CMSC 330: Organization of Programming Languages

Functional Programming with Lists

Lists in OCaml

- The basic data structure in OCaml
  - Lists can be of *arbitrary length*
    - Implemented as a linked data structure
  - Lists must be *homogeneous*
    - All elements have the same type

- Operations
  - Construct lists
  - Destruct them via pattern matching

Constructing Lists

Syntax

- `[]` is the empty list (pronounced "nil")
- ***e1*::*e2*** prepends element ***e1*** to list ***e2***
  - Operator `::` is pronounced "cons"
  - ***e1*** is the head, ***e2*** is the tail
- **[*e1*;*e2*;…;*en*]** is *syntactic sugar* for **_e1_::_e2_::…::_en_::[]**

> Both *cons* and *nil* are terms from LISP

Examples

```
3::[]          (* The list [3] *)
2::(3::[])     (* The list [2; 3] *)
[1; 2; 3]      (* The list 1::(2::(3::[])) *)
```

Constructing Lists

Evaluation

- `[]` is a value
- To evaluate `[`***e1***`;…;`***en***`]`, evaluate ***e1*** to a value ***v1***, …., evaluate ***en*** to a value ***vn***, and return `[`***v1***`;…;`***vn***`]`

Examples

```
# let y = [1; 1+1; 1+1+1] ;;
val y : int list = [1; 2; 3]

# let x = 4::y ;;
val x : int list = [4; 1; 2; 3]

# let z = 5::y ;;
val z : int list = [5; 1; 2; 3]

# let m = “hello”::”bob”::[];;
val m : string list = [“hello”; “bob”]
```

Typing List Construction

Nil:

`[]: 'a list`

Polymorphic type: like a generic type in Java

i.e., empty list has type ***t*** **list** for any type ***t***

Cons:

If ***e1*** : ***t*** and ***e2*** : ***t*** **list** then ***e1*** **::** ***e2*** : ***t*** **list**

With parens for clarity:

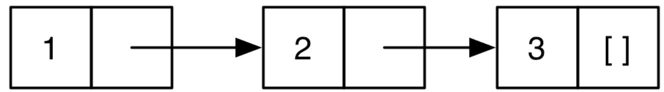
If ***e1*** : ***t*** and ***e2*** : (***t*** **list**) then (***e1*** **::** ***e2***) : (***t*** **list**)

Examples

```
# let x = [1;"world"] ;;
This expression has type string but an expression was
  expected of type int

# let m = [[1];[2;3]];;
val y : int list list = [[1]; [2; 3]]

# let y = 0::[1;2;3] ;;
val y : int list = [0; 1; 2; 3]

# let w = [1;2]::y ;;
This expression has type int list but is here used with
  type int list list
```

- The left argument of :: is an element, the right is a list
- Can you construct a list y such that [1;2]::y makes sense?

