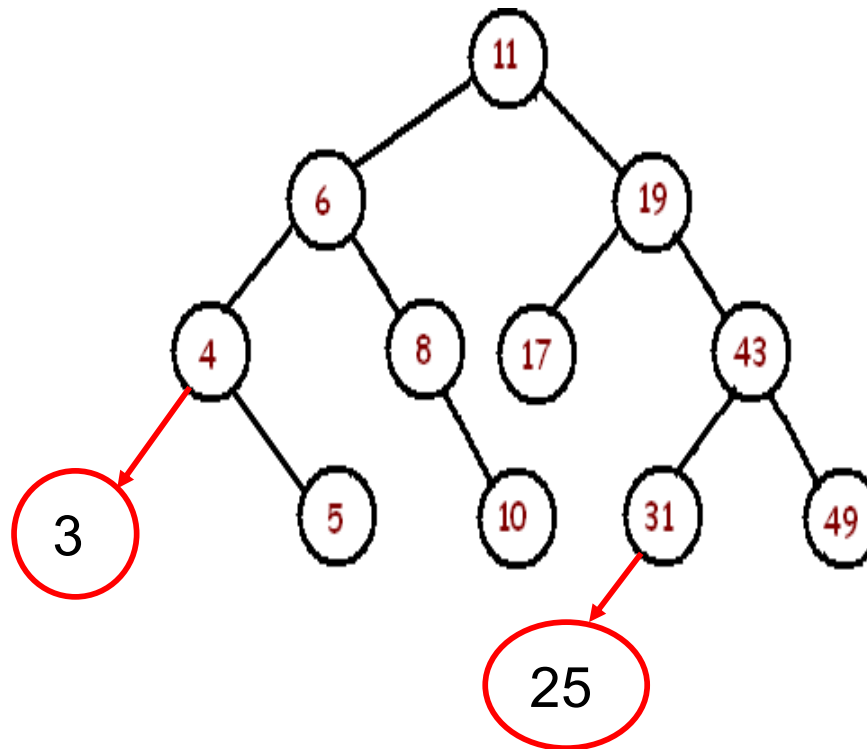


# Insert

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- ▶ The insertion procedure is quite similar to searching.
- ▶ We start at the root and recursively go down the tree searching for a location in a BST to insert a new node.
- ▶ New node always becomes a leaf node

insert(3)  
insert(25)



# Insert

---

```
public void put(Key key, Value val){
    root = put(root, key, val);
}

private Node put(Node x, Key key, Value val){
    if(x == null) {return new Node(key,val);}
    int cmp = key.compareTo(x.key);
    if(cmp < 0) {x.left = put(x.left, key, val);}
    else if(cmp>0) {x.right = put(x.right, key, val);}
    else {x.value = val;}
    return x;
}
```

# Delete

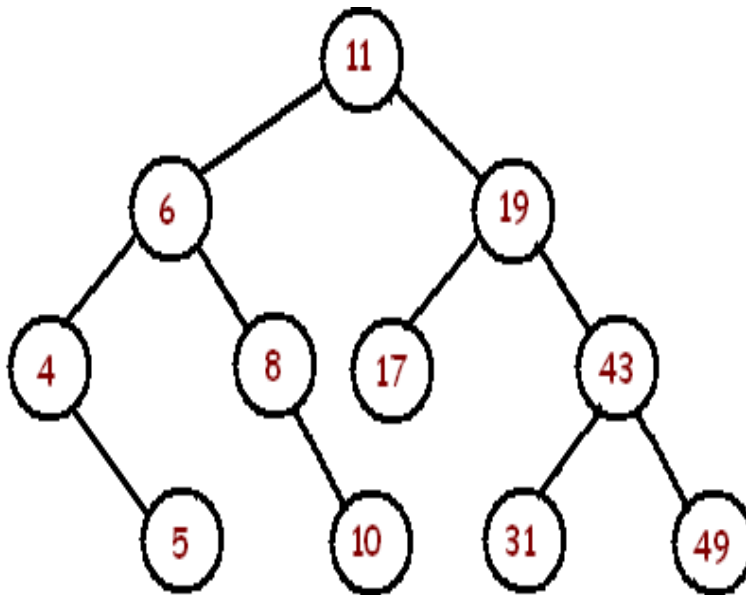
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- ▶ There are several cases to consider.
- ▶ A node to be deleted:
  - is not in a tree: there is nothing to delete
  - is a leaf: remove the node
  - has only one child: same as deleting a node from a linked list
  - has two children:
    - Replace the node with:
      - Largest of left subtree
      - Smallest of Right subtree

