

CMSC 330: Organization of Programming Languages

OCaml Data Types

OCaml Data

- So far, we've seen the following kinds of data
 - **Basic types** (int, float, char, string)
 - **Lists**
 - One kind of data structure
 - A list is either `[]` or `h::t`, deconstructed with pattern matching
 - **Tuples and Records**
 - Let you collect data together in fixed-size pieces
 - **Functions**
- How can we build other data structures?
 - Building everything from lists and tuples is awkward

User Defined Types

- `type` can be used to create new names for types
- Like `typedef` in C – a name might be more useful for communicating intent than just the type structure

User Defined Types

```
# type mylist = int*(int list);;
type mylist = int * int list

# let empty:mylist = (0, []);;
val empty : mylist = (0, [])

# let add x ((n,xs):mylist):mylist = (n+1,x::xs);;
val add : int -> mylist -> mylist = <fun>

# let length ((n,_):mylist) = n;;
val length : mylist -> int = <fun>

# let x = add 1 (add 2 empty);;
val x : mylist = (2, [1; 2])
```

Annotation required
to tell type inference
you want mylist,
not int*int list

(User-Defined) Variants

```
type coin = Heads | Tails
```

```
let flip x =
```

```
  match x with
```

```
    Heads -> Tails
```

```
  | Tails -> Heads
```

```
let rec count_heads x =
```

```
  match x with
```

```
    [] -> 0
```

```
  | (Heads::x') -> 1 + count_heads x'
```

```
  | (_::x') -> count_heads x'
```

In simplest form:
Like a C **enum**

Basic pattern
matching
resembles C
switch

Combined list
and variant
patterns
possible

Constructing and Destructing Variants

- Syntax

- **type** $t = C1 \mid \dots \mid Cn$
- the Ci are called **constructors**
 - Must begin with a capital letter

- Evaluation

- A constructor Ci is already a value
- Destructing a value v of type t is done by pattern matching on v ; the patterns are the constructors Ci

- Type Checking

- $Ci : t$ (for each Ci in t 's definition)

Data Types: Variants with Data

- We can define variants that “carry data” too
 - Not just a constructor, but a constructor *with values*

```
type shape =  
  Rect of float * float (* width*length *)  
  | Circle of float      (* radius *)
```

- **Rect** and **Circle** are constructors, so a **shape** is either
 - **Rect** (*w*, *l*) for any floats *w* and *l*, or
 - **Circle** *r* for any float *r*

Data Types: Pattern Matching

```
let area s =  
  match s with  
    | Rect (w, l) -> w *. l  
    | Circle r -> r *. r *. 3.14  
;;  
area (Rect (3.0, 4.0));; (* 12.0 *)  
area (Circle 3.0);;    (* 28.26 *)
```

- Use pattern matching to **deconstruct** values
 - Can bind pattern values to data parts

Data types are *aka* algebraic data types and tagged unions

Data Types: Pattern Matching

```
type shape =  
  Rect of float * float (* width*length *)  
  | Circle of float      (* radius *)  
  
let lst = [Rect (3.0, 4.0) ; Circle 3.0]
```

- What's the type of `lst`?
 - **shape list**
- What's the type of `lst`'s first element?
 - **shape**

Quiz 1

```
type foo = (int * (string list)) list
```

Which one of the following could match type `foo`?

- A. `[(3, "foo", "bar")]`
- B. `[(7, ["foo", "bar"])]`
- C. `[(5, ["foo"; "bar"])]`
- D. `[(9, [("foo", "bar")])]`

Quiz 1

```
type foo = (int * (string list)) list
```

Which one of the following could match type `foo`?

- A. `[(3, "foo", "bar")]`
- B. `[(7, ["foo", "bar"])]`
- C. `[(5, ["foo"; "bar"])]`
- D. `[(9, [("foo", "bar")])]]`

Quiz 2: What does this evaluate to?

```
type num = Int of int | Float of float;;  
  
let aux a =  
  match a with  
  | Int i -> float_of_int i  
  | Float j -> j  
;;  
  
aux (Int 2);;
```

- A. 4.0
- B. 2.0
- C. 2
- D. Type Error

Quiz 2: What does this evaluate to?

```
type num = Int of int | Float of float;;  
let aux a =  
  match a with  
  | Int i -> float_of_int i  
  | Float j -> j  
;;  
aux (Int 2);;
```