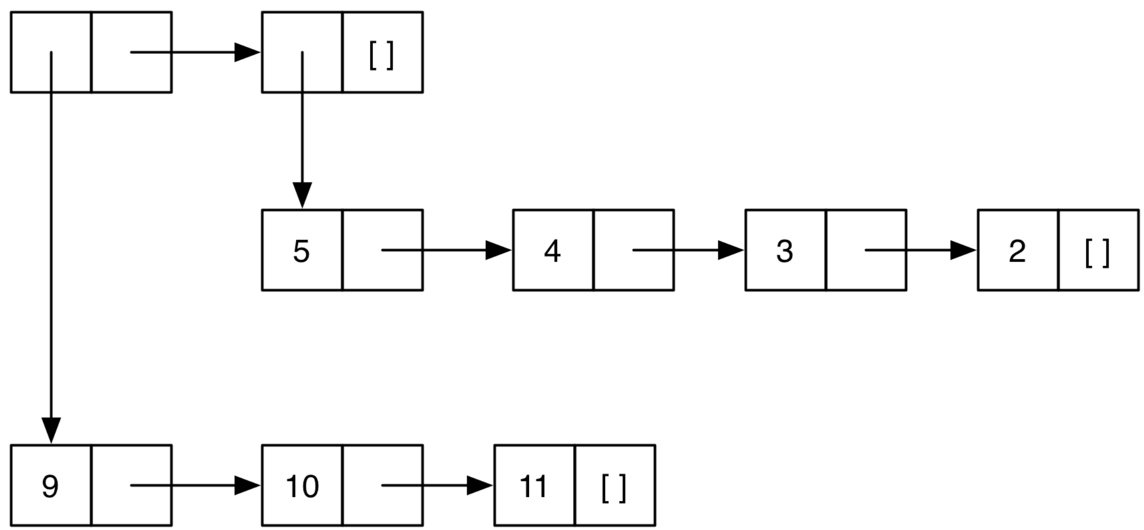
Lists in Ocaml are Linked



- [1;2;3] is represented above
  - A nonempty list is a pair (element, rest of list)
  - The element is the head of the list
  - The pointer is the tail or *rest* of the list
    - ...which is itself a list!
- Thus in math (i.e., inductively) a list is either
  - The empty list [ ]
  - Or a pair consisting of an element and a list
    - This recursive structure will come in handy shortly

Lists of Lists

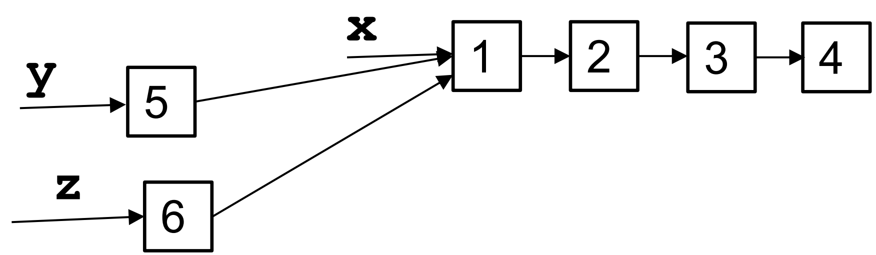
- Lists can be nested arbitrarily
  - Example: `[ [9; 10; 11]; [5; 4; 3; 2] ]`
    - (Type `int list list`)



Lists are Immutable

- No way to *mutate* (change) an element of a list
- Instead, build up new lists out of old, e.g., using **::**

```
let x = [1;2;3;4]
let y = 5::x
let z = 6::x
```



Quiz 1

What is the type of the following expression?

```
[1.0; 2.0; 3.0; 4.0]
```

A. `array`

B. `list`

C. `float list`

D. `int list`

Quiz 1

What is the type of the following expression?

```
[1.0; 2.0; 3.0; 4.0]
```

A. **`array`**
B. **`list`**
C. **`float list`**
D. **`int list`**

Quiz 2

What is the type of the following expression?

```
31::[3]
```

A. `int`
B. `int list`
C. `int list list`
D. `error`

Quiz 2

What is the type of the following expression?

```
31::[3]
```

A. **`int`**

B. **`int list`**

C. **`int list list`**

D. **`error`**

Quiz 3

What is the type of the following definition?

```
let f x = “alien”::[x]
```

A. `string -> string`

B. `string list`

C. `string list -> string list`

D. `string -> string list`

Quiz 3

What is the type of the following definition?

```
let f x = “alien”::[x]
```

A. `string -> string`

B. `string list`

C. `string list -> string list`

D. `string -> string list`

Pattern Matching

- To pull lists apart, use the **match** construct
- Syntax

```
match e with
| p1 -> e1
| …
| pn -> en
```

- Evaluate ***e*** to a value ***v***
- If ***p1*** matches ***v***, then evaluate ***e1*** to ***v1*** and return ***v1***

...