# Delete Example

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Delete 5

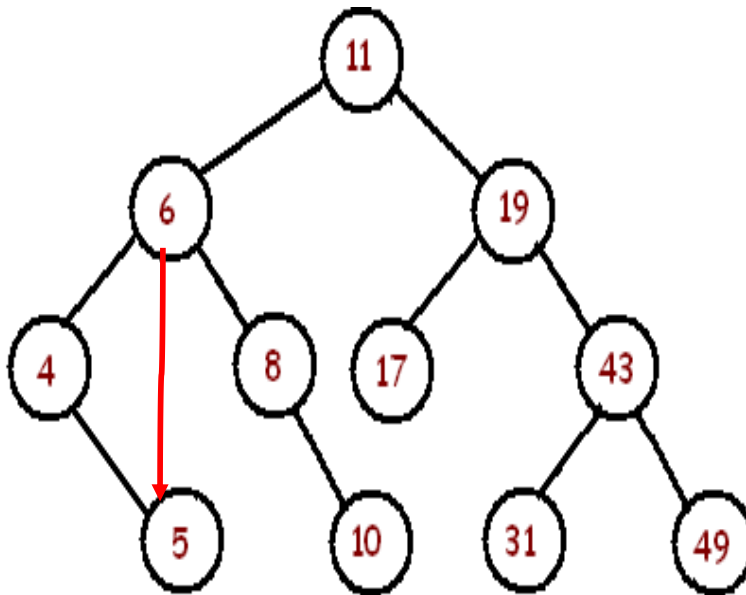


5 is a leaf node. It can be removed directly.

# Delete Example

---

Delete 8

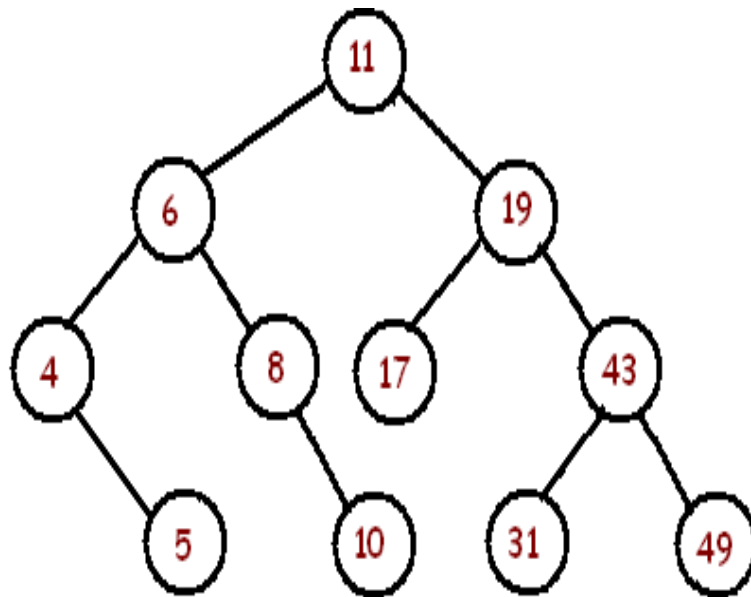


Node 8 has only one child. Parent, Node 6, will point to node 10, the child of node 8.

# Delete Example

---

Delete 11



Node 10, the max of left subtree  
or Node 17, the min of right  
subtree will replace Node 11.

