# BILL AND EMILY RECORD LECTURE!!!!

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UNTIMED PART OF FINAL IS MONDAY May 10 9:00AM. NO DEAD CAT

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FINAL IS MONDAY May 17 8:00PM-10:15PM

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# FILL OUT COURSE EVALS for ALL YOUR COURSES!!!

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# Solving One Non-Linear Recurrences

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Assume that  $\boxdot$  is a non-associative operation that is computable (e.g., Subtraction).

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Someone gives you a, b and wants you to evaluate

#### a 🖸 b

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There is **no problem** with doing this.

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There is **a problem** with doing this. They may want

## $(x \boxdot y) \boxdot z \text{ or } x \boxdot (y \boxdot z)$

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Someone gives you x, y, z and wants you to evaluate

#### $x \boxdot y \boxdot z$

There is a **problem** with doing this. They may want

#### $(x \boxdot y) \boxdot z \text{ or } x \boxdot (y \boxdot z)$

Which one? Can't tell. They need to PARENTHESIZE.

 $x_1$ . One:  $(x_1)$ .



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 $x_1$ . One:  $(x_1)$ .  $x_1 \boxdot x_2$ . One:  $(x_1 \boxdot x_2)$ 

 $x_1$ . One:  $(x_1)$ .  $x_1 \boxdot x_2$ . One:  $(x_1 \boxdot x_2)$  $x_1 \boxdot x_2 \boxdot x_3$ . Two:  $(x_1 \boxdot x_2) \boxdot (x_3)$  and  $(x_1) \boxdot (x_2 \boxdot x_3)$ .

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# Can We Make A Recurrence?

Let  $a_n$  be the number of ways to parenthesize

 $x_1 \boxdot \cdots \boxdot x_n$ .

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## Can We Make A Recurrence?

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For every way to parenthesize there exist an  $1 \leq i \leq n$  such that it looks like

$$(x_1 \odot \cdots \odot x_i) \boxdot (x_{i+1} \boxdot \cdots \boxdot x_n)$$

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where the left and right sides are also parenthesized.

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where the left and right sides are also parenthesized.

Hence

 $a_1 = 1$ 

$$(\forall n \geq 2)[a_n = a_1a_{n-1} + a_2a_{n-2} + \dots + a_{n-1}a_1]$$

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 $a_0 = 0$  $a_1 = 1$ .

 $(\forall n \geq 2)[a_n = a_0a_n + a_1a_{n-1} + a_2a_{n-2} + \dots + a_{n-1}a_1 + a_na_0]$ 

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How to solve? Define

$$A(x) = \sum_{n=0}^{\infty} a_n x^i$$

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**Goal** Want A(x) as a function we can take the Taylor Series of to recover  $a_n$  as a formula.

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**Goal** Want A(x) as a function we can take the Taylor Series of to recover  $a_n$  as a formula.

A(x) is called the **Generating Function** of the sequence  $a_0, a_1, \ldots$ 

The recurrence only works when  $n \ge 2$ . Hence we look at

$$\sum_{n=2}^{\infty} a_n x^n = \sum_{n=2}^{\infty} (a_0 a_n + a_1 a_{n-1} + \dots + a_n a_0) x^n$$

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We get both sides to be a function of A(x).

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We get both sides to be a function of A(x).

$$\sum_{n=2}^{\infty} a_n x^n = \left(\sum_{n=0}^{\infty} a_n x^n\right) - a_1 x^1 - a_0 x^0 = A(x) - a_1 x - a_0 = A(x) - x$$

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$$\sum_{n=2}(a_0a_n+a_1a_{n-1}+\cdots+a_na_0)x^n$$

This looks a little like  $A(x)^2$ .

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This looks a little like  $A(x)^2$ . Next Slide.

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$$\sum_{n=2}^{\infty} (a_0 a_n + a_1 a_{n-1} + a_2 a_{n-2} + \dots + a_{n-1} a_1 + a_n a_0) x^n$$

$$A(x) = \sum_{n=0}^{\infty} a_n x^n$$

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$$A(x)^2 = \sum_{n=0}^{\infty} (a_0 a_n + a_1 a_{n-1} + \dots + a_n a_0) x^n$$

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$$\sum_{n=2}^{\infty} (a_0 a_n + a_1 a_{n-1} + a_2 a_{n-2} + \dots + a_{n-1} a_1 + a_n a_0) x^n$$

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$$A(x)^2 = \sum_{n=0}^{\infty} (a_0 a_n + a_1 a_{n-1} + \dots + a_n a_0) x^n$$

$$\sum_{n=2}^{\infty} (a_0 a_n + \dots + a_n a_0) x^n = A(x)^2 - (a_0 a_1 + a_1 a_0) x^1 - a_0^2 x^0 = A(x)^2$$

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## Equate LHS and RHS

$$A(x) - x = A(x)^2$$

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## Equate LHS and RHS

$$A(x) - x = A(x)^2$$

$$A(x)^2 - A(x) + x = 0$$

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## Equate LHS and RHS

$$A(x) - x = A(x)^2$$

$$A(x)^2 - A(x) + x = 0$$

$$A(x) = \frac{1 \pm \sqrt{1 - 4x}}{2}$$

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#### Do we Use + or -?

The Taylor Series for  $\sqrt{1-4x}$  is:



Do we Use + or -?

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$$\sqrt{1-4x} = \sum_{n=0}^{\infty} -\frac{2}{n} {\binom{2n-2}{n-1}} x^n$$

Do we Use + or -?

The Taylor Series for  $\sqrt{1-4x}$  is:

$$\sqrt{1-4x} = \sum_{n=0}^{\infty} -\frac{2}{n} {\binom{2n-2}{n-1}} x^n$$

The Taylor series has neg coeffs but we need pos coeffs. We take

$$-\sqrt{1-4x} = \sum_{n=0}^{\infty} \frac{2}{n} \binom{2n-2}{n-1} x^n$$

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 $A(x) = \sum_{n=0}^{\infty} a_n x^n$ 

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Algebra shows:

$$\frac{1 - \sqrt{1 - 4x}}{2} = \sum_{n=1}^{\infty} \frac{1}{n} \binom{2n - 2}{n - 1}.$$

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$$A(x) = \sum_{n=0}^{\infty} a_n x^n$$
$$A(x) = \frac{1 - \sqrt{1 - 4x}}{2}$$

Algebra shows:

$$\frac{1 - \sqrt{1 - 4x}}{2} = \sum_{n=1}^{\infty} \frac{1}{n} \binom{2n - 2}{n - 1}.$$

SO we have our answer!

$$a_n = \frac{1}{n} \binom{2n-2}{n-1}$$

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UNTIMED PART OF FINAL IS MONDAY May 10 9:00AM. NO DEAD CAT

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# FILL OUT COURSE EVALS for ALL YOUR COURSES!!!

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