

CMSC 330: Organization of Programming Languages

Lets, Tuples, Records

Let Expressions

- Syntax
 - `let x = $e1$ in $e2$`
 - x is a *bound variable*
 - $e1$ is the *binding expression*
 - $e2$ is the *body expression*
- `let` expressions bind *local* variables
 - Different from `let definitions`, which are at the top-level

Let Expressions

- Syntax

- let $x = e1$ in $e2$

- Evaluation

- Evaluate $e1$ to $v1$
 - Substitute $v1$ for x in $e2$
 - yielding new expression $e2'$
 - Evaluate $e2'$ to $v2$, the final result

Example

let $z = 3+4$ in $3*z$

(evaluate $e1$)

➤ let $z = 7$ in $3*z$

(substitute for var z in $e2$)

➤ $3*7$

(compute the final result)

➤ 21

Let Expressions

- Syntax

- `let x = $e1$ in $e2$`

- Type checking

Example

What is the type of `let z = $3+4$ in $3*z$` ?

- `$3+4$: int`
- Assuming `z : int`, we have `$3*z$: int`
- So the type of `let z = $3+4$ in $3*z$` is `int`

Let Expressions

- Syntax

- `let $x = e1$ in $e2$`

- Type checking

- If `$e1 : t1$` and

- If assuming `$x : t1$` implies `$e2 : t$`

- Then `(let $x = e1$ in $e2$) : t`

Example

What is the type of `let $z = 3+4$ in $3*z$` ?

- `$3+4 : int$`
- Assuming `$z : int$` , we have `$3*z : int$`
- So the type of `let $z = 3+4$ in $3*z$` is `int`

Let Definitions vs. Let Expressions

- At the top-level, we write
 - `let x = e;;` (* no `in e2` part *)
 - This is called a let *definition*, not a let *expression*
 - Because it doesn't, itself, evaluate to anything
- Omitting `in` means “from now on”:
 - # `let pi = 3.14;;`
 - (* `pi` is now *bound* in the rest of the top-level scope *)

Top-level expressions

- We can write any expression at top-level, too

– `e;;`

– This says to evaluate `e` and then ignore the result

- Equivalent to `let _ = e;;`
- Useful when `e` has a side effect, such as reading/writing a file, printing to the screen, etc.

```
let x = 37;;  
let y = x + 5;;  
print_int y;;  
print_string "\n";;
```

- When run, outputs `42` to the screen

Let Expressions: Scope

- In `let x = e1 in e2`, var **x** is *not* visible outside of **e2**

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;  
print_float pi;;
```

error: **pi** not bound

bind **pi** (only) in body of **let**
(which is `pi *. 3.0 *. 3.0`)

Binding in other languages

- In `let x = e1 in e2`, var `x` is *not* visible outside of `e2`

```
let pi = 3.14 in pi *. 3.0 *. 3.0;;  
print_float pi;; (* pi unbound! *)
```

- Compare to similar usage in Java/C

```
{  
    float pi = 3.14;  
  
    pi * 3.0 * 3.0;  
}  
pi; /* pi unbound! */
```

Examples – Scope of Let bindings

- $x;;$
 - (* Unbound value x *)
- $\text{let } x = 1 \text{ in } x + 1;;$
 - (* 2 *)
- $\text{let } x = x \text{ in } x + 1;;$
 - (* Unbound value x *)

Examples – Scope of Let bindings

- `let x = 1 in (x + 1 + x) ;;`
 - `(* 3 *)`
- `(let x = 1 in x + 1) ;;`
`x;;`
 - `(* Unbound value x *)`
- `let x = 4 in (let x = x + 1 in x) ;;`
 - `(* 5 *)`

Nested Let Expressions

- Uses of `let` can be nested (last example on prev. slide)
 - Nested bound variables (`pi` and `r`) not visible outside
- Similar scoping possibilities C and Java