- Else if ***pn*** matches ***v***, then evaluate ***en*** to ***vn*** and return ***vn***
- Else, no patterns match: raise **Match_failure** exception

Pattern Matching Example

```
let is_empty l =
  match l with
    [] -> true
  | (h::t) -> false
```

- Example runs
  - `is_empty []    (* evaluates to true  *)`
  - `is_empty [1]   (* evaluates to false *)`
  - `is_empty [1;2](* evaluates to false *)`

Pattern Matching Example (cont.)

```
let hd l =
  match l with
   (h::t) -> h
```

- Example runs
  - `hd [1;2;3]` `(* evaluates to 1 *)`
  - `hd [2;3]` `(* evaluates to 2 *)`
  - `hd [3]` `(* evaluates to 3 *)`
  - `hd []` `(* Exception: Match_failure *)`

Quiz 4

To what does the following expression evaluate?

```
match [1;2;3] with
    [] -> [0]
  | h::t -> t
```

A. `[]`
B. `[0]`
C. `[1]`
D. `[2;3]`

Quiz 4

To what does the following expression evaluate?

```
match [1;2;3] with
    [] -> [0]
  | h::t -> t
```

A. `[]`
B. `[0]`
C. `[1]`
D. `[2;3]`

"Deep" pattern matching

- You can nest patterns for more precise matches
  - **a::b** matches lists with **at least one** element
    - Matches **[1;2;3]**, binding **a** to **1** and **b** to **[2;3]**
  - **a::[]** matches lists with **exactly one** element
    - Matches **[1]**, binding **a** to **1**
    - Could also write pattern **a::[]** as **[a]**
  - **a::b::[]** matches lists with **exactly two** elements
    - Matches **[1;2]**, binding **a** to **1** and **b** to **2**
    - Could also write pattern **a::b::[]** as **[a;b]**
  - **a::b::c::d** matches lists with **at least three** elements
    - Matches **[1;2;3]**, binding **a** to **1**, **b** to **2**, **c** to **3**, and **d** to **[]**
    - *Cannot* write pattern as **[a;b;c]::d** (why?)

Pattern Matching – Wildcards

- An underscore `_` is a wildcard pattern
  - Matches anything
  - But doesn't add any bindings
  - Useful to hold a place but discard the value
    - i.e., when the variable does not appear in the branch expression

- In previous examples
  - Many values of `h` or `t` ignored
  - Can replace with wildcard `_`

Pattern Matching – Wildcards (cont.)

- Code using _
  - `let is_empty l = match l with`
    `[] -> true   | (_::_) -> false`
  - `let hd l = match l with (h::_) -> h`
  - `let tl l = match l with (_::t) -> t`
- Outputs
  - `is_empty[1](* evaluates to false *)`
  - `is_empty[ ](* evaluates to true  *)`
  - `hd [1;2;3] (* evaluates to 1     *)`
  - `hd [1]     (* evaluates to 1     *)`
  - `tl [1;2;3] (* evaluates to [2;3] *)`
  - `tl [1]     (* evaluates to [ ]   *)`

Quiz 5

To what does the following expression evaluate?

```
match [1;2;3] with
  | 1::[]     -> [0]
  | _::_      -> [1]
  | 1::_::[] -> []
```

A. `[]`
B. `[0]`
C. `[1]`
D. `[2;3]`

Quiz 5

To what does the following expression evaluate?

```
match [1;2;3] with
  | 1::[]     -> [0]
  | _::_      -> [1]
  | 1::_::[] -> []
```

A. `[]`
B. `[0]`
C. `[1]`
D. `[2;3]`

Pattern Matching – An Abbreviation

- `let f p = e`, where *p* is a pattern
  - is shorthand for `let f x = match x with p -> e`
- Examples
  - `let hd (h::_) = h`
  - `let tl (_::t) = t`
  - `let f (x::y::_) = x + y`
  - `let g [x; y] = x + y`
- Useful if there's only one acceptable input