# Delete

---

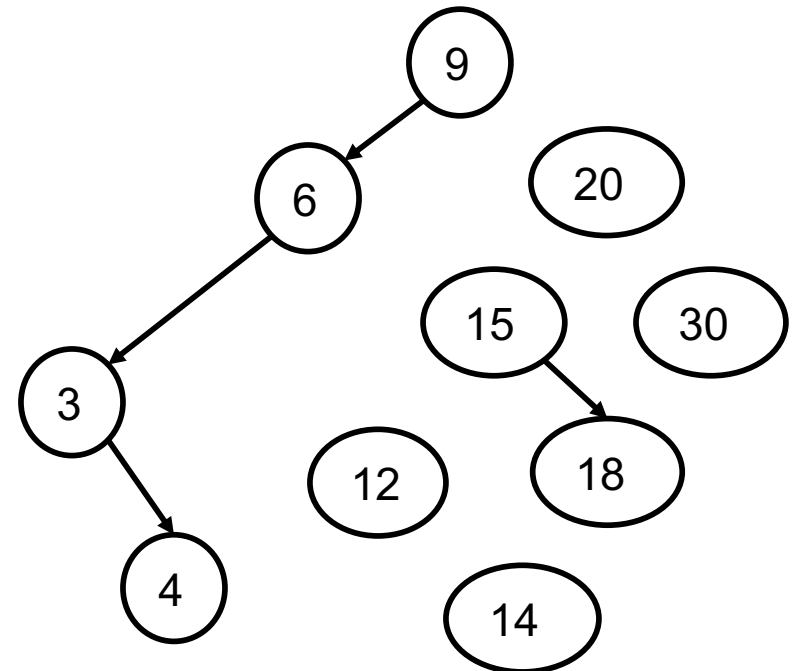
```
private Node delete(Node x, Key k) {
    if(x == null) return null;
    int cmp = k.compareTo(x.key);
    if(cmp < 0) x.left = delete(x.left, k);
    else if (cmp > 0) x.right = delete(x.right, k);
    else{
        if(x.right == null) return x.left;
        if(x.left == null) return x.right;
        Node t = x;
        x = min(t.right);
        x.right = deleteMin(t.right);
        x.left = t.left;
    }
    return x;
}
```

# Delete Example

---

```
private Node delete(Node x, Key k){
    if(x == null) return null;
    int cmp = k.compareTo(x.key);
    if(cmp < 0) x.left = delete(x.left, k);
    else if (cmp > 0) x.right = delete(x.right, k);
    else{
        if(x.right == null) return x.left;
        if(x.left == null) return x.right;
        Node t = x;
        x = min(t.right);
        x.right = deleteMin(t.right);
        x.left = t.left;
    }
    return x;
}
```

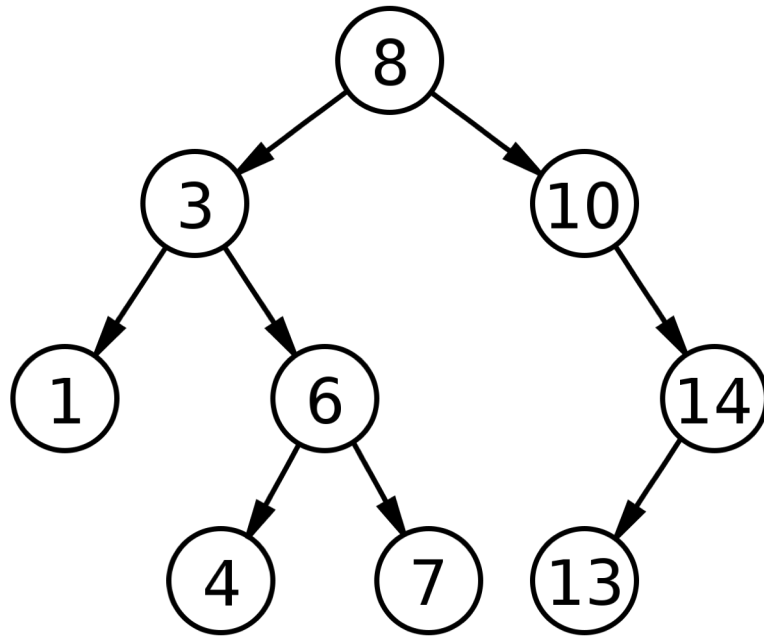
Delete 9





# isBST

---



# isBST

---

- ▶ 1) inOrder , check if sorted
- ▶ 2)

```
boolean isBST(Node x, Key min, Key max) {  
    if (x == null) return true;  
    if (min != null && x.key.compareTo(min) <= 0)  
        return false;  
    if (max != null && x.key.compareTo(max) >= 0)  
        return false;  
    return isBST(x.left, min, x.key) &&  
        isBST(x.right, x.key, max);  
}
```

```
isBST(root, null, null);
```

# Find min

---

# Find max

---

# Construct a BST

---

- ▶ **Construct a BST from given preorder traversal**
- ▶ preOrder: {10, 5, 1, 7, 40, 50}

# LCA (Least Common Ancestor)

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$\text{LCA}(5, 10) = 6$

$5 < 6 < 10$

