CMSC 132: Object-Oriented Programming II

Red & Black Tree

CMSC 132 Summer 2021

BST



Balanced BST Search: O(Log n)

Unbalanced BST Search: O(n)

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What is the worst case time complexity for search, insert and delete operations in a general Binary Search Tree?

- A. O(n) for all
- B. O(Logn) for all
- C. O(Logn) for search and insert, and O(n) for delete
- D. O(Logn) for search, and O(n) for insert and delete

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To delete a node X with 2 non-null children in a BST, we replace the node X with the minimum node Y from X's right subtree. Which of the following is true about the node Y?

- A. Y is always a leaf node
- B. Y is always either a leaf node or a node with empty left child
- C. Y may be an ancestor of the node
- D. Y is always either a leaf node or a node with empty right child

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We are given a set of n distinct elements and an unlabeled binary tree with n nodes. In how many ways can we populate the tree with the given set so that it becomes a binary search tree?

A. 0
B. 1
C. n!
D. n²

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Which of the following traversal outputs the data in sorted order in a BST?

- A. Preorder
- B. Inorder
- C. Postorder
- D. Level order

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Numbers 7, 5, 1, 8, 3, 6, 0, 9, 4, 2 are inserted in that order into an initially empty binary search tree. What is the in-order traversal sequence of the resultant tree?

- A. 7510324689
 B. 0243165987
- C. 0123456789
- D. 9864230157

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The following numbers are inserted into an empty binary search tree in the given order: 10, 1, 3, 5, 15, 12, 16. What is the height of the binary search tree (the height is the maximum distance of a leaf node from the root)?

A. 2B. 3C. 4D. 6

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The preorder traversal sequence of a binary search tree is 30, 20, 10, 15, 25, 23, 39, 35, 42. Which one of the following is the postorder traversal sequence of the same tree?

A. 10, 20, 15, 23, 25, 35, 42, 39, 30
B. 15, 10, 25, 23, 20, 42, 35, 39, 30
C. 15, 20, 10, 23, 25, 42, 35, 39, 30
D. 15, 10, 23, 25, 20, 35, 42, 39, 30

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A. 10, 20, 15, 23, 25, 35, 42, 39, 30
B. 15, 10, 25, 23, 20, 42, 35, 39, 30
C. 15, 20, 10, 23, 25, 42, 35, 39, 30
D. 15, 10, 23, 25, 20, 35, 42, 39, 30



Which of the following traversals is sufficient to construct BST from given traversals 1) Inorder 2) Preorder 3) Postorder

- A. Any one of the given three traversals is sufficient
- B. Either 2 or 3 is sufficient
- C. 2 and 3
- D. 1 and 3

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Balanced Binary Search Tree

- Red & Black Tree
- AVL Tree
- ▶ 2-3 Tree
- B-tree: Databases

Red & Black Tree

- A BST such that:
 - Tree edges have color: Red or Black
 - No node has two red edges connected to it.
 - Every path from root to null link has the same number of black links.
 - Red links lean left. (LLRB)
 - New node edge is Red



Search: red-black BSTs

 Observation. Search is the same as for elementary BST (ignore color).

```
public Val get(Key key)
{
    Node x = root;
    while (x != null)
    {
        int cmp = key.compareTo(x.key);
        if (cmp < 0) x = x.left;
        else if (cmp > 0) x = x.right;
        else if (cmp == 0) return x.val;
    }
    return null;
}
```



Red-black BST representation

```
private static final boolean RED
                                    = true;
private static final boolean BLACK = false;
private class Node
Ł
   Key key;
   Value val;
   Node left, right;
   boolean color; // color of parent link
}
private boolean isRed(Node x)
{
   if (x == null) return false;
   return x.color == RED;
}
                              null links are black
```



Elementary Operations

 Left rotation. Orient a (temporarily) right-leaning red link to lean left.



rotate E left (before)

rotate E left (after)

 Left rotation. Orient a (temporarily) right-leaning red link to lean left.



```
private Node rotateLeft(Node h)
{
    assert isRed(h.right);
    Node x = h.right;
    h.right = x.left;
    x.left = h;
    x.color = h.color;
    h.color = RED;
    return x;
}
```

Right rotation: Orient a left-leaning red link to (temporarily) lean right.



rotate E left (before)

rotate E left (after)

Right rotation: Orient a left-leaning red link to (temporarily) lean right.



```
private Node rotateRight(Node h)
{
    assert isRed(h.left);
    Node x = h.left;
    h.left = x.right;
    x.right = h;
    x.color = h.color;
    h.color = RED;
    return x;
}
```

Color flip:



Color flip(before)

Color flip (after)

Color flip.



```
private void flipColors(Node h)
{
    assert !isRed(h);
    assert isRed(h.left);
    assert isRed(h.right);
    h.color = RED;
    h.left.color = BLACK;
    h.right.color = BLACK;
}
```

Insertion in a LLRB tree

- Right child red, left child black: rotate left.
- Left child, left-left grandchild red: rotate right.
- Both children red: flip colors.



Insertion

```
Node put(Node h, Key key, Value val) {
  if (h == null) return new Node(key, val, RED, 1);
  int cmp = key.compareTo(h.key);
  if (cmp < 0) h.left = put(h.left, key, val);
  else if (cmp > 0) h.right = put(h.right, key, val);
  else h.val = val;
```

```
// fix-up any right-leaning links
if (isRed(h.right) && !isRed(h.left)) h = rotateLeft(h);
if (isRed(h.left) && isRed(h.left.left)) h = rotateRight(h);
if (isRed(h.left) && isRed(h.right)) flipColors(h);
return h;
```



}



Insert: M, X



Insert: M, X



Rotate Left

Insert: M, D, X



Flip Color

Insert: M, D, B



Rotate Right

Insert: M, D, B





M, D, B, H Insert S



M, D, B,H Insert S



Color flip

M, D, B,H Insert S













M, D, B,H,S,E,K



Color flip

R&B tree vs. BST

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