Pattern Matching Typing

```
match e with
| p1 -> e1
| …
| pn -> en
```

- If ***e*** and ***p1***, …, ***pn*** each have type ***ta***
- and ***e1***, …, ***en*** each have type ***tb***
- Then entire **`match`** expression has type ***tb***

- <u>Examples</u>

type: 'a list -> 'a

```
let hd l =
  match l with
  (h::_) -> h
```

ta = 'a list

tb

tb = 'a

type: int list -> int

```
let rec sum l =
  match l with
    [] -> 0
  | (h::t) -> h+sum t
```

tb

ta = int list

tb = int

Polymorphic Types

- The **`sum`** function works only for **`int lists`**
- But the **`hd`** function works for *any type of list*
  - `hd [1; 2; 3]                  (* returns 1 *)`
  - `hd ["a"; "b"; "c"]            (* returns "a" *)`
- OCaml gives such functions polymorphic types
  - `hd : 'a list -> 'a`
  - this says the function takes a list of *any* element type `'a`, and returns something of that same type
- These are basically generic types in Java
  - `'a list` is like `List<T>`

Examples Of Polymorphic Types

- `let tl (_::t) = t`

```
# tl [1; 2; 3];;
- : int list = [2; 3]
# tl [1.0; 2.0];;
- : float list = [2.0]
(* tl : 'a list -> 'a list *)
```

- `let fst x y = x`

```
# fst 1 “hello”;;
- : int = 1
# fst [1; 2] 1;;
- : int list = [1; 2]
(* fst : 'a -> 'b -> 'a *)
```

Examples Of Polymorphic Types

- `let eq x y = x = y    (* let eq x y = (x = y) *)`

```
# eq 1 2;;
- : bool = false
# eq "hello" "there";;
- : bool = false
# eq "hello" 1    -- type error
(* eq : 'a -> 'a -> bool *)
```

Quiz 6

What is the type of the following function?

```
let f x y =
  if x = y then 1 else 0
```

A. `'a -> 'b -> int`

B. `'a -> 'a -> bool`

C. `'a -> 'a -> int`

D. `int`

Quiz 6

What is the type of the following function?

```
let f x y =
  if x = y then 1 else 0
```

A. `'a -> 'b -> int`

B. `'a -> 'a -> bool`

C. `'a -> 'a -> int`

D. `int`

Pattern matching is ***AWESOME***

1. You can't forget a case
   - Compiler issues inexhaustive pattern-match warning
2. You can't duplicate a case
   - Compiler issues unused match case warning
3. You can't get an exception
   - Can't do something like `List.hd []`
4. Pattern matching leads to elegant, concise, beautiful code

Lists and Recursion

- Lists have a recursive structure
  - And so most functions over lists will be recursive

```
let rec length l = match l with
    [] -> 0
  | (_::t) -> 1 + (length t)
```

  - This is just like an inductive definition
    - *The length of the empty list is zero*
    - *The length of a nonempty list is 1 plus the length of the tail*
  - Type of length?
    - `‘a list -> int`

More Examples

- `sum l (* sum of elts in l *)`

```
let rec sum l = match l with
    [] -> 0
  | (x::xs) -> x + (sum xs)
```

- `negate l (* negate elements in list *)`

```
let rec negate l = match l with
    [] -> []
  | (x::xs) -> (-x) :: (negate xs)
```

- `last l  (* last element of l *)`

```
let rec last l = match l with
    [x] -> x
  | (x::xs) -> last xs
```

More Examples (cont.)

```
(* return a list containing all the elements in the list l
  followed by all the elements in list m *)
```

- `append l m`

```
let rec append l m = match l with
    [] -> m
  | (x::xs) -> x::(append xs m)
```

- `rev l  (* reverse list; hint: use append *)`

```
let rec rev l = match l with
     [] -> []
   | (x::xs) -> append (rev xs) [x]
```

- **`rev`** takes $O(n^2)$ time. Can you do better?