- A. 4.0
- B. 2.0
- C. 2
- D. Type Error

Variation: Shapes in Java

Compare this to OCaml

```
public interface Shape {  
    public double area();  
}
```

```
class Rect implements Shape {  
    private double width, length;  
  
    Rect (double w, double l) {  
        this.width = w;  
        this.length = l;  
    }  
  
    double area() {  
        return width * length;  
    }  
}
```

```
class Circle implements Shape {  
    private double rad;  
  
    Circle (double r) {  
        this.rad = r;  
    }  
  
    double area() {  
        return rad * rad * 3.14159;  
    }  
}
```

Option Type

```
type optional_int =  
  None  
  | Some of int  
  
let divide x y =  
  if y != 0 then Some (x/y)  
  else None  
  
let string_of_opt o =  
  match o with  
  Some i -> string_of_int i  
  | None -> "nothing"
```

```
let p = divide 1 0;;  
print_string  
  (string_of_opt p);;  
(* prints "nothing" *)  
  
let q = divide 1 1;;  
print_string  
  (string_of_opt q);;  
(* prints "1" *)
```

- Comparing to Java: **None** is like `null`, while **Some *i*** is like an **Integer (*i*)** object

Polymorphic Option Type

- A **Polymorphic** version of `option` type can work with *any kind of data*
 - As `int option`, `char option`, etc...

*Polymorphic parameter:
like `Option<T>` in Java*

```
type 'a option =  
  Some of 'a  
| None
```

```
let opthd l =  
  match l with  
  [] -> None  
  | x::_ -> Some x
```

In fact, this `option` type is built into OCaml

```
let p = opthd [];; (* p = None *)  
let q = opthd [1;2];; (* q = Some 1 *)  
let r = opthd ["a"];; (* r = Some "a" *)
```


Quiz 3: What does this evaluate to?

```
let foo f = match f with
  None -> 42.0
  | Some n -> n +. 42.0
;;
foo 3.3;;
```

- A. 45.3
- B. 42.0
- C. Some 45.3
- D. Error

Quiz 3: What does this evaluate to?

```
let foo f = match f with
  None -> 42.0
  | Some n -> n +. 42.0
;;
foo 3.3;;   foo (Some 3.3)
```

- A. 45.3
- B. 42.0
- C. Some 45.3
- D. Error**

Recursive Data Types

- We can build up lists with **recursive** variant types

```
type 'a mylist =  
  Nil  
  | Cons of 'a * 'a mylist  
  
let rec len x = match x with  
  Nil -> 0  
  | Cons (_, t) -> 1 + (len t)  
  
len (Cons (10, Cons (20, Cons (30, Nil))))  
(* evaluates to 3 *)
```

- Won't have nice `[1; 2; 3]` syntax for this kind of list

Variants (full definition)

• Syntax

- **type** $t = C1$ [of $t1$] | ... | Cn [of tn]
- the Ci are called **constructors**
 - Must begin with a capital letter; may include associated data - notated with brackets [] to indicate it's optional

• Evaluation

- A constructor Ci is a value if it has no assoc. data
 - $Ci\ vi$ is a value if it does
- Destructing a value of type t is by pattern matching
 - patterns are constructors Ci with data components, if any

• Type Checking

- Ci [vi] : t [if vi has type ti]

OCaml Exceptions

```
exception My_exception of int
let f n =
  if n > 0 then
    raise (My_exception n)
  else
    raise (Failure "foo")
let bar n =
  try
    f n
  with My_exception n ->
    Printf.printf "Caught %d\n" n
  | Failure s ->
    Printf.printf "Caught %s\n" s
```

OCaml Exceptions: Details

- Exceptions are declared with `exception`
 - They may appear in the signature as well
- Exceptions may take arguments
 - Just like type constructors
 - May also have no arguments
- Catch exceptions with `try...with...`
 - Pattern-matching can be used in `with`
 - If an exception is uncaught
 - Current function exits immediately
 - Control transfers up the call chain
 - Until the exception is caught, or until it reaches the top level

OCaml Exceptions: Useful Examples

- `failwith s`: Raises exception `Failure s` (s is a string).
- `Not_found`: Exception raised by library functions if the object does not exist
- `invalid_arg s`: Raises exception `Invalid_argument s`

```
let div x y =  
    if y = 0 then failwith "div by 0" else x/y;;
```

```
let lst = [(1, "alice"); (2, "bob"); (3, "cat")];;
```

```
let lookup key lst =
```

```
    try
```

```
        List.assoc key lst
```

```
    with
```

```
        Not_found -> "key does not exist"
```