```
let res =  
  (let area =  
    (let pi = 3.14 in  
      let r = 3.0 in  
        pi *. r *. r) in  
    area /. 2.0);;
```

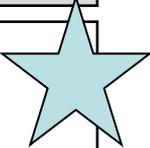
```
float res;  
{ float area;  
  { float pi = 3.14  
    float r = 3.0;  
    area = pi * r * r;  
  }  
  res = area / 2.0;  
}
```

Nested Let Style: Generally Avoid

- Sometimes a nested binding can be rewritten in a more linear style
 - Easier to understand
- Can go too far: namespace pollution
 - Avoiding adding unnecessary variable bindings to top-level

```
let res =  
  (let area =  
    (let pi = 3.14 in  
      let r = 3.0 in  
        pi *. r *. r) in  
    area /. 2.0);;
```

```
let res =  
  let pi = 3.14 in  
  let r = 3.0 in  
  let area = pi *. r *. r in  
  area /. 2.0;;
```



```
let pi = 3.14;;  
let r = 3.0;;  
let area = pi *. r *. r;;  
let res = area /. 2.0;;
```

Let Expressions in Functions

- You can use `let` inside of functions for local vars

```
let area r =  
  let pi = 3.14 in  
  pi *. r *. r
```

- And you can use many `lets` in sequence

```
let area d =  
  let pi = 3.14 in  
  let r = d /. 2.0 in  
  pi *. r *. r
```

Shadowing Names

- **Shadowing** is rebinding a name in an inner scope to have a different meaning
 - May or may not be allowed by the language

C

```
int i;

void f(float i) {
    {
        char *i = NULL;
        ...
    }
}
```

Java

```
void h(int i) {
    {
        float i; // not allowed
        ...
    }
}
```

OCaml

```
let x = 3;;
let g x = x + 3;;
```

Shadowing, by the Semantics

- What if **e2** is also a **let** for **x** ?
 - Substitution will **stop** at the **e2** of a shadowing **x**

Example

```
let x = 3+4 in let x = 3*x in x+1
```

```
➤ let x = 7 in let x = 3*x in x+1
```

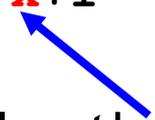
```
➤ let x = 3*7 in x+1
```

```
➤ let x = 21 in x+1
```

```
➤ 21+1
```

```
➤ 22
```

Will *not* be substituted,
since it is shadowed
by the inner let



Shadowing (of Locals) Discouraged

- You can use shadowing to *simulate* update (mutation)

```
let rec f x n =  
  if x = 0 then 1  
  else  
    let x = x - 1 in (* shadowed *)  
    n * (f x n)
```

- But **avoiding shadowing is clearer**
 - With no shadowing, if you see a variable **x**, you know it hasn't been “changed,” no matter where it appears
 - If you want to “update” **n**, use a new name **n1**, **n'**, etc.

Quiz 1

Which of these is **not** an expression that evaluates to 3?

A. `let x=2 in x+1`

B. `let x=3 in x`

C. `let x=3`

D. `3`

Quiz 1

Which of these is **not** an expression that evaluates to 3?

A. `let x=2 in x+1`

B. `let x=3 in x`

C. `let x=3 --->` not an expression

D. `3`

Quiz 2: What does this evaluate to?

```
let x = 2 in  
x = 3
```

- A. 3
- B. 2
- C. true
- D. false

Quiz 2: What does this evaluate to?

```
let x = 2 in  
x = 3
```

A. 3

B. 2

C. true

D. false

This expression is checking whether **x** is equal to 3

Quiz 3: What does this evaluate to?

```
let y = 3 in
let x = y+2 in
let y = 6 in
x+y
```

- A. 8
- B. 11
- C. 13
- D. 14

Quiz 3: What does this evaluate to?

```
let y = 3 in
let x = y+2 in
let y = 6 in
x+y
```

A. 8

B. 11

C. 13

D. 14

let Specializes match

More general form of `let` allows patterns:

- `let p = e1 in e2`
 - where *p* is a pattern. If *e1* fails to match that pattern then an exception is thrown

This pattern form of `let` is equivalent to

- `match e1 with p -> e2`

Examples

- `let [x] = [[1]] in 1::x (* evals to [1;1] *)`
- `let h::_ = [1;2;3] in h (* evaluates to 1 *)`
- `let () = print_int 5 in 3 (* evaluates to 3 *)`

Tuples

- **Constructed** using $(e1, \dots, en)$
- **Deconstructed** using pattern matching
 - Patterns involve parens and commas, e.g., $(p1, p2, \dots)$
- Tuples are similar to C structs
 - But without field labels
 - Allocated on the heap
- Tuples can be heterogenous
 - Unlike lists, which must be homogenous
 - $(1, ["string1"; "string2"])$ is a valid tuple

Tuple Types

- Tuple types use `*` to separate components
 - Type joins types of its components
- Examples
 - `(1, 2) :`
 - `(1, "string", 3.5) :`
 - `(1, ["a"; "b"], 'c') :`
 - `[(1,2)] :`
 - `[(1, 2); (3, 4)] :`
 - `[(1,2); (1,2,3)] :`

Tuple Types

- Tuple types use ***** to separate components
 - Type joins types of its components

- Examples

- (1, 2) : `int * int`
- (1, "string", 3.5) : `int * string * float`
- (1, ["a"; "b"], 'c') : `int * string list * char`
- [(1,2)] : `(int * int) list`
- [(1, 2); (3, 4)] : `(int * int) list`
- [(1,2); (1,2,3)] : `error`

Because the first list element has type `int * int`, but the second has type `int * int * int` – list elements must all be of the same type

Pattern Matching Tuples

```
# let plusThree t =  
  match t with  
    (x, y, z) -> x + y + z;;  
plusThree : int*int*int -> int = <fun>  
  
# let plusThree' (x, y, z) = x + y + z;;  
plusThree' : int*int*int -> int = <fun>  
  
# let addOne (x, y, z) = (x+1, y+1, z+1);;  
addOne : int*int*int -> int*int*int = <fun>  
  
# plusThree (addOne (3, 4, 5));;  
- : int = 15
```

Remember, **semicolon** for lists, **comma** for tuples

- `[1, 2] = [(1, 2)]` *which is a list of size one*
- `(1; 2)` *Warning: This expression should have type unit*

Tuples Are A Fixed Size

- This OCaml definition
 - `let foo x = match x with`
 - `(a, b) -> a + b`
 - `| (a, b, c) -> a + b + c`

has a type error. Why?

- Tuples of different size have different types
 - `(a, b)` has type: `'a * 'b`
 - `(a, b, c)` has type: `'a * 'b * 'c`
 - Patterns in the same `match` must have the same type

Quiz 4: What does this evaluate to?

```
let get a b = (a+b,0) in  
get 1 2
```

- A. (3,0)
- B. (2,0)
- C. 3
- D. type error

Quiz 4: What does this evaluate to?

```
let get a b = (a+b, 0) in  
get 1 2
```

- A. (3,0)
- B. (2,0)
- C. 3
- D. type error

Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in  
get (2,1) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

Quiz 5: What does this evaluate to?

```
let get (a,b) y = a+y in  
get (2,1) 1
```

- A. 3
- B. type error
- C. 2
- D. 1

Records

- Records: identify elements by **name**
 - Elements of a tuple are identified by **position**
- Define a **record type** before defining record values

```
type date = { month: string; day: int; year: int }
```

- Define a **record value**

```
# let today = { day=16; year=2017; month="f"^"eb" };;  
today : date = { day=16; year=2017; month="feb" };;
```

Destructing Records

```
type date = { month: string; day: int; year: int }  
let today = { day=16; year=2017; month="feb" };;
```

- **Access** by **field name** or **pattern matching**

```
print_string today.month;; (* prints feb *)  
(* patterns *)  
let { month=_; day=d } = today in  
let { year } = today in  
let _ = print_int d in      (* prints 16 *)  
print_int year;;           (* prints 2017 *)
```

- Notes:
 - In record patterns, you can skip or reorder fields
 - You can use the field name as the bound variable

Quiz 6: What is the type of `shift`?

```
type point = {x:int; y:int}
let shift { x=px } = [px]::[]
```

- A. `point -> int list`
- B. `int -> int list`
- C. `point -> point list`
- D. `point -> int list list`

Quiz 6: What is the type of `shift`?

```
type point = {x:int; y:int}
let shift { x=px } = [px]::[]
```

- A. `point -> int list`
- B. `int -> int list`
- C. `point -> point list`
- D. `point -> int list list`